

Kruskal's theorem and Nash-Williams theory

Ian Hodkinson, after Wilfrid Hodges

Version 3.6, 7 February 2003

This is based on notes I took at Wilfrid's seminar for Ph.D. students at Queen Mary College, around autumn 1985. His seminars drew to some extent on [1, 3, 4, 5].

I started this latex version of the notes in about March 1992, and revised and expanded it beyond the material in the 1985 notes in Jan 2003. Section 6 is still incomplete. Some of the improvements may be bad improvements, and there are surely still some errors due to my bad note-taking and transcribing. Use at your own risk.

I strongly recommend and request that *you do not cite these notes* in your own publications. The notes are made available because they may be helpful, but they are not authoritative or formally published. The original sources are [1, 3, 4] and these should be cited in preference. Any results given here that are not in these papers are essentially just exercises.

Thanks to Sz. Mikulás for helpful comments.

1 Well-quasi-orderings

Definition 1.1 *If S is any set and κ a cardinal, $[S]^\kappa$ is the set of subsets of S of size κ , and $[S]^{<\kappa}$ is the set of subsets of S of size less than κ .*

Definition 1.2 *A pre-order is a reflexive and transitive binary relation, usually written \leq, \sqsubseteq , etc. Given a pre-order \leq , we write $x < y$ to abbreviate $x \leq y \wedge y \not\leq x$.*

Definition 1.3 (wqo) *A well-quasi-ordering (wqo) is a pre-order such that (i) it is well-founded (it has no infinite strictly descending ($>$) sequences), and (ii) there is no infinite antichain (a set of pairwise incomparable elements).*

Lemma 1.4 *If \leq is a pre-order on a set I , then the following are equivalent:*

1. \leq is a wqo.
2. If $x_0, x_1, \dots \in I$ then there are $i < j$ with $x_i \leq x_j$.
3. If $x_0, x_1, \dots \in I$ then there is an infinite $X \subseteq \mathbb{N}$ such that $x_i \leq x_j$ for all $i < j$ in X .
4. Any $X \subseteq I$ has a finite subset of minimal elements: there is finite $Y \subseteq X$ such that $\forall x \in X \exists y \in Y (y \leq x)$.

Proof. For the equivalence of 1–3, it is enough to show $1 \Rightarrow 3$. Assume (1). Let $f : [\mathbb{N}]^2 \rightarrow \{\leq, >, \perp\}$ be such that $x_i f(i, j) x_j$ for all $i < j < \omega$. Here, $a \perp b$ means a and b are incomparable. By Ramsey's theorem (see corollary 4.4 below), there is infinite $X \subseteq \mathbb{N}$ such that $f \upharpoonright [X]^2$ is constant. Assuming (1), the constant value is not \perp or $>$; hence it is \leq , and we are done.

Clearly, $4 \Rightarrow 2$. To show $2 \Rightarrow 4$, assume $X \subseteq I$ fails (4). Define $x_0, x_1, \dots \in X$ by induction: given $n < \omega$, since $\{x_0, \dots, x_{n-1}\}$ is finite and X fails (4), there is $x_n \in X$ with $x_i \not\leq x_n$ for all $i < n$. Then the sequence x_0, x_1, \dots fails (2). \square

This is reminiscent of König's tree lemma: well-founded replaces the tree condition, and the lack of infinite antichains replaces the finitely-branching condition. The lemma's conclusion, that any infinite sequence of distinct elements contains an infinite increasing subsequence, corresponds to the infinite branch that König's lemma provides.

Now we see a little Nash-Williams theory.

Definition 1.5 *Given a pre-order \leq on I , a bad sequence is an infinite sequence x_0, x_1, \dots in I such that $x_i \not\leq x_j$ whenever $i < j$. We say such a sequence is minimal bad if it is bad, and for each n there is no bad sequence $x_0, \dots, x_{n-1}, y_n, y_{n+1}, \dots$ with $y_n < x_n$.*

By the lemma, a wqo is one without bad sequences. Any infinite subsequence of a bad sequence is bad.

Lemma 1.6 *If \leq is well-founded but is not a wqo then there is a minimal bad sequence.*

Proof. By induction. There are bad sequences. Choose $x_0 \in I$ minimal such that it is the first element of a bad sequence — this is possible as \leq is well-founded. Inductively, if x_i ($i < n$) are chosen, choose x_n such that it is minimal such that x_0, \dots, x_n extends to a bad sequence. Clearly the result x_0, x_1, \dots is a minimal bad sequence. \square

Lemma 1.7 *Let x_0, x_1, \dots be minimal bad. Put $Y = \{x \in I : x < x_i \text{ for some } i\}$. Then \leq is a wqo on Y .*

Proof. If not, then there is a bad sequence y_0, y_1, \dots in Y . Each y_i is $<$ some x_i . Choose $i < \omega$ to make i' as small as possible. By throwing away y_0, \dots, y_{i-1} we can assume $i = 0$.

Claim: $x_0, \dots, x_{i'-1}, y_0, y_1, \dots$ is bad.

Proof of Claim: If not, then as the x -part and y -part are from bad sequences, we must have $x_m \leq y_n$ for some $m < i'$ and some n . Now by definition of Y , we have $y_n < x_{n'}$ for some n' . By choice of i' as small as possible, $n' \geq i'$. So $x_m \leq y_n < x_{n'}$ and $m < n'$, contradicting x_0, x_1, \dots being bad. This proves the claim.

The claim contradicts the choice of x_0, x_1, \dots as minimal bad. Thus the lemma is proved. \square

Definition 1.8 *Given a pre-order (I, \leq) , define $I^{<\omega}$ to be the set of all finite sequences of elements of I , ordered by: $(x_0, \dots, x_{m-1}) \leq (y_0, \dots, y_{n-1})$ iff there is a one-one order-preserving map $f : m \rightarrow n$ such that $x_i \leq y_{f(i)}$ for all $i < m$.*

Lemma 1.9 *The relation \leq on $I^{<\omega}$ is a pre-order, and if $<$ is well-founded on I then $<$ is well-founded on $I^{<\omega}$.*

Proof. Reflexivity and transitivity of \leq on $I^{<\omega}$ are clear. If there is an infinite strictly decreasing chain \bar{x}_i ($i < \omega$) in $I^{<\omega}$, then as the lengths of the \bar{x}_i must form a non-increasing sequence of natural numbers, we can assume the \bar{x}_i all have the same length, m say. For each i there is $j < m$ such that $x_{i+1,j} < x_{ij}$. Hence there is j such that the above holds for infinitely many i . Thus $(x_{ij} : i < \omega)$ contains a strictly decreasing subsequence, contradicting the well-foundedness of $<$ on I . \square

Theorem 1.10 (Higman) *If \leq is a wqo on I , then \leq is a wqo on $I^{<\omega}$.*

Proof. By lemma 1.9, \leq is well-founded. If the theorem fails, then by lemma 1.9 there is a minimal bad sequence $\bar{x}_0, \bar{x}_1, \dots$ of sequences $\bar{x}_i = (x_{i,0}, \dots, x_{i,j_i-1})$. Now the null sequence is vacuously \leq any sequence, so each \bar{x}_i has length ≥ 1 . So we can write it as $(x_{i,0}, \text{tail}(\bar{x}_i))$. As $\text{tail}(\bar{x}_i)$ has length $< \text{len}(\bar{x}_i)$, we have $\text{tail}(\bar{x}_i) < \bar{x}_i$ in $I^{<\omega}$ (note the strict $<$). So by lemma 1.7, $Y = \{\text{tail}(\bar{x}_i) : i < \omega\}$ is a subset of a well-quasi-ordered set, and so itself well-quasi-ordered. So by lemma 1.4 we can assume that $\text{tail}(\bar{x}_0) \leq \text{tail}(\bar{x}_1) \leq \dots$ (This might not preserve *minimal* badness, but minimality is no longer needed.) Also, I is a wqo, so there are $i < j$ with $x_{i,0} \leq x_{j,0}$. Hence by “piecing together”, $\bar{x}_i \leq \bar{x}_j$, a contradiction. \square

Exercise 1.11 Let (I, \leq) be a wqo. Define a relation \leq on the set $\wp_{<\omega}(I)$ of finite subsets of I , by $s \leq t$ iff there is a map $f : s \rightarrow t$ with $i \leq f(i)$ for all $i \in s$. Show that $(\wp_{<\omega}(I), \leq)$ is a wqo (cf. [1, p. 32]). Repeat for the variant of \leq where f is required to be one-one.

