332
Advanced Computer Architecture
Chapter 2

Caches and Memory Systems

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These lecture notes are partly based on the course text, Hennessy and Patterson's *Computer Architecture, a quantitative approach (3rd and 4th eds),* and on the lecture slides of David Patterson and John Kubiatowicz's Berkeley course

Review: Cache performance

- Miss-oriented Approach to Memory Access:

$$CPUtime = IC \times \left(CPI_{Execution} + \frac{MemAccess}{Inst} \times MissRate \times MissPenalt\ y \right) \times CycleTime$$

$$CPUtime = IC \times \left(CPI_{Execution} + \frac{MemMisses}{Inst} \times MissPenalt\ y \right) \times CycleTime$$

  - $CPI_{Execution}$ includes ALU and Memory instructions

- Separating out Memory component entirely
  - AMAT = Average Memory Access Time
  - CPI_{ALUOps} does not include memory instructions

$$CPUtime = IC \times \left(\frac{AluOps}{Inst} \times CPI_{AluOps} + \frac{MemAccess}{Inst} \times AMAT \right) \times CycleTime$$

$$\begin{aligned} AMAT &= HitTime + MissRate \times MissPenalt\ y \\ &= \left(HitTime_{Inst} + MissRate_{Inst} \times MissPenalt\ y_{Inst}\right) + \\ &\quad \left(HitTime_{Data} + MissRate_{Data} \times MissPenalt\ y_{Data}\right) \end{aligned}$$

Average memory access time:

$$AMAT = HitTime + MissRate \times MissPenalt\ y$$

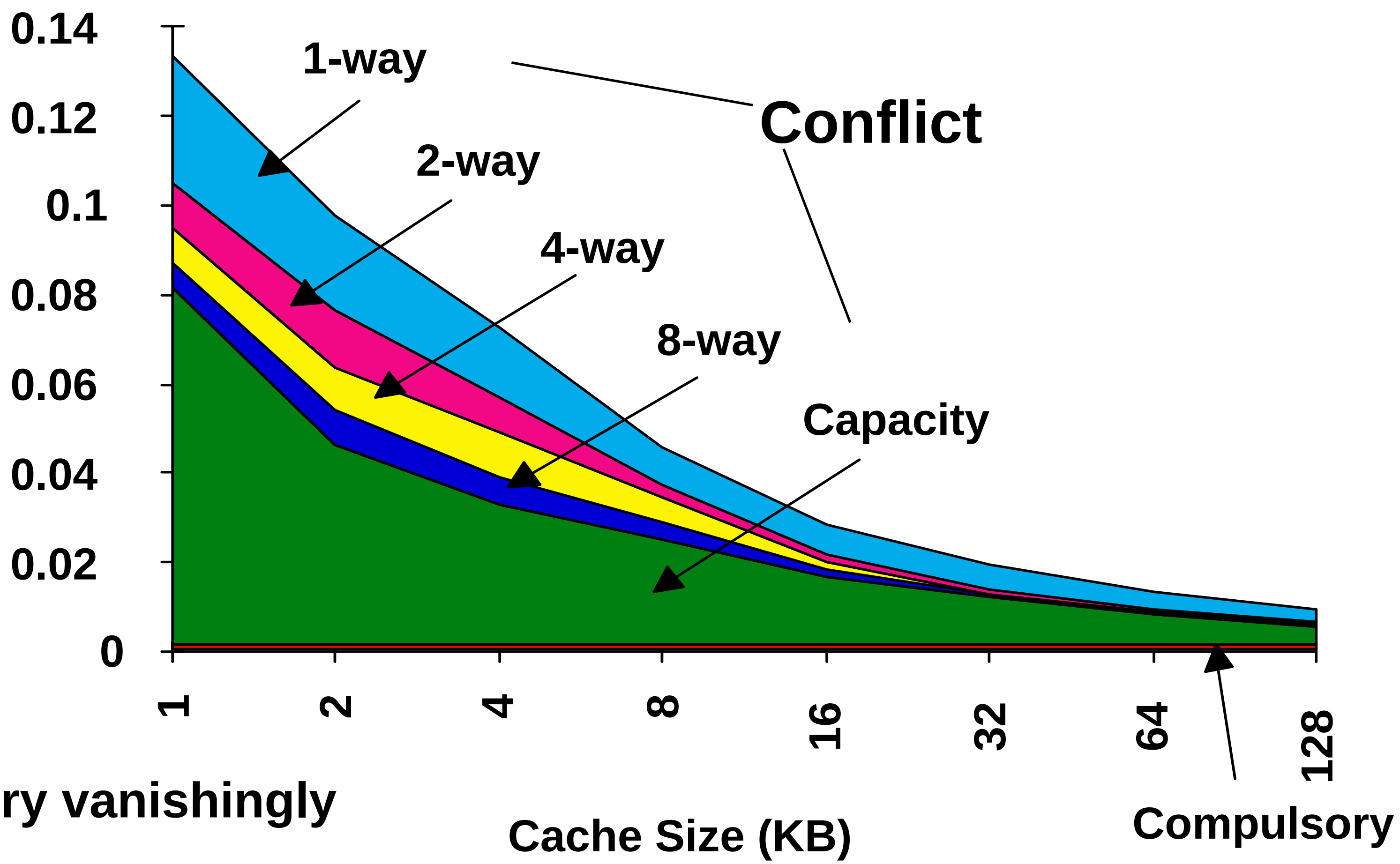
There are three ways to improve cache performance:

