



The AutoMed Query Processor



Previous Architecture

- n Query reformulator: reformulates the input query q by following the transformation pathways from the GS to LS_i
- n Fragment processor: replaces each scheme s with a wrapper containing s
- n Evaluator: evaluates q



New Architecture

- n Query reformulator: same functionality
- n Logical optimiser: performs various logical optimisations
- n Query annotator: detects the largest subtrees t_i translatable by the datasource wrappers and inserts a wrapper object as the root of t_i
- n Physical Optimiser: performs datasource specific optimisations
- n Evaluator: same functionality



Logical Optimisations

- n Rule application using templates
 - n Disjunction optimiser – see DBIS'04
 - n Nil optimiser – see DBIS'04
 - n Comprehensions optimiser. Example:
$$[\{h\}|q_1; \{x\}\mathcal{B}(DS_1 ++ DS_2); q_2] \dot{\rightarrow}$$
$$[\{h\}|q_1; \{x\}\mathcal{B}DS_1; q_2] ++ [\{h\}|q_1; \{x\}\mathcal{B}DS_2; q_2]$$
- n Java code to modify ASG structure
 - n Datasource reorganiser
 - n Common sub-expression elimination



Datasource Reorganiser

n Java-based optimisation

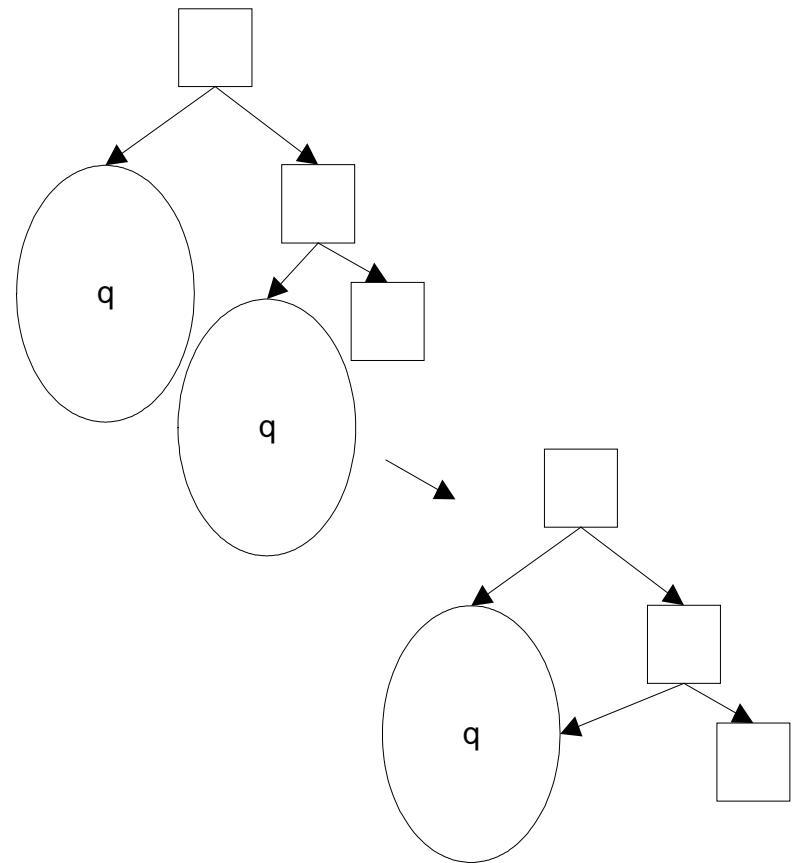
n Examples:

n $DS_1:A ++ DS_2:B ++ DS_1:C \rightarrow$
 $(DS_1:A ++ DS_1:C) ++ DS_2:B$

n $[\{h\} | DS_1:A; DS_2:C; DS_1:B; DS_2:D; p_1; p_2; p_{12}]$
 \rightarrow
 $[\{h\} | \{h_1\} \mathbf{B} [\{h_1\} | DS_1:A; DS_1:B; p_1];$
 $\{h_2\} \mathbf{B} [\{h_2\} | DS_2:C; DS_2:D; p_2];$
 $p_{12}]$

Common Sub-Expression Elimination

- n Input query may contain multiple identical subqueries
- n Transform input query into a DAG to avoid evaluation of the same subquery