Remark 1.12 (Rado) *Theorem 1.10 fails for the set $I^{<\omega_1}$, the set of countable sequences of elements of I .*

Proof. Let $I = \{(i, j) : i < j < \omega\}$, ordered by $(i, j) \leq (k, l)$ iff either $i = k$ and $j \leq l$, or else $i, j < k$. One can check that this is a wqo on I . Now for $i < \omega$ let α_i be the sequence $((i, i + 1), (i, i + 2), \dots)$. Then for all $i < j < \omega$, $\alpha_i \not\leq \alpha_j$. So the sequence $(\alpha_i : i < \omega)$ is bad. \square

Exercise 1.13 Show that exercise 1.11 fails for the full power set $\wp(I)$.

Exercise 1.14 (cf. exercise 1.11) Let (I, \leq) be a wqo. Define a relation \preceq on $\wp_{<\omega}(I)$ by $s \preceq t$ iff there is a map $g : t \rightarrow s$ with $g(i) \leq i$ for all $i \in t$. Find (I, \leq) such that $(\wp_{<\omega}(I), \preceq)$ is not a wqo. [I think this non-preservation is well-known but thanks to Sz. Mikulás for pointing this example out to me.]

2 Kruskal's theorem

Definition 2.1 *A tree is a finite connected graph without cycles. Let T, S be trees. We say that T immediately yields S if a tree isomorphic to S can be got from T by either removing one "leaf" vertex and its only attaching edge, or turning (a) into (b) below:*



We write $S > T$ if there are $n > 1$ and trees T_1, \dots, T_n such that $S \cong T_1, T_n \cong T$ and T_i immediately yields T_{i+1} . We write $S \leq T$ if $S < T$ or $S \cong T$.

Note that $|S| < |T|$ if $S < T$. Hence $<$ is well-founded on finite trees. Question: is there an infinite sequence of trees T_0, T_1, \dots such that if $i < j$ then $T_i \not\leq T_j$ (i.e., a bad sequence)? Kruskal answered "no". We'll prove this.

Definition 2.2 *A pointed tree is a tree in which some vertex is distinguished as the "root". A tree is then partially ordered by: $x \leq y$ if the path from the root to y passes through x — x is nearer to the root than y . For any vertices*

x, y there's a greatest lower bound $x \wedge y$, the furthest node from the root on both paths to x and to y .

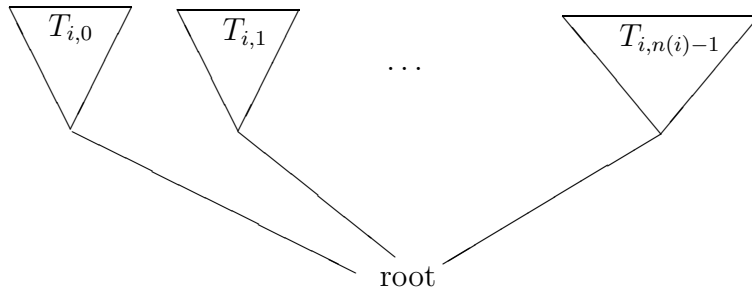
A decent embedding of pointed trees T, S is a 1-1 map $f : T \rightarrow S$ satisfying:

1. $u \leq v \Rightarrow f(u) \leq f(v)$
2. $f(u) \wedge f(v) = f(u \wedge v)$.

One can show that if there's a decent embedding $f : T \rightarrow S$ then $T \leq S$ as non-pointed trees. So it is enough to show that if we pre-order the class of pointed trees by $T \leq S$ iff there's a decent embedding from T to S , then \leq is a wqo.

Theorem 2.3 *The pointed trees are a wqo under this ordering.*

Proof. As trees are finite, it's clearly well-founded. If it's not a wqo, take a minimal bad sequence of trees, T_0, T_1, \dots . Since the sequence is bad, each T_i has at least two vertices (as the 1-point tree embeds decently into any pointed tree). So T_i can be written as in the diagram below.



Each T_{ij} is $< T_i$ (use the inclusion map; " \leq " fails because all trees are finite), and so $Y = \{T_{ij} : i < \omega, j < n(i)\}$ is a wqo by lemma 1.7. Hence by Higman's theorem, $Y^{<\omega}$ is also a wqo. Now replace T_i by the finite sequence $(T_{i,0}, \dots, T_{i,n(i)-1})$. We deduce that there's $i < j$ such that $(T_{i,0}, \dots, T_{i,n(i)-1}) \leq (T_{j,0}, \dots, T_{j,n(j)-1})$. We can now construct a decent embedding $T_i \rightarrow T_j$, contradicting badness. \square

Kruskal's theorem is now proved. The argument seems to me to be essentially to be proving that a sequence T_0, T_1, \dots is not bad by induction on the "slope" of a sequence T_0, T_1, \dots of trees, where T_0, T_1, \dots has slope $< U_0, U_1, \dots$ iff $T_i < U_i$ for all i . This ordering is well-founded in the sense that non-empty "closed" sets of sequences (e.g., the bad sequences) have minimal elements. A tree is shown to be decomposable into a finite sequence of smaller pieces,

in such a way that a decent embedding of sequences can be patched together to give a decent embedding of the original trees. Higman's theorem is used to extend the inductive hypothesis to *sequences* of smaller objects.

3 Prikry pairs and Ramsey sets

This section is central to what follows.

Definition 3.1

1. We say a set $\mathcal{S} \subseteq [\omega]^\omega$ is Ramsey if there is some infinite set $M \subseteq \omega$ such that either $[M]^\omega \subseteq \mathcal{S}$ or $[M]^\omega \cap \mathcal{S} = \emptyset$.
2. We say $\mathcal{S} \subseteq [\omega]^\omega$ is open if for every $X \in \mathcal{S}$ there is $n < \omega$ such that whenever $Y \subseteq \omega$ is infinite and $Y \cap n = X \cap n$ then $Y \in \mathcal{S}$.

Remark 3.2 $[\omega]^\omega$ is Ramsey. We can use the axiom of choice and diagonalisation to make a non-Ramsey set. $\mathcal{S} \subseteq [\omega]^\omega$ is open if whenever $X \in \mathcal{S}$ then some finite part of X puts it in \mathcal{S} .

Theorem 3.3 (Galvin, Prikry) *Every open set is Ramsey. (In fact, Borel sets are Ramsey.)*

Fix $\mathcal{S} \subseteq [\omega]^\omega$. First a definition.