1. Reduce the miss rate,
2. Reduce the miss penalty, or
3. Reduce the time to hit in the cache.

Reducing Misses

- Classifying Misses: 3 Cs
  - *Compulsory*—The first access to a block is not in the cache, so the block must be brought into the cache. Also called *cold start misses* or *first reference misses*.
    (Misses in even an Infinite Cache)
  - *Capacity*—If the cache cannot contain all the blocks needed during execution of a program, capacity misses will occur due to blocks being discarded and later retrieved.
    (Misses in Fully Associative Size X Cache)
  - *Conflict*—If block-placement strategy is set associative or direct mapped, conflict misses (in addition to compulsory & capacity misses) will occur because a block can be discarded and later retrieved if too many blocks map to its set. Also called *collision misses* or *interference misses*.
    (Misses in N-way Associative, Size X Cache)
- More recent, 4th "C":
  - *Coherence* - Misses caused by cache coherence.

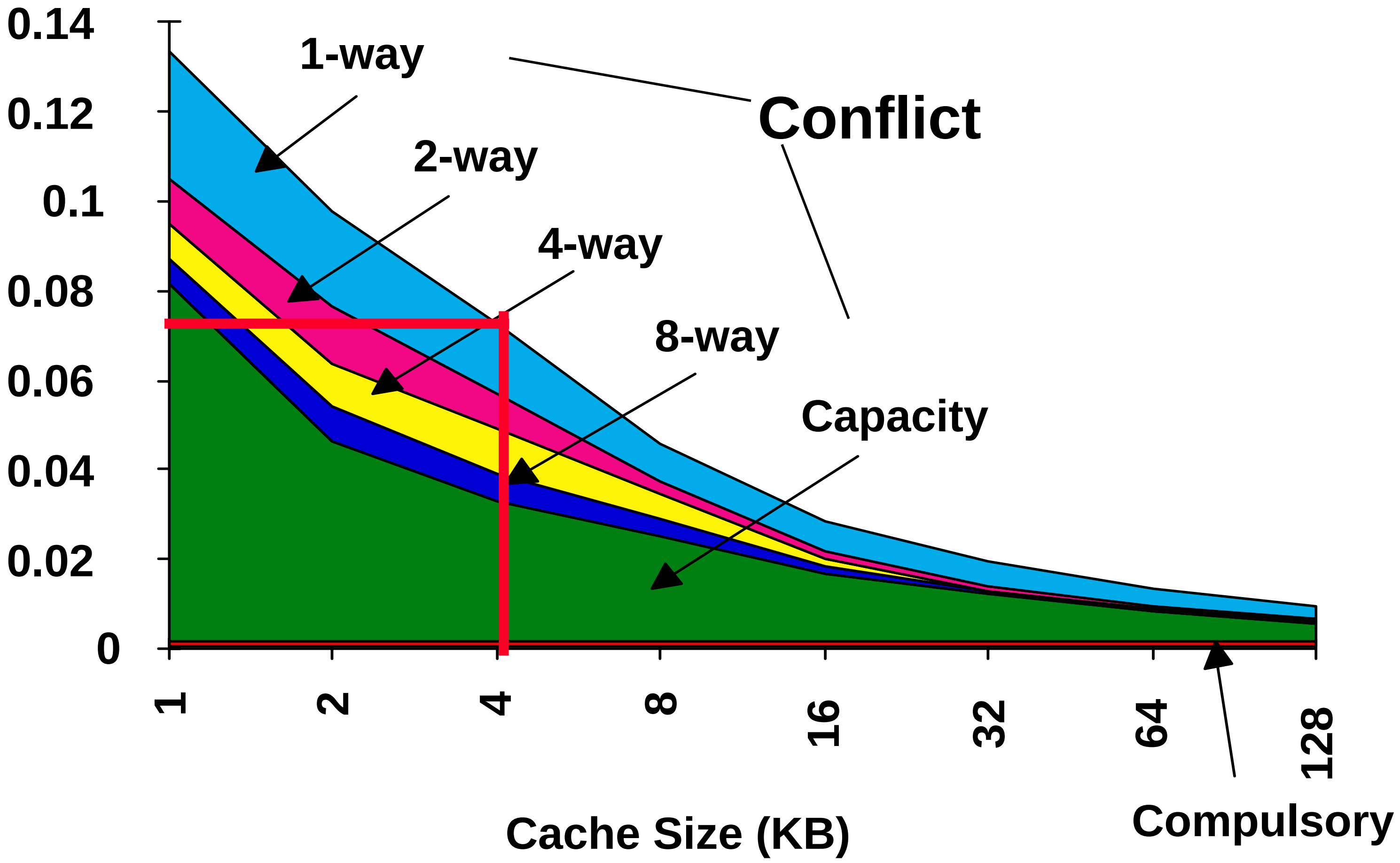
3Cs Absolute Miss Rate (SPEC92)



