## **Chapter 8**

# **Model-Based Design**



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## Design

Concepts: design process:

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requirements to **models** to implementations

Models: check properties of interest:

- safety on the appropriate (sub)system
- progress on the overall system

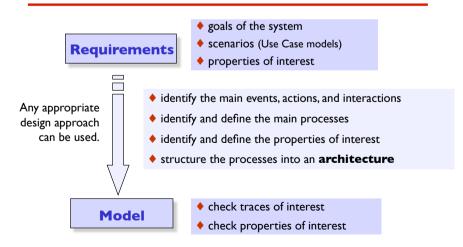
Practice: model interpretation - to infer actual system behavior

threads and monitors

Aim: rigorous design process.

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# 8.1 from requirements to models



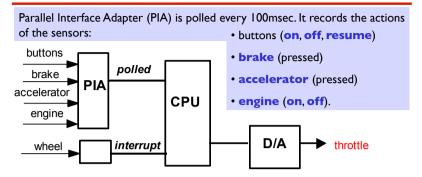
# a Cruise Control System - requirements



When the car ignition is switched on and the **on** button is pressed, the current speed is recorded and the system is enabled: it maintains the speed of the car at the recorded setting.

Pressing the brake, accelerator or off button disables the system. Pressing resume or on reenables the system.

## a Cruise Control System - hardware



Wheel revolution sensor generates interrupts to enable the car **speed** to be calculated.

Output: The cruise control system controls the car speed by setting the **throttle** via the digital-to-analogue converter.

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# model -design

♦ Main events, actions and interactions.

```
on, off, resume, brake, accelerator
engine on, engine off,
speed, setThrottle
clearSpeed, recordSpeed,
enableControl, disableControl

Prompts
```

Identify main processes.

Sensor Scan, Input Speed,
Cruise Controller, Speed Control and
Throttle

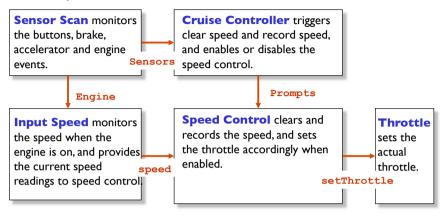
Identify main properties.

safety - disabled when off, brake or accelerator pressed.

Define and structure each process.

## model - outline design

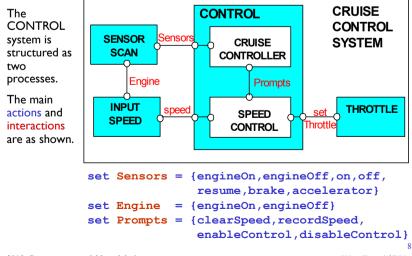
♦outline processes and interactions.



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# model - structure, actions and interactions



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## model elaboration - process definitions

```
SENSORSCAN = ({Sensors} -> SENSORSCAN).
     // monitor speed when engine on
INPUTSPEED = (engineOn -> CHECKSPEED),
CHECKSPEED = (speed -> CHECKSPEED
             |engineOff -> INPUTSPEED
             ١.
     // zoom when throttle set
THROTTLE = (setThrottle -> zoom -> THROTTLE).
     // perform speed control when enabled
SPEEDCONTROL = DISABLED,
DISABLED =({speed,clearSpeed,recordSpeed}->DISABLED
          | enableControl -> ENABLED
          ),
ENABLED = ( speed -> setThrottle -> ENABLED
          |{recordSpeed,enableControl} -> ENABLED
          | disableControl -> DISABLED
```

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## model - CONTROL subsystem

```
||CONTROL = (CRUISECONTROLLER
||SPEEDCONTROL
).
```

# Animate to check particular traces:

- Is control enabled after the engine is switched on and the on button is pressed?
- Is control disabled when the brake is then pressed?Is control re-enabled
- Is control re-enabled when resume is then pressed?

However, we need analysis to check exhaustively:

```
    Safety: Is the control disabled when off, brake or accelerator is pressed?
    Progress: Can every
```

action eventually be selected?

