# Finite model property for guarded fragments, and extending partial isomorphisms

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Joint work with M Otto

# Outline of talk

I aim to explain the context and proof of a recent theorem of M. Otto–I.H. on extending partial isomorphisms of relational structures.

- Definitions
- History
- Applications
- Guarded fragments
- Proof of theorem of Otto-I.H.

## Partial isomorphisms & automorphisms

**Definition 1** Let *L* be a relational signature, and *A* a *L*-structure.

1. A *partial isomorphism* of *A* is a partial 1–1 map  $p: A \rightarrow A$  such that for all *n*-ary  $R \in L$  and all  $a_1, \ldots, a_n \in \text{dom } p$ ,

 $A \models R(a_1, \ldots, a_n) \leftrightarrow R(p(a_1), \ldots, p(a_n)).$ 

- 2. An *automorphism* of *A* is a bijective partial isomorphism of *A*.
- 3. We write Aut *A* for the set (group) of automorphisms of *A*.

# The problem

Given:

- a finite relational signature *L*,
- a finite *L*-structure *A*,
- some partial isomorphisms  $p_1, \ldots, p_n$  of A.

Questions:

- 1. Can we find a *finite L*-structure  $B \supseteq A$  such that  $p_1, \ldots, p_n$  extend to automorphisms of *B*?
- 2. Can we find such a *B* with 'nice properties'?

We will see some applications later.

## **History**

- **0.** For sets  $(L = \emptyset)$ , this is easy.
- **1. J. Truss, 1992** Can extend a single partial isomorphism of a finite graph to an automorphism of a larger finite graph.
- 2. E. Hrushovski, 1992 Can extend all partial isomorphisms of a finite graph to automorphisms of a (single) larger finite graph.
- **3. B.** Herwig, 1995 For any finite relational signature *L* and any finite *L*-structure *A*, there is a finite *L*-structure  $B \supseteq A$  such that
  - 1. any partial isomorphism of A extends to an automorphism of B,
  - 2. for any  $b \in B$ , there is  $g \in Aut B$  with  $g(b) \in A$ ,
  - 3. if  $b_1, \ldots, b_n \in B$ , and  $B \models R(b_1, \ldots, b_n)$ for some *n*-ary  $R \in L$ , then there is some  $g \in Aut B$  with  $g(b_1), \ldots, g(b_n) \in A$ .

# Gaifman graph

To explain a later result of Herwig, we need a definition.

**Definition 2** Let *L* be a relational signature, and *A* an *L*-structure. The *Gaifman graph* Gaif(A) of *A* is the (undirected loop-free) graph defined by:

- its set of nodes is dom *A*,
- (x, y) is an edge iff there are *n*-ary  $R \in L$  and  $a_1, \ldots, a_n \in A$  with  $A \models R(a_1, \ldots, a_n),$  $x, y \in \{a_1, \ldots, a_n\}.$

Herwig's 1998 theorem

- **4. B. Herwig, 1998** For any finite *L*-structure *A*, can extend all partial isomorphisms of *A* to automorphisms of a finite *L*-structure  $B \supseteq A$  such that
  - 1. the 1995 properties hold,
  - 2. if *S* is an *L*-structure, Gaif(S) is a clique, and there is a homomorphism  $h : S \rightarrow B$ , then there is a homomorphism  $g : S \rightarrow A$ .

## Corollaries

- 1. For any  $n \ge 3$ , can extend all partial isomorphisms of a finite  $K_n$ -free graph to automorphisms of a larger finite  $K_n$ -free graph.
- 2. For any class  $\mathcal{T}$  of finite tournaments, can extend all partial isomorphisms of a finite digraph omitting all  $T \in \mathcal{T}$  to automorphisms of a larger finite digraph omitting all  $T \in \mathcal{T}$ .

# History ctd.

Hrushovski's proof was group-theoretic/combinatorial.

Herwig's papers greatly extended these methods.

5. Herwig–Lascar, 2000 — gave simpler and purely combinatorial proofs of Hrushovski's 1992 and Herwig's 1995 results, connected them to equivalent results in free groups, and extended the results.

A purely combinatorial account of Herwig's 1998 result was missing.

Also, even this result is still not strong enough for some applications.