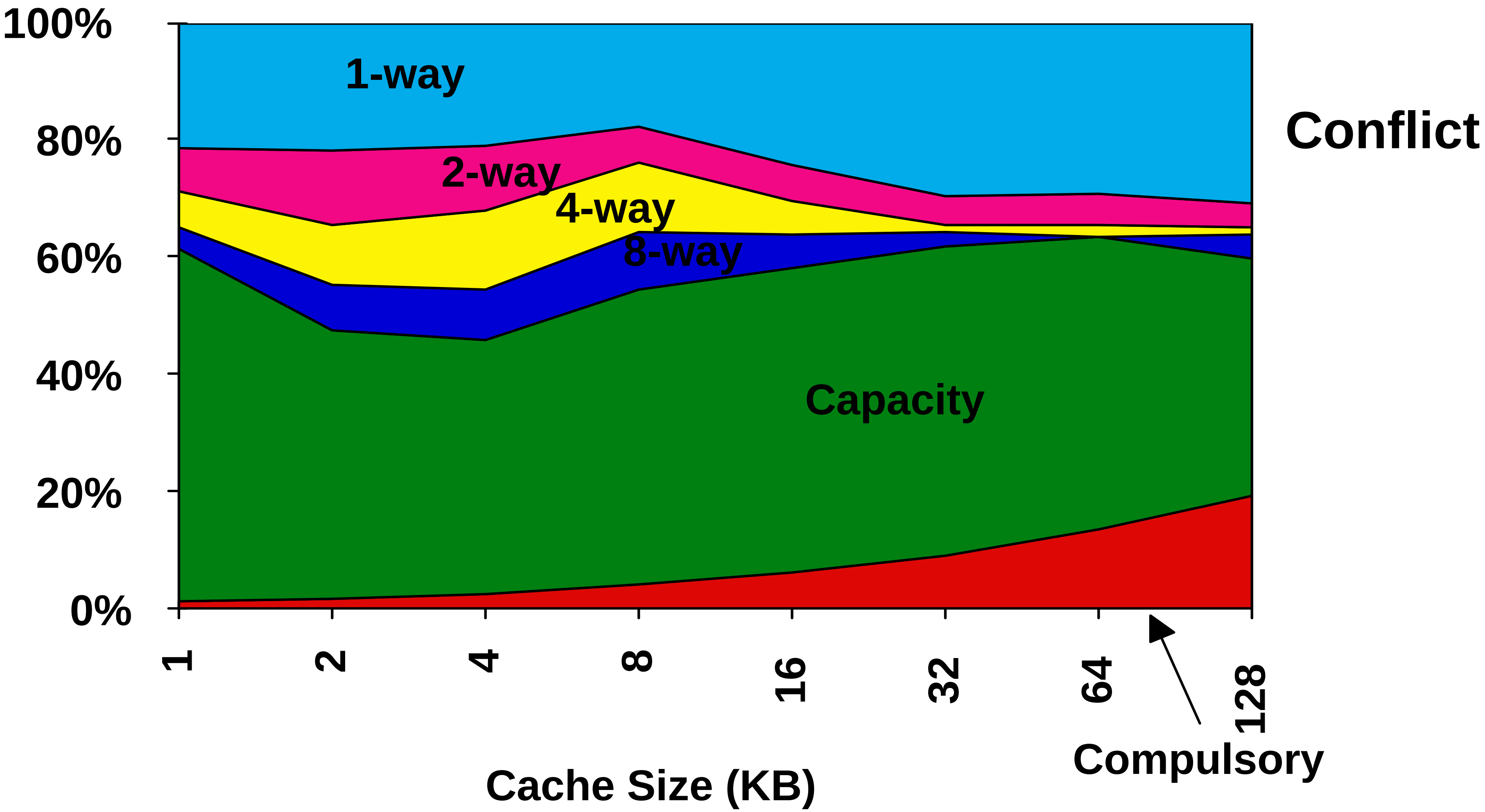
Compulsory vanishingly small

2:1 Cache Rule (of thumb!)

miss rate 1-way associative cache size X
= miss rate 2-way associative cache size X/2



3Cs Relative Miss Rate



Flaws: for fixed block size
Good: insight => invention

How We Can Reduce Misses?