Definition 3.4

1. If $X, Y \subseteq \omega$ we write $X < Y$ if for all $x \in X$ and $y \in Y$ we have $x < y$.
2. We say that an infinite set $M \subseteq \omega$ accepts a finite set $X \subseteq \omega$ if $X \cup N \in \mathcal{S}$ for every infinite $N \subseteq M$ with $X < N$.
3. We say that M rejects X if no infinite $N \subseteq M$ accepts X .
4. We say that M determines X if it either accepts or rejects X .

Remark 3.5

1. If M accepts X , and $N \subseteq M$ is infinite, then N accepts X .
2. The same holds for "rejects".
3. By definition of "rejects", for any finite X and infinite M , there is infinite $N \subseteq M$ that determines X .

Definition 3.6 A Prikry pair is a pair (X, N) where $X, N \subseteq \omega$, X is finite, N is infinite and $X < N$. We partially order Prikry pairs by: $(X, N) < (X', N')$ iff X' is an end extension of X , $N' \subseteq N$, $X' \setminus X \subseteq N$.

Lemma 3.7 There is an infinite $M \subseteq \omega$ that determines every finite subset of itself.

Proof. Inductively choose Prikry pairs (X_i, M_i) ($i < \omega$) so that M_i determines every subset of X_i . We let $X_0 = \emptyset$ and take M_0 to be any infinite set determining \emptyset . Inductively, let $X_{i+1} = X_i \cup \min(M_i)$ and choose $M_{i+1} \subseteq M_i \setminus \min(M_i)$ that determines every subset of X_{i+1} (use (3) of the remark repeatedly). Let $M = \bigcup_{i < \omega} X_i$. Observe that $M \subseteq X_i \cup M_i$ for all i . We claim that M works. If $X \subseteq M$ is finite, then for all $i < \omega$, $X \subseteq X_i$ iff $M_i > X$. Let i be least such that this condition holds. Then if $N \subseteq M$ and $N > X$, we have $N \subseteq M_i$. But M_i determines X , and hence so does N . \square

Lemma 3.8 There is infinite $N \subseteq \omega$ that either accepts each finite subset of itself, or rejects each finite subset of itself.

Proof. Take M as in lemma 3.7. Select elements $a_0 < a_1 < \dots$ of M by induction, so that for each i , if $X \subseteq \{a_0, \dots, a_{i-1}\}$ and M rejects X then M rejects $X \cup \{a_i\}$. For this, suppose that a_0, \dots, a_{i-1} are all chosen. Let X_1, \dots, X_k list all subsets of $\{a_0, \dots, a_{i-1}\}$ that are rejected by M . We want a_i so that M rejects all of $X_1 \cup \{a_i\}, \dots, X_k \cup \{a_i\}$.

Suppose there's no such a_i . Then for all $a_{i-1} < a \in M$ there's $j_a \leq k$ such that M doesn't reject $X_{j_a} \cup \{a\}$. By lemma 3.7, M accepts $X_{j_a} \cup \{a\}$. Clearly there is some $j \leq k$ such that $j = j_a$ for infinitely many a . Put $N = \{a : j_a = j\}$. Then $N \subseteq M$ is infinite, and accepts X_j . For let $P \subseteq N$ be infinite with $X_j < P$. If $p = \min(P)$, then by choice of N , M accepts $X_j \cup \{p\}$. So as $P \setminus \{p\} \subseteq M$, we have $X_j \cup \{p\} \cup (P \setminus \{p\}) \in \mathcal{S}$. That is, $X_j \cup P \in \mathcal{S}$ for all such P ; this contradicts the fact that M rejects X_j .

So we can choose a_0, a_1, \dots as stated. Put $N = \{a_i : i < \omega\}$. There are two cases. If N accepts \emptyset then by definition, $[N]^\omega \subseteq \mathcal{S}$. So N accepts any finite subset of itself. If not, then it is easily seen by induction on $|X|$ that N rejects any $X \subseteq N$. \square

Proof of theorem 3.3: By lemma 3.8, we can take an infinite N that uniformly decides all its finite subsets. If N accepts them, then as above, it accepts \emptyset so $[N]^\omega \subseteq \mathcal{S}$. If it rejects them, we claim $[N]^\omega \cap \mathcal{S} = \emptyset$. For if not, there's infinite $X \subseteq N$ such that $X \in \mathcal{S}$. Now as \mathcal{S} is open, we can take $n < \omega$ such that if $Y \cap n = X \cap n$ and Y is infinite then $Y \in \mathcal{S}$. We can increase n as we like; so as X is infinite, we can assume that $n-1 \in X$. But now, if $P \subseteq N$

is infinite and $P > X \cap n$ then $\min(P) \geq n$. So $((X \cap n) \cup P) \cap n = X \cap n$, and $(X \cap n) \cup P \in \mathcal{S}$. Hence, N accepts $X \cap n$, a contradiction. \square

It's a good exercise to prove Ramsey's theorem now. We will wait a little (see corollary 4.4).

4 Barriers

Definition 4.1 *Let X be an infinite subset of ω .*

1. *We say that $\mathcal{B} \subseteq [X]^{<\omega}$ is a barrier on X if:
 - *for every infinite $Y \subseteq X$, there is an initial segment of Y in \mathcal{B} ,*
 - *\mathcal{B} is an antichain with respect to \subseteq .**
2. *A barrier is a barrier on some infinite $X \subseteq \omega$.*
3. *Clearly, for a barrier \mathcal{B} on X and infinite $Y \subseteq X$, there is a unique initial segment of Y in \mathcal{B} . We write this initial segment as $Y \upharpoonright \mathcal{B}$.*
4. *The base of a barrier \mathcal{B} is defined to be $\bigcup \mathcal{B}$.*

Remark 4.2 It can be checked that if \mathcal{B} is a barrier on X then $X = \bigcup \mathcal{B}$. Bearing in mind that X is order-isomorphic to ω , we see that any barrier is 'isomorphic' to a barrier on ω . We seem more interested in barrier-ness than in what the base of a barrier is.

For any $n < \omega$, $[\omega]^n$ is a barrier on ω . Thinking of $n = 1$, where $\{\{0\}, \{1\}, \dots\}$ is a barrier, we will perhaps see (especially in definition 5.1) that a barrier is a kind of generalised sequence, or rather, the index set of such a sequence.

4.1 The Nash-Williams Ramsey theorem

Theorem 4.3 (Nash-Williams Ramsey theorem) *Let \mathcal{B} be a barrier on $X \in [\omega]^\omega$. Suppose that $\mathcal{B} = \mathcal{B}_1 \cup \mathcal{B}_2$. Then there is an infinite subset Y of X such that one of $\mathcal{B}_i \cap [Y]^{<\omega}$ ($i = 1, 2$) is a barrier on Y .*

Proof. [Rewritten Jan 2003]. X is in order-preserving bijection with ω , so we may assume without loss of generality that $X = \omega$. Let

$$\mathcal{S} = \{Z \in [\omega]^\omega : Z \text{ has an initial segment in } \mathcal{B}_1\}.$$