# **Applications**

1. Small index property.

A countable structure M has this if any subgroup of Aut M of index  $< 2^{\omega}$  is open in Aut M(in the topology of pointwise convergence).

Hrushovski's result  $\vdash$  the 'random graph' has the small index property. Herwig  $\vdash$  s.i.p. for universal homogeneous  $K_n$ free graphs, and for Henson digraphs.

- Finite model theory: hierarchy theorems for fixedpoint logics (Grohe, 1996).
  Other work of Grohe too.
- Finite model property for *guarded fragments* of first-order logic, and classes of 'relativised' algebras in algebraic logic. Crs<sub>n</sub>, WA, etc.: Andréka, I.H., Németi, 1999. Guarded fragment: Grädel, 1999. More if we can strengthen Herwig's theorem...

## Guarded fragments: a rough guide

'Find out why modal logic is well-behaved (decidable etc), and generalise.'

Guarded fragments (Andréka, van Benthem, Németi, 1998) are 'modal' fragments of first-order logic.

- Any atomic formula is guarded.
- Guarded formulas are closed under booleans.
- If  $\varphi(\bar{x}, \bar{y})$  is guarded, and  $\gamma(\bar{x}, \bar{y})$  is a guard, then  $\exists \bar{y}(\gamma(\bar{x}, \bar{y}) \land \varphi(\bar{x}, \bar{y}))$  is guarded.

In the basic *guarded fragment*,  $\gamma$  must be atomic.

In the *loosely guarded fragment [van Benthem 1997]*,  $\gamma$  can be a conjunction of atomic formulas, if every y in  $\overline{y}$  and z in  $\overline{x}\overline{y}$  co-occur in a single conjunct.

In the *packed fragment [Marx, 2001],*  $\gamma$  can be a conjunction of atomic and existentially-quantified atomic formulas, if all distinct u, v in  $\overline{x}\overline{y}$  co-occur free in a single conjunct.

The guard enforces that  $\overline{x}\overline{y}$  is a clique in the Gaifman graph. The *clique-guarded fragment* [*Grädel*, 1999] does this directly.

Finite model property for guarded fragments

Guarded fragments are well-behaved: decidable in 2EXPTIME, 'back-and-forth' characterisation, some interpolation results, etc.

E. Grädel (1999) showed that the *basic guarded fragment* has the *finite model property:* any guarded sentence with a model has a finite model.

He used Herwig's 1995 theorem. This gives finite  $B \supseteq A$  such that whenever  $\overline{b} \in B$  satisfies a guard of the basic guarded fragment, there is  $g \in Aut B$  with  $g(\overline{b}) \in A$ .

But for LGF and PF/CGF, we need more.

## **Strengthening Herwig's 1998 theorem**

- 6. M. Otto–I.H., 2001 Let L be a relational signature, and A a finite L-structure. There exists a finite L-structure  $C \supseteq A$  such that:
  - 1. any partial isomorphism of A extends to an automorphism of C,
  - 2. if  $S \subseteq C$  is a clique in Gaif(C), then there is  $g \in Aut \ C$  with

$$g(S) \stackrel{\mathsf{def}}{=} \{g(x) : x \in S\} \subseteq A.$$

#### **Consequences for finite model property**

Whenever  $\overline{c} \in C$  satisfies a packed fragment guard, there is  $g \in Aut \ C$  with  $g(\overline{c}) \in A$ .

Hence can generalise Grädel's FMP proof to the *loosely guarded and packed fragments.* 

[Remark: I.H. 2001 proved these fragments have FMP by unpleasant hack of Herwig's proof.]

#### **Proof**

Take L, A as stated. By Herwig's 1995 theorem, there is a finite *L*-structure  $B \supseteq A$  such that

- 1. every partial isomorphism of A extends to an automorphism of B,
- 2. any  $x \in B$  is mapped into A by some  $g \in Aut B$ .
- If B = A, we are done. Assume that  $B \supset A$ .

**Definition 3** A set  $U \subseteq B$  is *small* if there is some  $g \in Aut B$  with  $g(U) \subseteq A$ , and *large* otherwise.

- *B* is large.
- If U is large then  $|U| \ge 2$ .
- If U is large and  $g \in Aut B$  then g(U) is large.

Write  $\mathcal{U}$  for the set of large subsets of B.