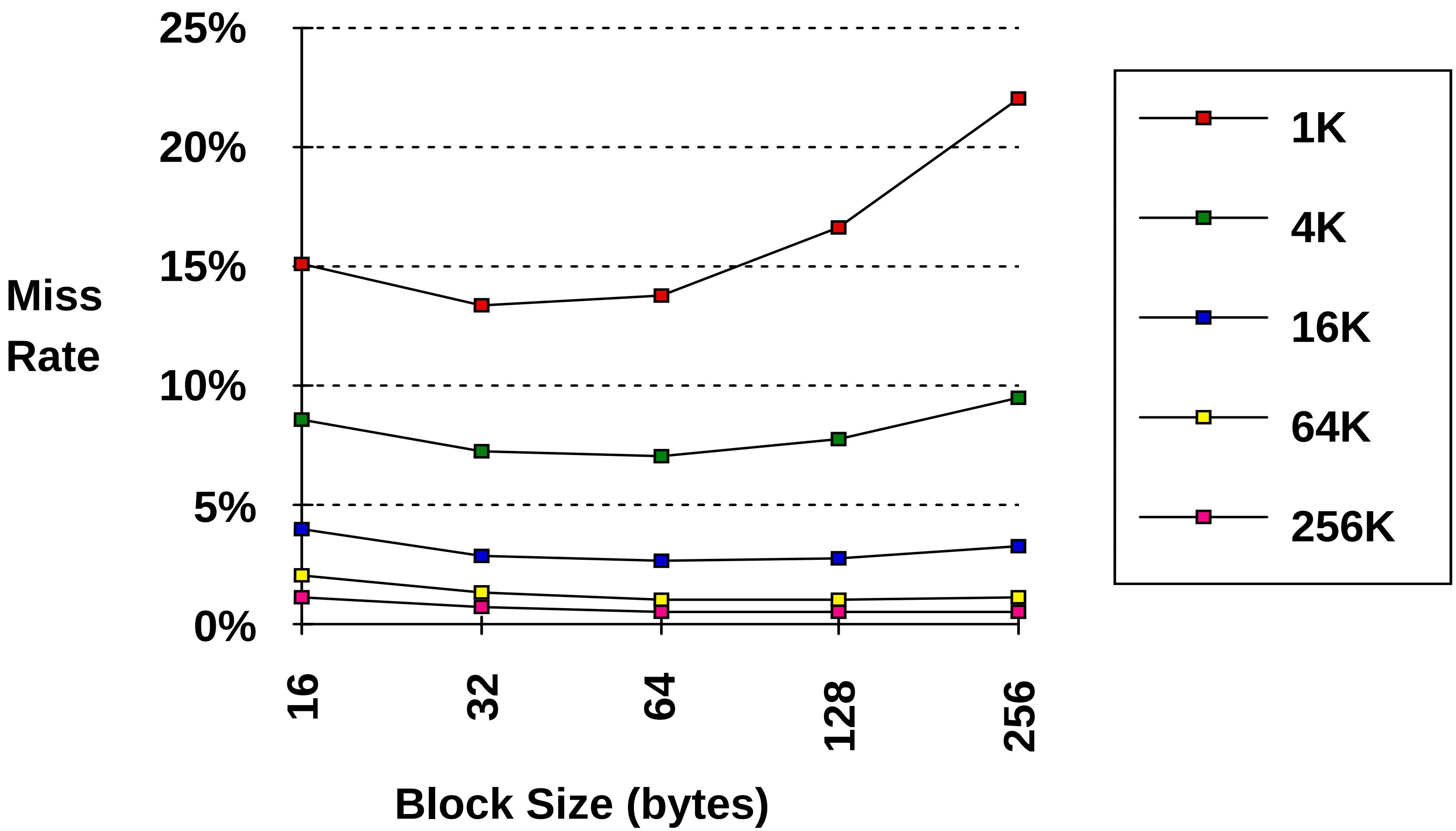
- 3 Cs: Compulsory, Capacity, Conflict
- In all cases, assume total cache size not changed:
- What happens if:

1) Change Block Size:
Which of 3Cs is obviously affected?

2) Change Associativity:
Which of 3Cs is obviously affected?

3) Change Compiler:
Which of 3Cs is obviously affected?

1. Reduce Misses via Larger Block Size



2. Reduce Misses via Higher Associativity

- 2:1 Cache Rule of thumb:
  - The Miss Rate of a direct-mapped cache of size N
  - Is the same as the Miss Rate of a 2-way set-associative cache size of size N/2
    - on average, over a large suite of benchmarks
- Beware: Execution time is the only thing that matters in the end
  - Will Clock Cycle time increase?
  - Hill [1988] suggested hit time for 2-way vs. 1-way
external cache +10%,
internal + 2%