Clearly, \mathcal{S} is open, and hence (by theorem 3.3) Ramsey. Take $Y \in [\omega]^\omega$ such that $[Y]^\omega \subseteq \mathcal{S}$ or $[Y]^\omega \cap \mathcal{S} = \emptyset$. If $[Y]^\omega \subseteq \mathcal{S}$, then clearly $\mathcal{B}_1 \cap [Y]^{<\omega}$ is a

barrier on Y . If on the other hand $[Y]^\omega \cap \mathcal{S} = \emptyset$, then take any $Z \in [Y]^\omega$. Since \mathcal{B} is a barrier on ω , Z has an initial segment in \mathcal{B} . But $Z \notin \mathcal{S}$. So that initial segment must be in \mathcal{B}_2 . Hence, $\mathcal{B}_2 \cap [Y]^{<\omega}$ is a barrier on Y . \square

Corollary 4.4 (Ramsey's theorem) *If $f : [\omega]^n \rightarrow k$ where $k, n < \omega$, then there is infinite $N \subseteq \omega$ such that $f \upharpoonright [N]^n$ is constant.*

Proof. For some fixed n , let $\mathcal{B} = [\omega]^n$. Obviously, \mathcal{B} is a barrier on ω . If $\mathcal{B} = \mathcal{B}_1 \cup \mathcal{B}_2 \cup \dots \cup \mathcal{B}_k$, iterated application of the theorem gives an infinite $Y \subseteq \omega$ such that some $\mathcal{B}_i \cap [Y]^{<\omega}$ ($i \leq k$) is a barrier on Y . If e.g., $\mathcal{B}_1 \cap [Y]^{<\omega}$ is a barrier on Y , then every n -element subset s of Y is in \mathcal{B}_1 . For if we end-extend s to an infinite subset S of Y , then S has an initial segment in \mathcal{B}_1 . As this must have size n , it must be s . The theorem now follows by letting $\mathcal{B}_i = f^{-1}(i-1)$ for each $1 \leq i \leq k$, since then $[Y]^n \subseteq f^{-1}(j)$ for some $j < k$. \square

4.2 The barrier $\mathcal{B}(2)$

Definition 4.5

1. If s, t are finite subsets of ω , we write $s \triangleleft t$ to mean that there are $i_1 < \dots < i_k$ and j ($1 \leq j < k$) such that $s = \{i_1, \dots, i_j\}$ and $t = \{i_2, \dots, i_k\}$.
2. Given a barrier \mathcal{B} and a pre-ordered set Q , we say that a map $f : \mathcal{B} \rightarrow Q$ is bad if there are no $s, t \in \mathcal{B}$ with $s \triangleleft t$ and such that $f(s) \leq f(t)$.
3. A map $f : \mathcal{B} \rightarrow Q$ is perfect if $f(s) \leq f(t)$ for all $s \triangleleft t$ in \mathcal{B} .

Note that $j = 1$ is allowed; $k > j$ by barrier-ness. \triangleleft is not transitive in general.

Example 4.6 Suppose \mathcal{B} is a barrier on ω . For any infinite $X \subseteq \omega$, temporarily write $X_0 = X$ and $X_{n+1} = X_n \setminus \{\min(X_n)\}$. Each X_n has an initial segment $s_n \in \mathcal{B}$. Then $s_0 \triangleleft s_1 \triangleleft \dots$. Moreover, if $s_0 \triangleleft s_1 \triangleleft \dots \triangleleft s_n$ in \mathcal{B} , choose infinite $X \subseteq \omega$ having $s_0 \cup \dots \cup s_n$ as an initial segment. Then $s_i = X_i \upharpoonright \mathcal{B}$ for all $i \leq n$. This shows that s_0, \dots, s_n are recoverable from $s_0 \cup \dots \cup s_n$.

If a barrier \mathcal{B} 'is' the index set of a generalised sequence, and $s, t \in \mathcal{B}$, then $s \triangleleft t$ means that t is a 'higher' index than s .

Definition 4.7 *Let \mathcal{B} be a barrier on X . Write $\mathcal{B}(2) = \{s \cup t : s \triangleleft t \text{ in } \mathcal{B}\}$.*

Theorem 4.8 $\mathcal{B}(2)$ is a barrier on X .

Proof. We show that if $s, t, s', t' \in \mathcal{B}$, $s \triangleleft t$, $s' \triangleleft t'$, and $s \cup t \subseteq s' \cup t'$, then $s = s'$ and $t = t'$. For let $s \cup t$ be $\{i_1, \dots, i_k\}$ and $s' \cup t'$ be $\{j_1, \dots, j_m\}$, in increasing order. Then $t = \{i_2, \dots, i_k\}$ and $t' = \{j_2, \dots, j_m\}$. Clearly $t \subseteq t'$; so as \mathcal{B} is a barrier, we have $t = t'$. Hence $i_2 = j_2$ (as $i_1 = j_1$). Hence s, s' are initial segments of $s \cup t$, and so they are equal as \mathcal{B} is a barrier. So there are no proper inclusions between elements of $\mathcal{B}(2)$.

Let Y be any infinite subset $\{y_0, y_1, \dots\}$ of X , listed in increasing order. If \mathcal{B} is a barrier then Y has an initial segment $s \in \mathcal{B}$. Similarly, $Y \setminus \{y_0\}$ has an initial segment $t \in \mathcal{B}$. Then $s \triangleleft t$ by barrier-ness. So $s \cup t$ is an initial segment of Y in $\mathcal{B}(2)$. \square

Remark 4.9 The proof (and example 4.6) shows that each set $s \cup t \in \mathcal{B}(2)$ allows s, t to be recovered uniquely. So we will always write elements of $\mathcal{B}(2)$ in the form $s \cup t$, where (implicitly) $s, t \in \mathcal{B}$.

If $s \cup t, s' \cup t' \in \mathcal{B}(2)$ and $s \cup t \triangleleft s' \cup t'$, then clearly, t is an initial segment of $s' \cup t'$. Since s' is as well, we have $s' \subseteq t$ or $t \subseteq s'$. Since $s', t \in \mathcal{B}$, we have $t = s'$. So $s \triangleleft s'$ and $t \triangleleft t'$.

Exercise 4.10 Define $\mathcal{B}(n)$ for $2 \leq n < \omega$, and generalise the above to it.

Corollary 4.11 Given a barrier \mathcal{B} on a set $X \in [\omega]^\omega$, and a map $f : \mathcal{B} \rightarrow Q$, there exists a barrier \mathcal{D} on some $Y \in [X]^\omega$ with $\mathcal{D} \subseteq \mathcal{B}$, and such that the restriction $f \upharpoonright \mathcal{D}$ of f to \mathcal{D} is either bad or perfect.

Proof. Put

$$\begin{aligned} \mathcal{A}_1 &= \{s \cup t \in \mathcal{B}(2) : f(s) \leq f(t)\}, \\ \mathcal{A}_2 &= \{s \cup t \in \mathcal{B}(2) : f(s) \not\leq f(t)\}. \end{aligned}$$

Since $\mathcal{A}_1 \cup \mathcal{A}_2 = \mathcal{B}(2)$, by the Nash-Williams Ramsey theorem, there is infinite $Y \subseteq X$ such that some $\mathcal{C} \subseteq \mathcal{A}_1$ or $\mathcal{C} \subseteq \mathcal{A}_2$ is a barrier on Y .

Suppose $\mathcal{C} \subseteq \mathcal{A}_1$. Put

$$\mathcal{D} = \{s \in \mathcal{B} : s \cup t \in \mathcal{C} \text{ for some } t \text{ with } s \triangleleft t\}.$$

Then \mathcal{D} is a barrier on Y — certainly \mathcal{D} is an antichain since $\mathcal{D} \subseteq \mathcal{B}$, and further, if $Z \subseteq Y$ is infinite, there is an initial segment $s \cup t$ of Z with $s \cup t \in \mathcal{C}$, so $s \in \mathcal{D}$ is an initial segment of Z .

