#### Introducing myself



- Heterogeneous data management.
- Imperfect data management.

#### Dealing with Imperfection in Schema Integration Some preliminary ideas

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#### Outline

Imperfect data: an overview.

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- Imperfect data: an overview.
- Uncertainty in data integration based on semantic schema matching.
  - Semantic Schema Integration.
  - Adding uncertainty.

#### **Imperfect Data: An Overview**

## A taxonomy of imperfection

To identify the main classes of imperfection, we use the following scenario:



#### (Complete) Absence



#### Absence

- All our belief is committed to the known set of alternatives.
- For example: "We do not know the age of John".



## **Imprecision: Non-specificity**

- Imprecision concerns the cardinality of our believes.
- When we believe in a crisp set, imprecision is called non-specificity.
- For example: "John is between 170 and 180 cm. tall".



#### **Imprecision: Vagueness**

- If a set representing our belief is not crisp, imprecision is called vagueness.
- For example: "John is not very tall"



#### Uncertainty

- We have uncertainty when we do not commit all our belief.
- For example, "John should be 27 y.o.".



Туре	Abbr	Example (John's tallness)	
Absence	ABS	Not known.	
Non-Specificity	NS	Between 180 and 190 cm.	
		183 or 187 cm.	
Vagueness	VAG	Not very tall.	
Uncertainty	UN	Perhaps, 183 cm.	

#### Uncertainty in Schema Integration based on Semantic Schema Matching

with Nikos Rizopoulos

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=  $\cap$   $\subset$   $\supset$   $\cup$   $\neq$ 

#### SEMANTIC RELATIONSHIPS







=  $\cap$   $\subset$   $\supset$   $\cup$   $\neq$ 



#### **Dealing with uncertainty**



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#### **Dealing with uncertainty**





Identify uncertain relationships.

- Identify uncertain relationships.
- Produce uncertain partial integrated schemas.

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- Produce uncertain partial integrated schemas.
- Put together the uncertain partial integrated schemas, to obtain an uncertain integrated schema.
- Query the database.

- Dempster–Shafer's theory to represent believes in relationships.

$$\Theta = \{=, \cap, \subset, \supset, \cup, \neq\}$$

- $m(\{=\}) = .4$
- $m(\{\supset\}) = .3$
- $m(\{ \subset \}) = .3$

$$\Theta = \{=, \cap, \subset, \supset, \cup, \neq\}$$

• 
$$m(\{\cup, \neq\}) = 1$$

$$\Theta = \{=, \cap, \subset, \supset, \cup, \neq\}$$

• 
$$m(\{=\}) = .2$$

• 
$$m(\Theta) = .8$$

$$\Theta = \{=, \cap, \subset, \supset, \cup, \neq\}$$

• 
$$m(\Theta) = 1$$











- Scalable.
- Experts can be software agents or humans.

 $\mathbf{Exp}_1$  (Cardinality):

 $\ \, \blacksquare \ \, m(\{\cap,\supset,\cup,\neq\})=1$ 

**Exp**<sub>1</sub> (Cardinality):  $m(\{\cap, \supset, \cup, \neq\}) = 1$  **Exp**<sub>2</sub> (Thesaurus):  $m(\{\supset\}) = .5$  $m(\{=\}) = .2$ 

• 
$$m(\Theta) = .3$$

**Exp**<sub>1</sub> (Cardinality):  $m(\{\cap, \supset, \cup, \neq\}) = 1$  **Exp**<sub>2</sub> (Thesaurus):  $m(\{\supset\}) = .5$   $m(\{=\}) = .2$  $m(\Theta) = .3$ 

**Exp** $_3$  (Human):

● 
$$m(\{=, \supset, \cup\}) = 1$$

**Exp**<sub>1</sub> (Cardinality):  $m(\{\cap, \supset, \cup, \neq\}) = 1$  **Exp**<sub>2</sub> (Thesaurus):  $m(\{\supset\}) = .5$   $m(\{=\}) = .2$  $m(\Theta) = .3$ 

**Exp** $_3$  (Human):

- $m(\{=, \supset, \cup\}) = 1$
- $\oplus$  (Combination):
  - $m(\{ \supset \}) = \frac{5}{8}$
  - $\ \, {} \quad m(\{\supset,\cup\})=\tfrac{3}{8}$

#### **Unc. partial integrated schemas**

Given two objects, some belief is committed to each possible relationship between them.