Example: Avg. Memory Access Time vs. Miss Rate

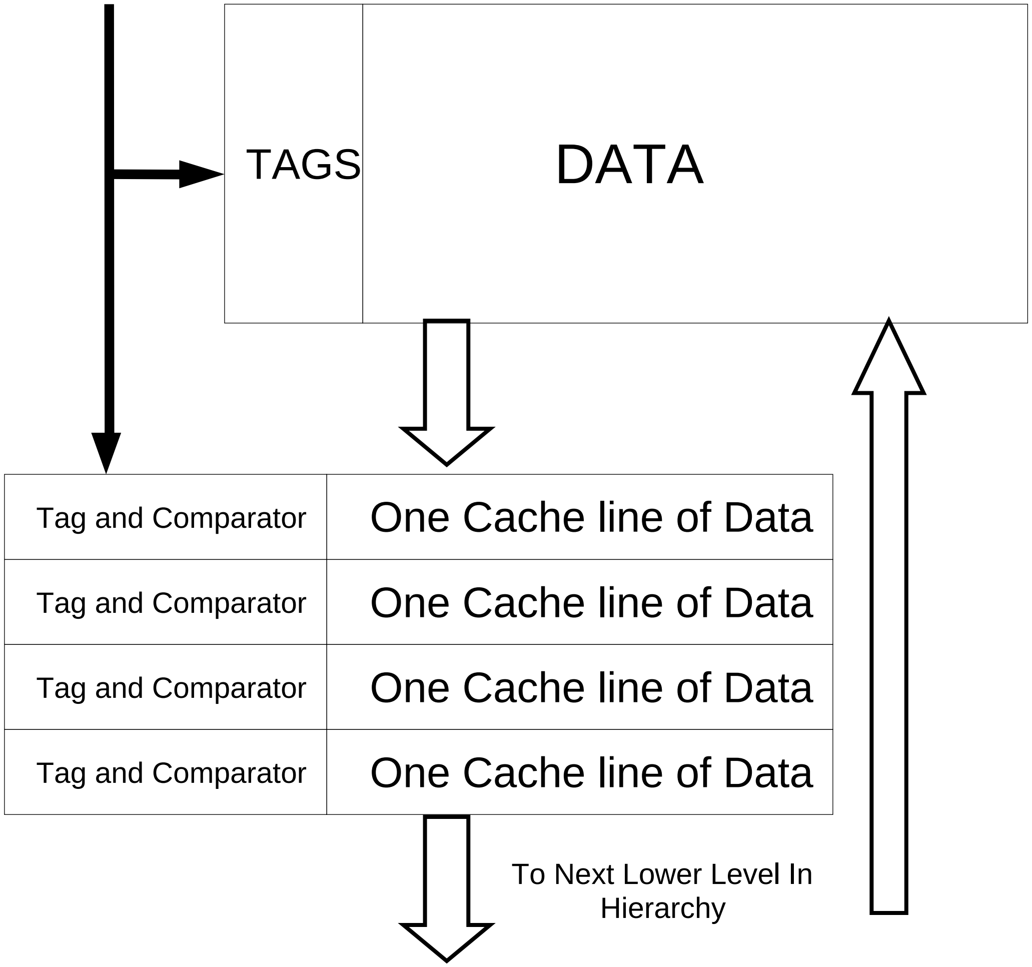
- Example: assume clock cycle time (CCT) = 1.10 for 2-way, 1.12 for 4-way, 1.14 for 8-way vs. CCT direct mapped

Cache Size		Associativity		
(KB)	1-way	2-way	4-way	8-way
1	2.33	2.15	2.07	2.01
2	1.98	1.86	1.76	1.68
4	1.72	1.67	1.61	1.53
8	1.46	1.48	1.47	1.43
16	1.29	1.32	1.32	1.32
32	1.20	1.24	1.25	1.27
64	1.14	1.20	1.21	1.23
128	1.10	1.17	1.18	1.20

(Red means A.M.A.T. not improved by more associativity)