Let $g = f \upharpoonright \mathcal{D}$; then g is perfect. For suppose $s \triangleleft t$ in \mathcal{D} . End-extend $s \cup t$ to an infinite set in Y . This has an initial segment $s' \cup t' \in \mathcal{C}$. As in the proof of theorem 4.8, $s = s'$ and $t = t'$. So $s \cup t \in \mathcal{C} \subseteq \mathcal{A}_1$.

The case $\mathcal{C} \subseteq \mathcal{A}_2$ is similar; we get a bad g . \square

5 Better-quasi-orderings

This is a long and interesting section.

Definition 5.1 (bqo) Let (Q, \leq) be a pre-order. We say that \leq is a better-quasi-ordering (bqo) if for every barrier \mathcal{B} on ω , there is no bad map $f : \mathcal{B} \rightarrow Q$.

Because any barrier is isomorphic to a barrier on ω , if \leq is a bqo then for every barrier \mathcal{B} , there is no bad map $\mathcal{B} \rightarrow Q$.

Example 5.2 (ω, \leq) is a bqo. For let \mathcal{B} be a barrier on ω , and $f : \mathcal{B} \rightarrow \omega$. Let s_n be an initial segment of $\omega \setminus n$ in \mathcal{B} . Then $s_0 \triangleleft s_1 \triangleleft \dots$. If f is bad, then $f(s_0) \not\leq f(s_1) \not\leq \dots$, which is impossible.

Considering the barrier $[\omega]^1 = \{\{n\} : n < \omega\}$, we see that a bqo is a wqo.

5.1 Combining bqos

Lemma 5.3 Let \leq, \sqsubseteq be bqos on the same set Q . Then $\preceq =_{def.} \leq \cap \sqsubseteq$ is also a bqo on Q .

Proof. Certainly, \preceq is a pre-order. Let \mathcal{B} be a barrier on ω , and $f : \mathcal{B} \rightarrow Q$. By corollary 4.11, there exists a barrier $\mathcal{D} \subseteq \mathcal{B}$ such that $f \upharpoonright \mathcal{D}$ is either bad or perfect with respect to (Q, \leq) . Since \leq is a bqo, $f \upharpoonright \mathcal{D}$ cannot be bad, so is perfect. Since \sqsubseteq is a bqo, $f \upharpoonright \mathcal{D}$ cannot be bad with respect to (Q, \sqsubseteq) , so there are $s \triangleleft t$ in \mathcal{D} with $f(s) \sqsubseteq f(t)$. By perfection, $f(s) \leq f(t)$. So $f(s) \preceq f(t)$, showing that f is not bad. \square

5.2 Power sets of bqos

Definition 5.4 Let (Q, \leq) be a pre-order. Define a pre-order \leq on $\wp(Q)$ by $\Gamma \leq \Delta$ iff for all $\delta \in \Delta$ there is $\gamma \in \Gamma$ with $\gamma \leq \delta$.

The following fails for wqos; cf. exercise 1.14.

Proposition 5.5 If (Q, \leq) is a bqo then so is $(\wp(Q), \leq)$.

Proof. Let \mathcal{B} be a barrier on ω , and $f : \mathcal{B} \rightarrow \wp(Q)$. Assume for contradiction that f is bad. So for each $s \cup t \in \mathcal{B}(2)$, $f(s) \not\leq f(t)$, so there is an element $h(s \cup t) \in f(t)$ such that for every $q \in f(s)$ we have $h(s \cup t) \not\leq q$. (Recall from remark 4.9 that s, t are uniquely recoverable from $s \cup t \in \mathcal{B}(2)$.) We have therefore defined a map $h : \mathcal{B}(2) \rightarrow Q$.

By theorem 4.8, $\mathcal{B}(2)$ is a barrier on ω . As Q is a bqo, h is not bad, and there are $s \cup t \triangleleft s' \cup t'$ in $\mathcal{B}(2)$ with $h(s \cup t) \leq h(s' \cup t')$. By remark 4.9, $t = s'$. But now, $h(s' \cup t') \in f(t')$ is $\geq h(s \cup t) \in f(t) = f(s')$. This contradicts the definition of h . \square

5.3 Minimal bad maps

The presentation of the results of this and the following two sections is similar to that of [1].

Let Q be pre-ordered by \leq . We assume that Q also carries a transitive well-founded relation \prec such that $x \prec y \Rightarrow x < y$ for all $x, y \in Q$. We write $x \preceq y$ to mean $x \prec y$ or $x = y$.

Definition 5.6

1. For sets $s, t \subseteq \omega$, we write $t \preceq s$ to mean that t is an initial segment of s , and $t \prec s$ to mean that t is a proper initial segment of s .
2. For barriers \mathcal{B}, \mathcal{C} on X, Y , we write $\mathcal{B} \sqsubseteq \mathcal{C}$ (\mathcal{B} foreruns \mathcal{C}) if $Y \subseteq X$ and every element of \mathcal{C} has an initial segment in \mathcal{B} .
3. We say $\mathcal{B} \sqsubset \mathcal{C}$ (strictly foreruns) if $\mathcal{B} \sqsubseteq \mathcal{C}$ and some element of \mathcal{C} has a proper initial segment in \mathcal{B} . Example: $\mathcal{B} \sqsubset \mathcal{B}(2)$.
4. Given $f : \mathcal{B} \rightarrow Q$ and $g : \mathcal{C} \rightarrow Q$, we write $f \sqsubseteq g$ if
 - (a) $\mathcal{B} \sqsubseteq \mathcal{C}$,
 - (b) $f(s) = g(s)$ for all $s \in \mathcal{B} \cap \mathcal{C}$,
 - (c) for all $b \in \mathcal{B}, c \in \mathcal{C}$ with $b \prec c$, we have $g(c) \prec f(b)$.

We write $f \sqsubset g$ if $f \sqsubseteq g$ and $\mathcal{B} \sqsubset \mathcal{C}$.

5. We say that a bad map $f : \mathcal{B} \rightarrow Q$ is minimal if there is no bad $g \sqsubset f$. (Minimality is with respect to \preceq , not \sqsubseteq !)
6. If $f : \mathcal{B} \rightarrow Q$ is bad but not minimal bad, then there are $\mathcal{C} \sqsubset \mathcal{B}$ and a bad $g : \mathcal{C} \rightarrow Q$ with $f \sqsubset g$. So there are $b \in \mathcal{B}, c \in \mathcal{C}$ with $b \prec c$. Write k for the greatest element of b . Fixing f , choose \mathcal{C}, g, b, c to make k minimal. We will write this least value of k as $k(f)$.

Remark 5.7 It can be checked that both \sqsubseteq s are reflexive and transitive. (Take care with condition 4b.)

If $\mathcal{C} \sqsubseteq \mathcal{B}$ are barriers and $\mathcal{C} \sqsubseteq \mathcal{D}$, then $\mathcal{B} \sqsubseteq \mathcal{D}$. If $f : \mathcal{B} \rightarrow Q, g : \mathcal{D} \rightarrow Q$, and $f|_{\mathcal{C}} \sqsubseteq g$, then $f \sqsubseteq g$. So if $f : \mathcal{B} \rightarrow Q$ is minimal bad, then so is $f|_{\mathcal{C}}$.