Note:  $\mathcal{U} = \{g(U) : U \in \mathcal{U}\}$  for all  $g \in Aut B$ .

#### **Domain of** C

**Definition 4** Let  $b \in B$ . A map  $\chi : \mathcal{U} \to \omega$  is said to be a *b*-valuation if for all  $U \in \mathcal{U}$ :

- if  $b \notin U$  then  $\chi(U) = 0$ ,
- if  $b \in U$  then  $1 \leq \chi(U) < |U|$ . (Note  $|U| \geq 2$ .)

View  $\chi$  as a notion  $[b \in U] = \chi(U)$ , for large U.

Value is 0 if  $b \notin U$ . Value is positive (many-valued logic!) if  $b \in U$ .

**Definition 5** We let *C* have domain

 $\{(b, \chi) : b \in B, \chi \text{ a } b\text{-valuation}\}.$ 

We'll define the *L*-structure of C in a minute.

**Definition 6** Also define the projection  $\pi : C \to B$ by  $\pi(b, \chi) = b$ .

#### **Generic sets**

**Definition 7** A set  $S \subseteq C$  is said to be *generic* if for all distinct  $(b, \chi), (c, \psi) \in S$ , we have

1.  $b \neq c$ , 2.  $\chi(U) \neq \psi(U)$  for all  $U \in \mathcal{U}$  with  $b, c \in U$ .

A set is generic iff each two-element subset is generic.

**Lemma 8** Let  $S \subseteq C$  be generic. Then  $\pi(S)$  is small.

**Proof.** If  $\pi(S) = U \in \mathcal{U}$ , then by genericity,

- $\pi \upharpoonright S$  is 1–1, so |U| = |S|,
- the map  $\theta:S \to \{1,2,\ldots,|U|-1\}$  given by  $\theta(b,\chi)=\chi(U)$

is 1–1.

This is impossible.

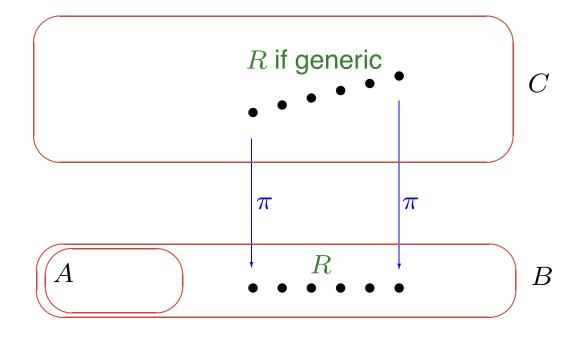
#### The structure of C

**Definition 9** Define *C* as an *L*-structure as follows.

If  $R \in L$  is *n*-ary, and  $(b_1, \chi_1), \ldots, (b_n, \chi_n) \in C$ , then let  $C \models R((b_1, \chi_1), \ldots, (b_n, \chi_n))$  iff

1.  $\{(b_1, \chi_1), \dots, (b_n, \chi_n)\}$  is generic,

**2.** 
$$B \models R(b_1, ..., b_n).$$



#### **Embedding** A into C

Lemma 10 *A embeds into C*.

**Proof.** Let  $U \in \mathcal{U}$ . So there is no  $g \in Aut B$  with  $g(U) \subseteq A$ .

Then  $U \not\subseteq A$ . So  $|U \cap A| < |U|$ .

Enumerate  $U \cap A$  as  $\{a_1^U, \ldots, a_n^U\}$ , with n < |U|. Do this for all  $U \in \mathcal{U}$ .

For  $a \in A$ , define an *a*-valuation  $\chi_a : \mathcal{U} \to \omega$  by  $\chi_a(U) = \begin{cases} 0, & \text{if } a \notin U, \\ \text{the } i \text{ such that } a = a_i^U, & \text{otherwise.} \end{cases}$ 

Now define  $\nu : A \to C$  by  $\nu(a) = (a, \chi_a)$ .

Note that  $\nu(A) = \{\nu(a) : a \in A\}$  is generic. So  $\nu : A \to C$  is an *L*-embedding.

We can therefore *replace* A by  $\nu(A) \cong A$ , and prove the theorem for it.