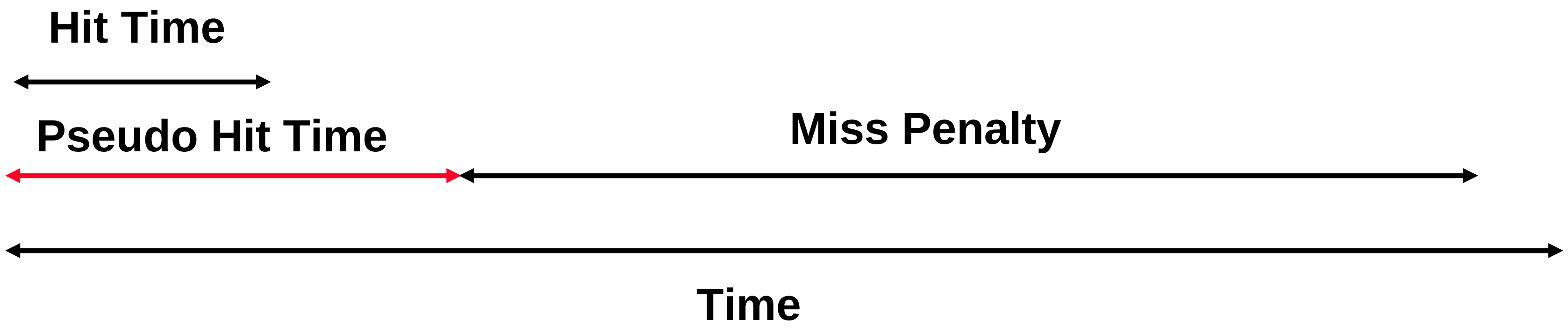
3. Reducing Misses via a "Victim Cache"

- How to combine fast hit time of direct mapped yet still avoid conflict misses?
- Add buffer to place data discarded from cache
- Jouppi [1990]: 4-entry victim cache removed 20% to 95% of conflicts for a 4 KB direct mapped data cache
- Used in AMD Opteron, Alpha, HP machines



4. Reducing Misses via "Pseudo-Associativity"

- How to combine fast hit time of Direct Mapped and have the lower conflict misses of 2-way SA cache?
- Divide cache: on a miss, check other half of cache to see if there, if so have a pseudo-hit (slow hit)



- Drawback: CPU pipeline is hard if hit takes 1 or 2 cycles
  - Better for caches not tied directly to processor (L2)
  - Used in MIPS R10000 L2 cache, similar in UltraSPARC

5. Reducing Misses by Hardware Prefetching of Instructions & Data

- E.g., Instruction Prefetching
  - Alpha 21064 fetches 2 blocks on a miss
  - Extra block placed in "stream buffer"
  - On miss check stream buffer
- Works with data blocks too:
  - Jouppi [1990] 1 data stream buffer got 25% misses from 4KB cache; 4 streams got 43%
  - Palacharla & Kessler [1994] for scientific programs for 8 streams got 50% to 70% of misses from 2 64KB, 4-way set associative caches
- Prefetching relies on having extra memory bandwidth that can be used without penalty

6. Reducing Misses by Software Prefetching Data

- Data Prefetch
  - Load data into register (HP PA-RISC loads)
  - Cache Prefetch: load into cache (MIPS IV, PowerPC, SPARC v. 9)
  - Special prefetching instructions cannot cause faults; a form of speculative execution
- Prefetching comes in two flavors:
  - Binding prefetch: Requests load directly into register.
    - Must be correct address and register!
  - Non-Binding prefetch: Load into cache.
    - Can be incorrect. Frees HW/SW to guess!
- Issuing Prefetch Instructions takes time
  - Is cost of prefetch issues < savings in reduced misses?
  - Higher superscalar reduces difficulty of issue bandwidth

7. Reducing Misses by Compiler Optimizations

- McFarling [1989]* reduced instruction cache misses by 75% on 8KB direct mapped cache, 4 byte blocks in software
- Instructions
  - By choosing instruction memory layout based on callgraph, branch structure and profile data
  - Reorder procedures in memory so as to reduce conflict misses
- Similar (but different) ideas work for data
  - *Merging Arrays*: improve spatial locality by single array of compound elements vs. 2 arrays
  - *Loop Interchange*: change nesting of loops to access data in order stored in memory
  - *Loop Fusion*: Combine 2 independent loops that have same looping and some variables overlap
  - *Blocking*: Improve temporal locality by accessing "blocks" of data repeatedly vs. going down whole columns or rows

* "Program optimization for instruction caches", ASPLOS89, http://doi.acm.org/10.1145/70082.68200

Merging Arrays Example

```
/* Before: 2 sequential arrays */
int val[SIZE];
int key[SIZE];

/* After: 1 array of stuctures */
struct merge {
  int val;
  int key;
};
struct merge merged_array[SIZE];
```