Logical Optimiser

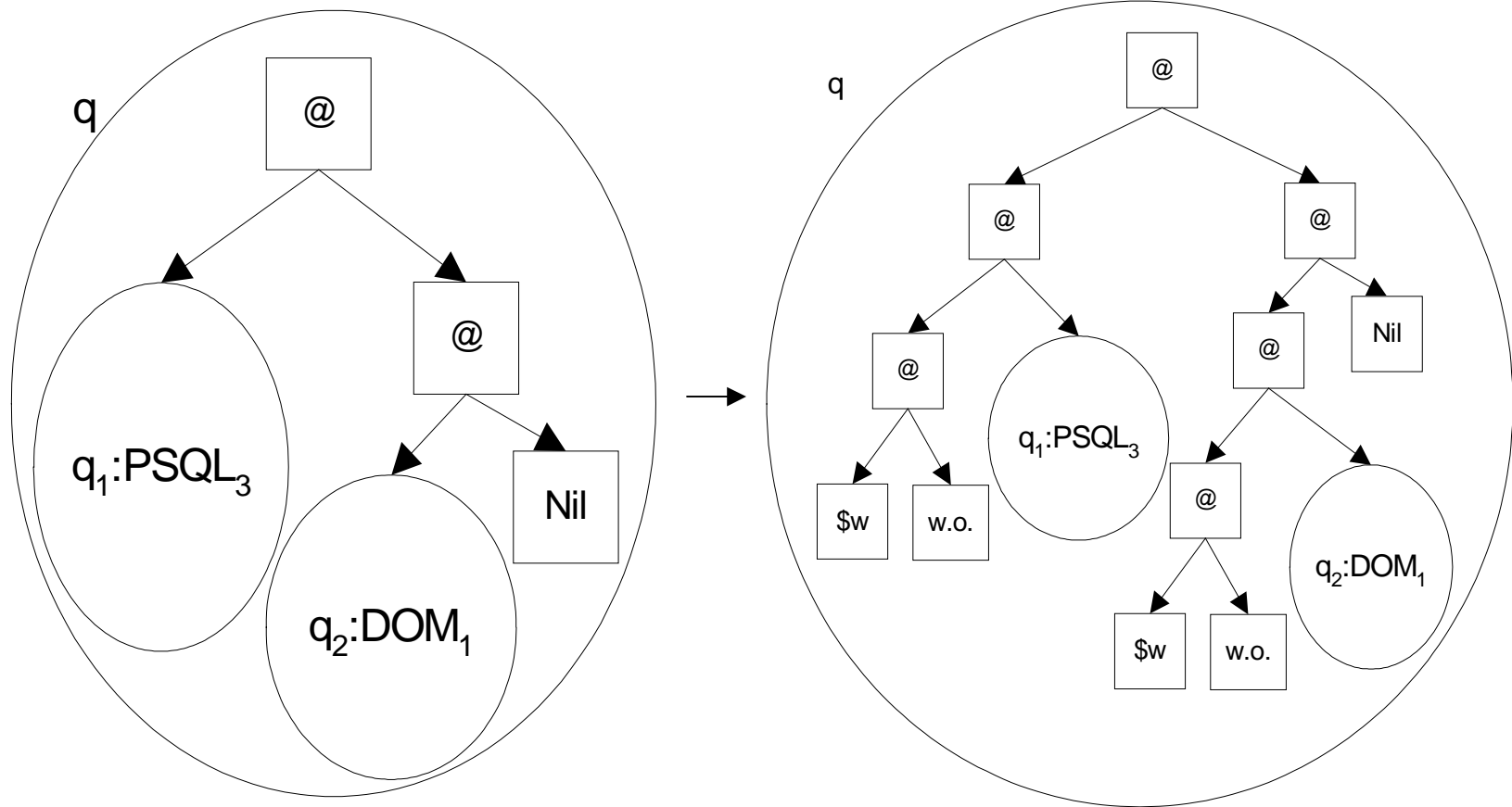
- n Applies logical optimisations using the following policy:
 - n Step 1: apply each logical optimisation until an application does not modify the query
 - n Step 2: apply step 1 until the query is not modified by any logical optimisation
 - n Step 3: apply common sub-expression elimination



Query Annotator

- n Detects the largest subqueries which can be translated by the datasource wrappers
- n Once the annotator detects a translatable subtree, it inserts a wrapper object

Query Annotator - Example





Query Annotator – Detection

- n Each wrapper is capable of translating a subset of IQL
- n Each subset of IQL is represented in the query processor by a parser p
- n When a query q is submitted to p , if it is not part of this subset of IQL, a syntax error is thrown
- n Each wrapper defines the subset of IQL it can translate by selecting a parser; if no parser is selected, the default parser is used



Query Annotator – Detection

- n Each Cell in the input query defines a subtree t
- n The Query Annotator traverses the input query once for every datasource wrapper w and for every t checks whether it is translatable by w



Physical Optimiser

- n Currently consists of a single optimiser, the dual model optimiser:
 - n Some datasources are modeled using two modeling layers: datasource-oriented & AutoMed-oriented à may cause unnecessary self-joins of schema constructs:

Original:

$[\{id,name\} | \{i,id\} \mathcal{B} \langle \langle person, id \rangle \rangle ; \{i,name\} \mathcal{B} \langle \langle person, dname \rangle \rangle]$

Reformulated:

$[\{id,name\} | \{i,id\} \mathcal{B} [\{k1,k1\} | \{k1,a1,a2\} \mathcal{B} \langle \langle person, 3 \rangle \rangle] ;$
 $\{i,name\} \mathcal{B} [\{k1,a2\} | \{k1,a1,a2\} \mathcal{B} \langle \langle person, 3 \rangle \rangle]]$

Optimised: $[\{id,name\} | \{k,id,name\} \mathcal{B} \langle \langle person, 3 \rangle \rangle]$



External Functions

- n Currently, the Evaluator handles only built-in and user-defined functions – all written in Java
- n We are currently extending the evaluator to use external functions written in other programming languages, such as C, C++ and Perl



Lazy Evaluation

- n The Evaluator currently fully evaluates a query submitted to it.
- n This may be inadequate for queries that return large result sets which do not fit into the available memory.
- n Thus, once the above functionality is fully implemented and tested, we will modify the evaluator to incrementally evaluate queries and to return fragments of result sets.



Type System – Type Checker

- n Devise a unified type system for AutoMed
- n Implement type checker which type checks input queries and queries supplied with transformations



Open Issues

- n Dual model optimiser
- n External functions
- n Lazy evaluation
- n Type system – type checker