Bel: =  $\cap$   $\subset$   $\supset$   $\cup$   $\neq$ 

#### **Unc. partial integrated schemas**

- Given two objects, some belief is committed to each possible relationship between them.
- From each possible relationship we can obtain a partial integrated schema.

#### Bel: = $\cap$ $\subset$ $\supset$ $\cup$ $\neq$ $\downarrow$ $\downarrow$ $\downarrow$ $\downarrow$ $\downarrow$ $\downarrow$ $\downarrow$ $\downarrow$ PI(=) PI(\cap) PI(\subset) PI(\supset) PI(\cup) PI(\neq)

#### **Unc. partial integrated schemas**

- Given two objects, some belief is committed to each possible relationship between them.
- From each possible relationship we can obtain a partial integrated schema.
- The belief committed to each partial integrated schema is the same previously committed to the corresponding possible relationship.



Uncertain relationship (Student-UG):

• 
$$m(\{\supset\}) = \frac{5}{8}$$

• 
$$m(\{\supset,\cup\}) = \frac{3}{8}$$

Uncertain partial integrated schema:

• 
$$m({\operatorname{PI}(\supset)}) = \frac{5}{8}$$

• 
$$m({\operatorname{PI}(\supset), \operatorname{PI}(\cup)}) = \frac{3}{8}$$



Main idea: take all possible combinations of uncertain partial integrated schemas.

$$= (A,B), \neq (A,B)$$

 $= (B, C), \neq (B, C)$ 



Some issues:

Compact representation.

Dependencies.



Uncertain relationship (Student–UG):

• 
$$m(\{\supset\}) = \frac{5}{8}, m(\{\supset,\cup\}) = \frac{3}{8}$$

Uncertain relationship (res-res):

• 
$$m(\{=\}) = \frac{1}{3}, m(\{\neq\}) = \frac{2}{3}$$



Uncertain partial integrated schema:

- $m(\{\mathsf{PI}(\supset)\} \times \{\mathsf{PI}(=)\}) = \frac{5}{24}$
- $m(\{\mathsf{PI}(\supset),\mathsf{PI}(\cup)\}\times\{\mathsf{PI}(=)\})=\frac{3}{24}$
- $m(\{\mathsf{PI}(\supset)\} \times \{\mathsf{PI}(\neq)\}) = \frac{5}{12}$
- $m(\{\mathsf{PI}(\supset),\mathsf{PI}(\cup)\} \times \{\mathsf{PI}(\neq)\}) = \frac{3}{12}$

- $= (A, B), \neq (A, B)$
- $= (B,C), \neq (B,C)$
- $= (A, C), \neq (A, C)$

- $= (A, B), \neq (A, B)$
- $= (B,C), \neq (B,C)$
- $= (A, C), \neq (A, C)$

A, B	B, C	A, C	Allowed
=	=	=	
=		$\neq$	
=	$\neq$		
$\neq$			
$\neq$		$\neq$	
=	$\neq$	$\neq$	
$\neq$	$\neq$		
$\neq$	$\neq$	$\neq$	

- $= (A, B), \neq (A, B)$
- $= (B,C), \neq (B,C)$
- $= (A, C), \neq (A, C)$

A, B	B, C	A, C	Allowed
=	=	=	Y
=	=	$\neq$	Ν
=	$\neq$	=	Ν
$\neq$	_	_	Ν
$\neq$	=	$\neq$	Y
=	$\neq$	$\neq$	Y
$\neq$	$\neq$	=	Y
$\neq$	$\neq$	$\neq$	Y

#### **Querying the database**

- No idea...
- It should not be very difficult to define the semantics of a query.
- Efficiency problems.
  - Cardinality reduction.
  - Compact query plans.

#### **Concluding remarks**

- Uncertainty is one of many possible types of imperfection/ignorance.
- We start our investigation from a method of schema integration based on semantic schema matching.
- In real cases of data integration, it can be difficult to identify certain semantic relationships.
- We have presented some preliminary ideas on how to extend this method to deal with uncertainty.

#### **Discussion**

# Dealing with Imperfection in Schema Integration

#### Some preliminary ideas

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