5.4 Nash-Williams/Laver theorem

Lemma 5.8 *Given $f : \mathcal{B} \rightarrow Q$ which is bad but not minimal bad, there are $\mathcal{D} \sqsupset \mathcal{B}$ and a bad $h : \mathcal{D} \rightarrow Q$ with $h \sqsupset f$, such that*

1. *for some $b \in \mathcal{B}$ with maximal element $k(f)$, we have $b \notin \mathcal{D}$,*
2. *every $m \leq k(f)$ in the base $\bigcup \mathcal{B}$ of \mathcal{B} is in the base of \mathcal{D} .*

Proof. Take bad $f : \mathcal{B} \rightarrow Q$, not minimal bad, and $g \sqsupset f$ bad, such that $g : \mathcal{C} \rightarrow Q$ for some $\mathcal{C} \sqsupset \mathcal{B}$, and for some $s \in \mathcal{B}$ and $t \in \mathcal{C}$ we have $s \triangleleft t$ and $\max(s) = k(f) = n$, say. Put

$$\begin{aligned} S &= (\{0, 1, \dots, n\} \cap \bigcup \mathcal{B}) \setminus \bigcup \mathcal{C}, \\ \mathcal{D} &= \mathcal{C} \cup \{b \in \mathcal{B} : b \subseteq \bigcup \mathcal{C} \cup S \text{ and } b \cap S \neq \emptyset\}. \end{aligned}$$

Claim: \mathcal{D} is a barrier on $\bigcup \mathcal{C} \cup S$.

Proof of Claim: Certainly, $\mathcal{D} \subseteq [\bigcup \mathcal{C} \cup S]^{<\omega}$. We check that \mathcal{D} is an \subseteq -antichain. Since this is true for \mathcal{B} and for \mathcal{C} , it suffices to take $b \in \mathcal{B}$ with $b \subseteq \bigcup \mathcal{C} \cup S$ and $b \cap S \neq \emptyset$, and $c \in \mathcal{C}$, and check that $b \not\subseteq c \not\subseteq b$. Since $c \cap S = \emptyset$, we have $b \not\subseteq c$. If $c \subseteq b$, then $c \subset b$. But c has an initial segment $b' \in \mathcal{B}$, so $b' \subset b$, contradicting that \mathcal{B} is a barrier.

Now let $X \subseteq \bigcup \mathcal{C} \cup S$ be infinite; we want to find an initial segment of X in \mathcal{D} . If X has an initial segment in \mathcal{C} , we are done. Assume it doesn't. Let $b = X \upharpoonright \mathcal{B}$ (see definition 4.1). Certainly, $b \subseteq \bigcup \mathcal{C} \cup S$. We claim that $b \cap S \neq \emptyset$, so that $b \in \mathcal{D}$.

Assume for contradiction that $b \cap S = \emptyset$. As X has no initial segment in \mathcal{C} , a barrier on $\bigcup \mathcal{C}$, we have $X \not\subseteq \bigcup \mathcal{C}$, and so $X \cap S \neq \emptyset$. Take $s \in X \cap S$. Now $b \cap S = \emptyset$, and $b \triangleleft X$. So $\max(b) < s \leq n$.

Let $Y = X \cap \bigcup \mathcal{C}$. Then Y is infinite, so has an initial segment c in \mathcal{C} . Since $b \triangleleft X$ and $b \subseteq \bigcup \mathcal{C}$, we have $b \triangleleft Y$ as well. So b, c are \leq -comparable. If $c \triangleleft b$, then c has an initial segment b' in \mathcal{B} , and $b' \subset b$, which is impossible as \mathcal{B} is a barrier. So $b \leq c$. But $\max(b) < n$, so by minimality of n we must have $b = c$. So X has an initial segment c in \mathcal{C} , a contradiction. This proves the claim.

We now show that \mathcal{D} has the two properties cited in the lemma. If $b \in \mathcal{B}$, $c \in \mathcal{C}$, and $b \triangleleft c$, then $b \subset c \in \mathcal{D}$; so by the claim, $b \notin \mathcal{D}$. Since \mathcal{B} contains such a b with $\max(b) = n$, the first property is established. Also, every $x \in \{0, 1, \dots, n\} \cap \bigcup \mathcal{B}$ is the least element of some infinite subset of $\bigcup \mathcal{C} \cup S$, which by the claim has an initial segment $d \in \mathcal{D}$. So $x \in d$, whence $\{0, 1, \dots, n\} \cap \bigcup \mathcal{B} \subseteq \bigcup \mathcal{D}$.

It is clear that $\mathcal{B} \sqsubset \mathcal{D}$. Define $h : \mathcal{D} \rightarrow Q$ by

$$h(d) = \begin{cases} f(d), & \text{if } d \in \mathcal{B}, \\ g(d), & \text{if } d \in \mathcal{C}. \end{cases}$$

This is well-defined, because $f \sqsubseteq g$. Trivially, if $d \in \mathcal{B} \cap \mathcal{D}$, $h(d) = f(d)$. Also, if $b \in \mathcal{B}$, $d \in \mathcal{D}$, and $b \triangleleft d$, then $d \notin \mathcal{B}$. So $h(d) = g(d) \prec f(b)$ since $f \sqsubseteq g$. So $f \sqsubset h$.

Moreover, h is bad. For let $s \triangleleft t$ in \mathcal{D} . There are four cases.

1. If $s, t \in \mathcal{B}$, then as f is bad, $h(s) = f(s) \not\leq f(t) = h(t)$.
2. The case where $s, t \in \mathcal{C}$ is similar, using badness of g .
3. Assume that $s \in \mathcal{B}$ and $t \in \mathcal{C}$. Let $b \in \mathcal{B}$ be an initial segment of t . Since $f \sqsubseteq g$, we have $g(t) \preceq f(b)$, and so $g(t) \leq f(b)$. Since $s \triangleleft t$, $\min(s) \notin b$, so $s \neq b$, and hence as \mathcal{B} is a barrier, $b \not\sqsubseteq s$. It follows that $s \triangleleft b$. If $h(s) \leq h(t)$ — i.e., $f(s) \leq g(t)$ — then by transitivity, $f(s) \leq f(b)$, contradicting badness of f . So $h(s) \not\leq h(t)$ as required.
4. Finally assume for contradiction that $s \in \mathcal{C} \setminus \mathcal{B}$ and $t \in \mathcal{B} \setminus \mathcal{C}$. Then $t \cap S \neq \emptyset$. As $s \cap S = \emptyset$, and $s \triangleleft t$, there is $x \in S \cap t$ larger than all elements of s . So $\max(s) < n$. But s has an initial segment b in \mathcal{B} , which must also have maximum $< n$. By minimality of n we have $s = b \in \mathcal{B}$, a contradiction.

□

The main theorem is next. The idea is that \mathcal{B} is to be refined in a step-by-step manner to $\mathcal{C} \sqsupset \mathcal{B}$ by reducing values of f . We may need to extend sequences to do this.

Theorem 5.9 (Nash-Williams, interpreted by Laver) *For every barrier \mathcal{B} and bad map $f : \mathcal{B} \rightarrow Q$, there are a barrier $\mathcal{C} \sqsupset \mathcal{B}$ and a minimal bad $g : \mathcal{C} \rightarrow Q$ with $f \sqsubseteq g$.*

Proof. If \mathcal{B}_0 is a barrier and $f_0 : \mathcal{B}_0 \rightarrow Q$ is bad but not minimal bad, we can find a barrier \mathcal{B}_1 with $\mathcal{B}_0 \sqsubseteq \mathcal{B}_1$, and a bad map $f_1 : \mathcal{B}_1 \rightarrow Q$ with $f_1 \sqsupset f_0$ as in lemma 5.8. If f_1 is not minimal bad, we can repeat to get \mathcal{B}_2 and $f_2 : \mathcal{B}_2 \rightarrow Q$, etc. If this process stops after a finite number of iterations, we're done.