## model elaboration - process definitions

```
set DisableActions = {off,brake,accelerator}
     // enable speed control when cruising, disable when a disable action occurs
CRUISECONTROLLER = INACTIVE.
INACTIVE =(engineOn -> clearSpeed -> ACTIVE
          |DisableActions -> INACTIVE ),
ACTIVE
         =(engineOff -> INACTIVE
          |on->recordSpeed->enableControl->CRUISING
          |DisableActions -> ACTIVE ),
CRUISING = (engineOff -> INACTIVE
          |DisableActions->disableControl->STANDBY
          lon->recordSpeed->enableControl->CRUISING ),
STANDBY = (engineOff -> INACTIVE
          |resume -> enableControl -> CRUISING
          |on->recordSpeed->enableControl->CRUISING
           |DisableActions -> STANDBY
```

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### model - Safety properties

Safety checks are compositional. If there is no violation at a subsystem level, then there cannot be a violation when the subsystem is composed with other subsystems.

This is because, if the **ERROR** state of a particular safety property is unreachable in the LTS of the subsystem, it remains unreachable in any subsequent parallel composition which includes the subsystem. Hence...

Safety properties should be composed with the appropriate system or subsystem to which the property refers. In order that the property can check the actions in its alphabet, these actions must not be hidden in the system.

## model - Safety properties

```
property CRUISESAFETY =
  ({DisableActions, disableControl} -> CRUISESAFETY
  | {on,resume} -> SAFETYCHECK
 ),
SAFETYCHECK =
  ({on,resume} -> SAFETYCHECK
  |DisableActions -> SAFETYACTION
  |disableControl -> CRUISESAFETY
 ),
SAFETYACTION = (disableControl->CRUISESAFETY) .
                                                   LTS?
 | | CONTROL = (CRUISECONTROLLER
              | | SPEEDCONTROL
              | | CRUISESAFETY
              ) .
                        Is CRUISESAFETY violated?
```

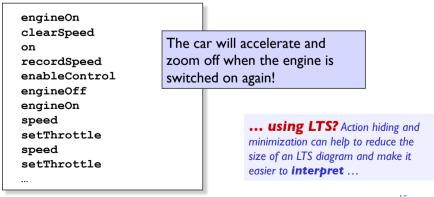
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# model - Safety properties

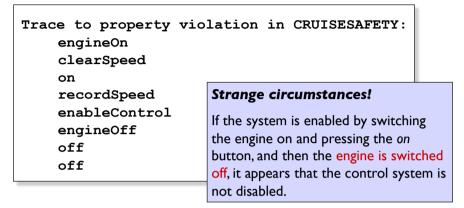
What if the engine is switched on again?

We can investigate further using animation ...



## model - Safety properties

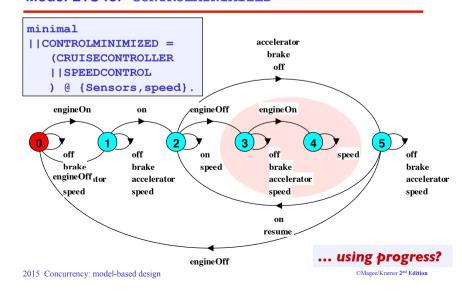
Safety analysis using LTSA produces the following violation:



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## Model LTS for CONTROLMINIMIZED

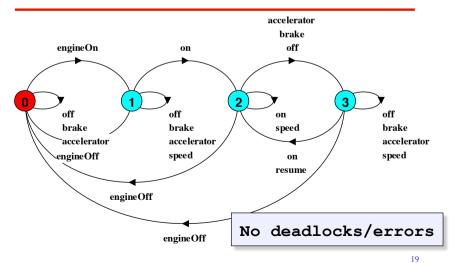


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## model - Progress properties

```
Progress violation for actions:
 {accelerator, brake, clearSpeed, disableControl,
enableControl, engineOff, engineOn, off, on,
recordSpeed, resume}
Trace to terminal set of states:
      engineOn
      clearSpeed
      on
      recordSpeed
      enableControl
                                       Check the model for
      engineOff
                                       progress properties with
      engineOn
Cycle in terminal set:
                                       no safety property and no
      speed
                                       hidden actions.
      setThrottle
Actions in terminal set:
      {setThrottle, speed}
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```

#### revised CONTROLMINIMIZED



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#### model - revised cruise controller

Modify CRUISECONTROLLER so that control is disabled when the engine is switched off:

```
CRUISING =(engineOff -> disableControl -> INACTIVE
          |DisableActions -> disableControl -> STANDBY
          |on->recordSpeed->enableControl->CRUISING
          ),
```

Modify the safety property:

```
property IMPROVEDSAFETY =
  {DisableActions, disableControl, engineOff} -> IMPROVEDSAFETY
 |{on,resume} -> SAFETYCHECK
 ),
SAFETYCHECK = ({on.resume} -> SAFETYCHECK
              |{DisableActions.engineOff} -> SAFETYACTION
              |disableControl
                                 -> IMPROVEDSAFETY
              ),
SAFETYACTION = (disableControl
                                 -> IMPROVEDSAFETY).
                                                     OK now?
```

## model analysis

We can now proceed to compose the whole system:

```
|| CONTROL =
   (CRUISECONTROLLER | | SPEEDCONTROL | | CRUISESAFETY
   )@ {Sensors, speed, setThrottle}.
||CRUISECONTROLSYSTEM =
     (CONTROL | | SENSORSCAN | | INPUTSPEED | | THROTTLE)
```

# **Deadlock?** Safety?

No deadlocks/errors

# **Progress?**

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## model - Progress properties

Progress checks are not compositional. Even if there is no violation at a subsystem level, there may still be a violation when the subsystem is composed with other subsystems.

This is because an action in the subsystem may satisfy progress yet be unreachable when the subsystem is composed with other subsystems which constrain its behaviour. Hence...

Progress checks should be conducted on the complete target system after satisfactory completion of the safety checks.

**Progress?** 

No progress violations detected.

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## model interpretation

Models can be used to indicate system sensitivities.

If it is possible that erroneous situations detected in the model may occur in the implemented system, then the model should be revised to find a design which ensures that those violations are avoided.

However, if it is considered that the real system will not exhibit this behavior, then no further model revisions are necessary.

Model interpretation and correspondence to the implementation are important in determining the relevance and adequacy of the model design and its analysis.

# model - system sensitivities

What about progress under **adverse** conditions? Check for system sensitivities.

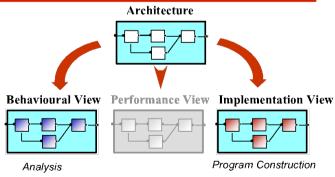
| | SPEEDHIGH = CRUISECONTROLSYSTEM << {speed}.

Progress violation for actions:
 {accelerator, brake, engineOff, engineOn, off, on, resume, setThrottle, zoom}
Trace to terminal set of states:
 engineOn
Cycle in terminal set:
 speed
Actions in terminal set:
 speed
The system may be sensitive to the priority of the action speed.

## The central role of design architecture

Design architecture describes the gross organization and global structure of the system in terms of its constituent components.

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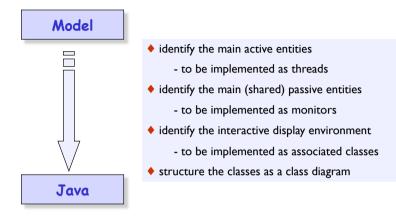


We consider that the models for analysis and the implementation should be considered as elaborated views of this basic design structure.

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## 8.2 from models to implementations



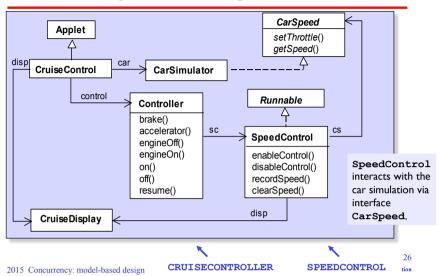
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## cruise control system - class Controller