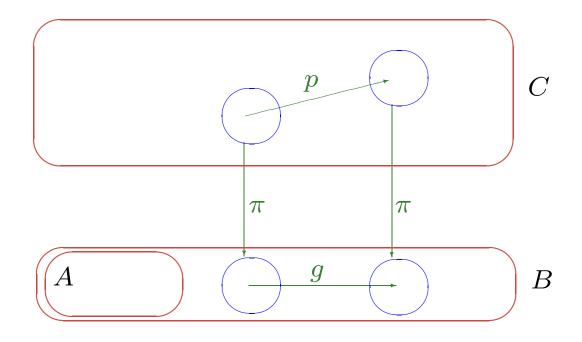
This will be easy after a definition and a lemma.

#### **Compatible maps**

**Definition 11** Let  $p : C \to C$  be a 1–1 partial map, and let  $g \in Aut B$ . We say that p is *g*-compatible if for all  $(b, \chi) \in dom p$  we have

 $p(b,\chi) = (g(b),\chi')$  for some  $\chi'$ .

(That is,  $\pi \circ p \subseteq g \circ \pi$ .)



#### Main lemma

**Lemma 12** Let  $p : C \to C$  be a 1–1 partial map with generic domain and range. Let  $g \in Aut B$ , and suppose that p is g-compatible. Then p extends to some g-compatible  $\hat{p} \in Aut C$ .

**Proof.** As dom p is generic, can write its elements as  $(b, \chi_b)$ . Suppose  $p(b, \chi_b) = (g(b), \chi'_{g(b)})$ , say. We need to define  $\hat{p} : (b, \chi) \mapsto (g(b), \chi')$  for all  $(b, \chi) \in C$ .

Fix a large set  $U \in \mathcal{U}$ . Then the set of pairs

 $\{ \langle \chi_b(U), \chi'_{g(b)}(g(U)) \rangle : (b, \chi_b) \in \text{dom } p \}$ is a 1–1 partial map on  $|U| = \{0, 1, \dots, |U| - 1\}$ , fixing 0 if defined on it.

Extend it to a permutation  $\theta_U$  of |U|, fixing 0. Do this for all  $U \in \mathcal{U}$ .

Define  $\hat{p}(b,\chi) = (g(b),\chi')$ , where

 $\chi'(g(U)) = \theta_U(\chi(U))$  for  $U \in \mathcal{U}$ .

Then  $\hat{p}$  extends p, and  $\hat{p}$  is g-compatible. Also,  $\hat{p} \in$  Aut C (because  $\hat{p}$  is g-compatible and preserves generic sets).

## Checking that C is as required

- 1. Certainly,  $C \supseteq \nu(A)$  and C is finite.
- 2. Let p be a partial isomorphism of  $\nu(A)$ . We need to extend it to  $\hat{p} \in \text{Aut } C$ .

Let  $p \downarrow = \nu^{-1} \circ p \circ \nu$  be the corresponding partial isomorphism of *A*. By Herwig's theorem, we can extend  $p \downarrow$  to some  $g \in \text{Aut } B$ .

Clearly, *p* is *g*-compatible. And dom *p*, rng *p* are generic (as  $\subseteq \nu(A)$ ).

By lemma 12, p extends to some (g-compatible)  $\hat{p} \in Aut C$ .

## Mapping cliques back into $\nu(A)$

**3.** Let  $S \subseteq C$  be a clique in Gaif(C). We want  $g \in Aut \ C$  with  $g(S) \subseteq \nu(A)$ .

S is generic.

So by lemma 8,  $\pi(S)$  is small.

So there is  $g \in Aut B$  with  $g(\pi(S)) \subseteq A$ .

The map

 $p: x \mapsto \nu(g(\pi(x))) \in \nu(A) \quad (\text{for } x \in S)$ 

is 1–1, and has generic domain (S) and range  $(\subseteq \nu(A))$ , and is *g*-compatible.

By lemma 12, p extends to  $\hat{p} \in \text{Aut } C$ , and  $\hat{p}(S) \subseteq \nu(A)$ .

# **Conclusion**

- The theorem strengthens Herwig's 1998 results.
- Combined with the combinatorial proof of Herwig's 1995 result by Herwig–Lascar, it gives a purely combinatorial proof of them.
- New and simple proof of finite model property for loosely guarded and packed (and clique-guarded) fragments.
  (Otto has a variant argument to give this, using finite model property of the basic guarded frag-

finite model property of the basic guarded fragment.)