Otherwise, it goes on infinitely many times. In this case, we claim that $k(f_0) \leq k(f_1) \leq \dots$ and the sequence rises infinitely often. Certainly the sequence cannot fall, as $k(f_i)$ is minimal. Each application of the lemma removes from \mathcal{B}_i some set b with maximum $k(f_i)$, and \mathcal{B}_{i+1} contains some c with $b \triangleleft c$. No later \mathcal{B}_j contains b , because $\mathcal{B}_{i+1} \sqsubseteq \mathcal{B}_j$, so there would be $b' \leq b$ in \mathcal{B}_{i+1} , so $b' \subset c$, contradicting that \mathcal{B}_{i+1} is a barrier. There are finitely many such sets b , so this cannot happen infinitely often. The claim is proved.

Let X_i be the base of \mathcal{B}_i ($i < \omega$). Since $\mathcal{B}_0 \sqsubseteq \mathcal{B}_1 \sqsubseteq \dots$, we have $X_0 \supseteq X_1 \supseteq \dots$. Define $X = \bigcap_{n < \omega} X_n$. Since $k(f_i) \in X$ for all i (see lemma 5.8), we see by the claim that X is infinite. Let $\mathcal{D} = \{b : b \text{ is in cofinitely many } \mathcal{B}_i\}$.

Claim: \mathcal{D} is a barrier on X .

Proof of Claim: Certainly, if $b \in \mathcal{D}$ then $b \subseteq X_i$ for cofinitely many $i < \omega$, so $b \subseteq X$. Also, if $b, c \in \mathcal{D}$ and $b \subseteq c$, then $b = c$ (as this holds in cofinitely many \mathcal{B}_i).

Consider any infinite $Y \subseteq X$. For each \mathcal{B}_i , we have $Y \subseteq X_i$, so some $b_i \in \mathcal{B}_i$ is an initial segment of Y . Clearly, $b_i \leq b_{i+1}$ or $b_{i+1} \triangleleft b_i$. If $b_{i+1} \triangleleft b_i$, then as $\mathcal{B}_i \sqsubseteq \mathcal{B}_{i+1}$, b_{i+1} has an initial segment $c \in \mathcal{B}_i$. So $c \subset b_i$, which is impossible since \mathcal{B}_i is a barrier. So $b_i \leq b_{i+1}$. If $b_i \triangleleft b_{i+1}$, then $f_i(b_i) \succ f_{i+1}(b_{i+1})$. Otherwise, $b_i = b_{i+1}$, so since $f_{i+1} \sqsupseteq f_i$, we have $f_i(b_i) = f_{i+1}(b_{i+1})$. But \triangleleft is well-founded, so the sequence $(b_i : i < \omega)$ becomes constant at j , say. Then $b_j \in \mathcal{D}$ is an initial segment of Y . The claim is proved.

We note that $\mathcal{D} \sqsupseteq \mathcal{B}_i$ for each i .

We now define $g : \mathcal{D} \rightarrow Q$ to agree eventually with all f_i — i.e., for all $d \in \mathcal{D}$, $g(d) = f_i(d)$ for all i with $d \in \mathcal{B}_i$. This is well-defined since $f_0 \sqsubseteq f_1 \sqsubseteq \dots$. We note that $g \sqsupseteq f_i$ for each i .

Claim: g is minimal bad.

Proof of Claim: If $d \triangleleft e$ in \mathcal{D} then $d \triangleleft e$ in some \mathcal{B}_i . So $g(d) = f_i(d) \not\leq f_i(e) = g(e)$. So g is bad.

If there are $\mathcal{E} \sqsupseteq \mathcal{D}$ and bad $h : \mathcal{E} \rightarrow Q$ with $h \sqsupseteq g$, then there are $s \in \mathcal{D}$ and $t \in \mathcal{E}$ with $s \triangleleft t$. We may choose $i < \omega$ with $s \in \mathcal{B}_i$ and $\max(s) < k(f_i)$ (as the $k(f_i)$ rise arbitrarily high). Then (cf. remark 5.7) $\mathcal{E} \sqsupseteq \mathcal{B}_i$ and $h \sqsupseteq f_i$, so h contradicts the value of $k(f_i)$. This proves theorem 5.9. \square

5.5 Sequences

Now for some practical applications. Fix a set Q and a pre-order \leq on it.

Definition 5.10

1. We can extend \leq to $\bigcup_{\alpha \text{ ordinal}} Q^\alpha$ as follows. For $x \in Q^\alpha$, $y \in Q^\beta$, put $x \leq y$ iff there is a one-one order-preserving map $\varphi : \alpha \rightarrow \beta$ such that for all $i < \alpha$ we have $x_i \leq y_{\varphi(i)}$. (Regarding $Q = Q^1$, this is an extension of \leq on Q .) As usual, $x < y$ means $x \leq y \not\leq x$.
2. For $x, y \in \bigcup_{\alpha} Q^\alpha$, define $x \prec y$ to mean x is isomorphic to a subsequence of y and $\text{dom}(x) < \text{dom}(y)$ as ordinals. Clearly, \prec is transitive and well-founded (as ordinals are), and $x \prec y \Rightarrow x < y$ ($<$ as above).

Theorem 5.11 *If (Q, \leq) is a bqo then so is $(\bigcup_{\alpha} Q^\alpha, \leq)$.*

Proof. Suppose for contradiction that there is a barrier \mathcal{B} and a bad map $f : \mathcal{B} \rightarrow \bigcup_{\alpha} Q^\alpha$. By theorem 5.9, we can assume that f is minimal bad. By the Nash-Williams Ramsey theorem (theorem 4.3), we can also assume that either all $f(b)$ have successor length, or they all have limit length. By remark 5.7, this does not affect minimality.

Case I: the $f(b)$ have successor length. Write $f(b) = f_1(b) \wedge f_2(b)$ with $f_2(b) \in Q$ (we ‘remove’ the last element). By corollary 4.11, we can assume that each of f_1, f_2 is either bad or perfect. Now f_2 can’t be bad, as Q is a bqo; hence it is perfect. But f is bad, so f_1 must be bad.

Define $g : \mathcal{B}(2) \rightarrow \bigcup_{\alpha} Q^\alpha$ by $g(s \cup t) = f_1(s)$. Then $\mathcal{B}(2) \sqsupset \mathcal{B}$; and because $f_1(s) \prec f(s)$ for all s , we have $g \sqsupset f$. Also, g is bad: for if $s \cup t \triangleleft t \cup v$ in $\mathcal{B}(2)$ (see remark 4.9), then because $s \triangleleft t$ and f_1 is bad, $g(s \cup t) = f_1(s) \not\leq f_1(t) = g(t \cup v)$. This contradicts the minimality of f .

Case II: each $f(b)$ has limit length. Consider $s \triangleleft t$ in \mathcal{B} . Because of the poor quality of f we have $f(s) \not\leq f(t)$. Define an order-preserving map $\varphi : \text{dom}(f(s)) \rightarrow \text{dom}(f(t))$ by induction: at each stage, φ maps the next element $i < \text{dom}(f(s))$ to the least $j < \text{dom}(f(t))$ such that $f(t)_j \geq f(s)_i$ and $j > \varphi(k)$ for all $k < i$. If this succeeded, we’d have $f(s) \leq f(t)$; so there must be a proper initial segment of $f(s)$ that’s $\not\leq f(t)$. Take (say) the shortest such — $f(s)_t$, say.