```
class Controller {
  final static int INACTIVE = 0; // cruise controller states
  final static int ACTIVE = 1;
                                                       Controller
  final static int CRUISING = 2;
                                                       is a passive
  final static int STANDBY = 3;
                                                       entity - it
  private int controlState = INACTIVE: //initial state
  private SpeedControl sc;
                                                       reacts to
  Controller (CarSpeed cs, CruiseDisplay disp)
                                                       events. Hence
    {sc=new SpeedControl(cs,disp);}
                                                       we implement
  synchronized void brake(){
                                                       it as a
    if (controlState==CRUISING )
                                                       monitor
      {sc.disableControl(); controlState=STANDBY; }
  synchronized void accelerator(){
    if (controlState==CRUISING )
      {sc.disableControl(); controlState=STANDBY; }
synchronized void engineOff() {
    if (controlState!=INACTIVE) {
      if (controlState==CRUISING) sc.disableControl();
      controlState=INACTIVE;
  }
```

# cruise control system - class diagram



## cruise control system - class Controller

```
synchronized void engineOn(){
    if (controlState==INACTIVE)
      {sc.clearSpeed(); controlState=ACTIVE;}
                                                       This is a direct
                                                       translation
 synchronized void on(){
                                                       from the
    if (controlState!=INACTIVE) {
                                                       model.
      sc.recordSpeed(); sc.enableControl();
      controlState=CRUISING;
 }
  synchronized void off() {
    if(controlState==CRUISING)
      {sc.disableControl(); controlState=STANDBY;}
 synchronized void resume(){
    if(controlState==STANDBY)
     {sc.enableControl(); controlState=CRUISING;}
 }
```

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#### cruise control system - class SpeedControl

```
class SpeedControl implements Runnable {
                                                      SpeedControl
  final static int DISABLED = 0: //speed control states
  final static int ENABLED = 1:
                                                      is an active entity
  private int state = DISABLED; //initial state
                                                       - when enabled, a
 private int setSpeed = 0:
                                   //target speed
 private Thread speedController;
                                                      new thread is
  private CarSpeed cs;
                              //interface to control speed
                                                      created which
 private CruiseDisplay disp;
                                                      periodically
  SpeedControl(CarSpeed cs, CruiseDisplay disp){
                                                      obtains car speed
    this.cs=cs; this.disp=disp;
    disp.disable(); disp.record(0);
                                                      and sets the
                                                      throttle
  synchronized void recordSpeed(){
    setSpeed=cs.getSpeed(); disp.record(setSpeed);
  synchronized void clearSpeed() {
    if (state==DISABLED) {setSpeed=0;disp.record(setSpeed);}
  synchronized void enableControl(){
    if (state==DISABLED) {
      disp.enable(); speedController= new Thread(this);
      speedController.start(); state=ENABLED;
```

## **Summary**

- Concepts
  - design process:

from requirements to models to implementations

- design architecture
- Models
  - check properties of interest
     safety: compose safety properties at appropriate (sub)system
     progress: apply progress check on the final target system model
- Practice
  - model interpretation to infer actual system behavior
  - threads and monitors

Aim: rigorous design process.

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## cruise control system - class SpeedControl

```
synchronized void disableControl(){
   if (state==ENABLED) { disp.disable(); state=DISABLED;}
}

public void run() { // the speed controller thread
   try {
    while (state==ENABLED) {
        double error = (float) (setSpeed-cs.getSpeed())/6.0;
        double steady = (double) setSpeed/12.0;
        cs.setThrottle(steady+error);//simplified feed back control
        wait(500);
    }
    } catch (InterruptedException e) {}
    speedController=null;
}
```

**SpeedControl** is an example of a class that combines both synchronized access methods (to update local variables) and a thread.

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#### **Course Outline**

- 2. Processes and Threads
- 3. Concurrent Execution
- 4. Shared Objects & Interference
- . Monitors & Condition Synchronization
- 6. Deadlock
- 7. Safety and Liveness Properties
- 8. Model-based Design

The main basic

Concepts

Models

**Practice** 

#### Advanced topics ...

- 9. Dynamic systems
- 10. Message Passing
- II. Concurrent Software Architectures
- 12. Timed Systems
- 13. Program Verification
- 14. Logical Properties

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