Define $g : \mathcal{B}(2) \rightarrow \bigcup_{\alpha} Q^\alpha$ by $g(s \cup t) = f(s)_t$. This is a non-trivial use of $\mathcal{B}(2)$. As $\mathcal{B} \sqsubset \mathcal{B}(2)$, and $f(s)_t \prec f(s)$ whenever $s \triangleleft t$, we have $f \sqsubset g$. But g is bad, for if $s \cup t \triangleleft t \cup v$ then $g(s \cup t) = f(s)_t \not\leq f(t)$, so $g(s \cup t) \not\leq f(t)_v = g(t \cup v)$. This is a contradiction. \square

5.6 Matrices

Fix $1 \leq d < \omega$, and write Q for ω^d . Recall from example 5.2 that (ω, \leq) is a bqo. Extend \leq to $\bigcup_{\alpha} \omega^\alpha$, as in definition 5.10. By theorem 5.11, $(\bigcup_{\alpha} \omega^\alpha, \leq)$

is a bqo, and hence so is its subset (Q, \leq) . Note that $x \leq y$ in Q iff $x_i \leq y_i$ for all $i < d$.

By definition 5.10 and theorem 5.11 again, we may extend the order \leq on Q to an order \leq on $\bigcup_\alpha Q^\alpha$, and it is a bqo.

Definition 5.12 Define \sqsubseteq on $\bigcup_\alpha Q^\alpha$ by: if $\xi \in Q^\alpha$, $\eta \in Q^\beta$, then $\xi \sqsubseteq \eta$ iff $\forall j < \beta \exists i < \alpha (\xi_i \leq \eta_j)$

Lemma 5.13 $(\bigcup_\alpha Q^\alpha, \sqsubseteq)$ is a bqo.

Proof. For a sequence $\xi = (\xi_i : i < \alpha) \in Q^\alpha$, write $\text{rng}(\xi)$ for the set $\{\xi_i : i < \alpha\} \in \wp(Q)$. Given a barrier \mathcal{B} on ω and a map $f : \mathcal{B} \rightarrow \bigcup_\alpha Q^\alpha$, define $g : \mathcal{B} \rightarrow \wp(Q)$ by $g(s) = \text{rng}(f(s))$. By proposition 5.5, there are $s \triangleleft t$ in \mathcal{B} such that for all $q \in g(t)$ there is $p \in g(s)$ with $p \leq q$. So clearly, $f(s) \sqsubseteq f(t)$. \square

Definition 5.14 Define \preceq on $\bigcup_\alpha Q^\alpha$ by: $\xi \preceq \eta$ iff $\xi \leq \eta$ and $\xi \sqsubseteq \eta$.

By lemma 5.3, \preceq is a bqo on $\bigcup_\alpha Q^\alpha$.

Definition 5.15

1. For an ordinal α , let \mathcal{M}_α be the set of all maps $m : \alpha \times d \rightarrow \omega$.
2. Define a bijection $\theta_\alpha : Q^\alpha \rightarrow \mathcal{M}_\alpha$ by $\theta_\alpha(\xi) = m$, where $m(i, j) = (\xi_i)_j$ for $i < \alpha$, $j < d$.
3. Let $\mathcal{M} = \bigcup_\alpha \mathcal{M}_\alpha$. Define a bijection $\theta = (\bigcup_\alpha \theta_\alpha) : \bigcup_\alpha Q^\alpha \rightarrow \mathcal{M}$.
4. For $m, n \in \mathcal{M}$, write $m \leq n$ if, given that $m \in \mathcal{M}_\alpha$ and $n \in \mathcal{M}_\beta$, there is surjective $f : \beta \rightarrow \alpha$ such that $m(f(i), j) \leq n(i, j)$ for all $i < \beta$, $j < d$.

Theorem 5.16 For any sequence m_0, m_1, \dots in \mathcal{M} , there are $n < k < \omega$ with $m_n \leq m_k$. Indeed, (\mathcal{M}, \leq) is a (class) bqo.

Proof. Let $\xi^0, \xi^1, \dots \in \bigcup_\alpha Q^\alpha$ be such that $\theta(\xi^i) = m_i$ for $i < \omega$. Since $(\bigcup_\alpha Q^\alpha, \preceq)$ is a bqo and $\mathcal{B} = \{\{n\} : n < \omega\}$ is a barrier on ω , the map $(\{n\} \mapsto \xi^n) : \mathcal{B} \rightarrow \bigcup_\alpha Q^\alpha$ is not \preceq -bad, and so there are $n < k < \omega$ with $\xi^n \preceq \xi^k$.

Write ξ for ξ^n and η for ξ^k ; suppose that $\xi \in Q^\alpha$ and $\eta \in Q^\beta$. By $\xi \leq \eta$, there is a one-one order-preserving map $\varphi : \alpha \rightarrow \beta$ with $\xi_i \leq \eta_{\varphi(i)}$ for all

$i < \alpha$. By $\xi \sqsubseteq \eta$, for all $i < \beta$ there is $\psi(i) < \alpha$ with $\xi_{\psi(i)} \leq \eta_i$. Define $f : \beta \rightarrow \alpha$ by

$$f(i) = \begin{cases} \varphi^{-1}(i), & \text{if } i \in \text{rng}(\varphi), \\ \psi(i), & \text{otherwise.} \end{cases}$$

Then f is surjective, and $\xi_{f(i)} \leq \eta_i$ in Q for all $i < \beta$. Since $\theta(\xi) = m_n$ and $\theta(\eta) = m_k$, that is

$$m_n(f(i), j) = (\xi_{f(i)})_j \leq (\eta_i)_j = m_k(i, j) \text{ for all } i < \beta, j < d.$$

So $m_n \leq m_k$. The bqo part is an exercise. □

Probably one can generalise to higher dimensions.

6 Countable linear orderings

Notation 6.1 We will write $\alpha, \beta, \gamma, \dots$ for linear orderings. i, j, k, \dots will denote ordinals. We write $\alpha \leq \beta$ if α embeds into β in an order-preserving fashion. We write $\alpha < \beta$ if $\alpha \leq \beta$ and $\alpha \not\leq \alpha$.

Theorem 6.2 (Laver) *The class of countable linear orderings is a bqo under $<$.*

(Transcription of notes is incomplete. Please see [1, 3, 4] for the proofs.)

References

- [1] R. Laver, *Better-quasi-orderings and a class of trees*, Studies in Foundations and Combinatorics (Gian-Carlo Rota, ed.), Advances in Mathematics Supplementary Studies, vol. 1, Academic Press, 1978, pp. 31–48.
- [2] R. Mansfield and G. Weitzkamp, *Recursive aspects of descriptive set theory*, Oxford Logic Guides, vol. 11, Oxford University Press, 1985.
- [3] C. St. J. A. Nash-Williams, *On well-quasi-ordering infinite trees*, Proc. Cambr. Philos. Soc. **61** (1965), 697–720.
- [4] ———, *On better-quasi-ordering transfinite sequences*, Proc. Cambr. Philos. Soc. **64** (1968), 273–290.
- [5] S. G. Simpson, *Bqo theory and Fraïssé’s conjecture*, 1985, chapter 9 of [2].