Reducing conflicts between val & key; improve spatial locality

Loop Interchange Example

```
/* Before */
for (k = 0; k < 100; k = k+1)
  for (j = 0; j < 100; j = j+1)
      for (i = 0; i < 5000; i = i+1)
            x[i][j] = 2 * x[i][j];
/* After */
for (k = 0; k < 100; k = k+1)
  for (i = 0; i < 5000; i = i+1)
      for (j = 0; j < 100; j = j+1)
            x[i][j] = 2 * x[i][j];
```

Sequential accesses instead of striding through memory every 100 words; improved spatial locality

Loop Fusion Example

```
/* Before */
for (i = 0; i < N; i = i+1)
  for (j = 0; j < N; j = j+1)
      a[i][j] = 1/b[i][j] * c[i][j];
for (i = 0; i < N; i = i+1)
  for (j = 0; j < N; j = j+1)
      d[i][j] = a[i][j] + c[i][j];
/* After fusion */
for (i = 0; i < N; i = i+1)
  for (j = 0; j < N; j = j+1)
  {   a[i][j] = 1/b[i][j] * c[i][j];
      d[i][j] = a[i][j] + c[i][j];}
```

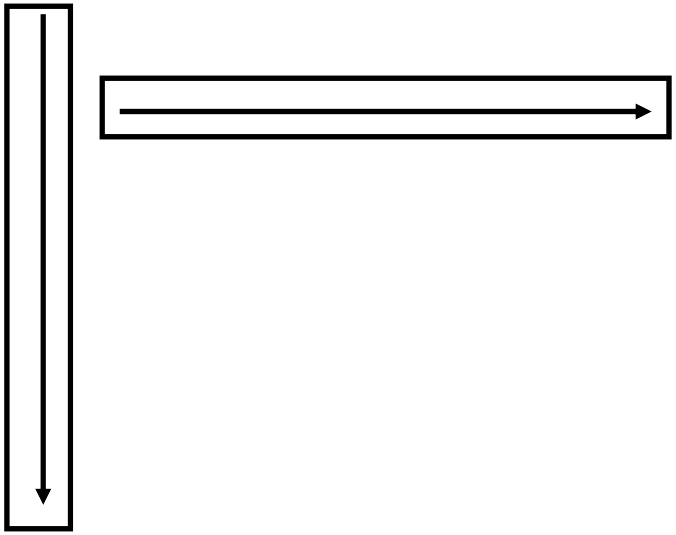
2 misses per access to a & c vs. one miss per access; improve spatial locality

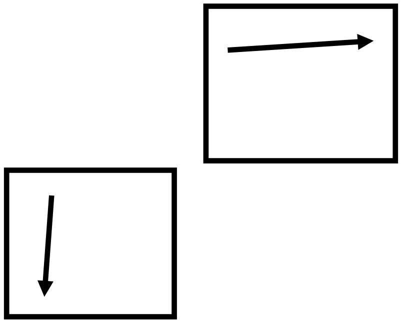
```
/* After array contraction */
for (i = 0; i < N; i = i+1)
  for (j = 0; j < N; j = j+1)
  {   c = c[i][j];
      a = 1/b[i][j] * c;
      d[i][j] = a + c;}
```

The real payoff comes if fusion enables Array Contraction: values transferred in scalar instead of via array

Blocking Example

```
/* Before */
for (i = 0; i < N; i = i+1)
  for (j = 0; j < N; j = j+1)
     {r = 0;
      for (k = 0; k < N; k = k+1){
        r = r + y[i][k]*z[k][j];};
      x[i][j] = r;
     };
```





- Two Inner Loops:
  - Read all NxN elements of z[]
  - Read N elements of 1 row of y[] repeatedly
  - Write N elements of 1 row of x[]
- Capacity Misses a function of N & Cache Size:
  - $2N^3 + N^2$ => (assuming no conflict; otherwise …)
- Idea: compute on BxB submatrix that fits

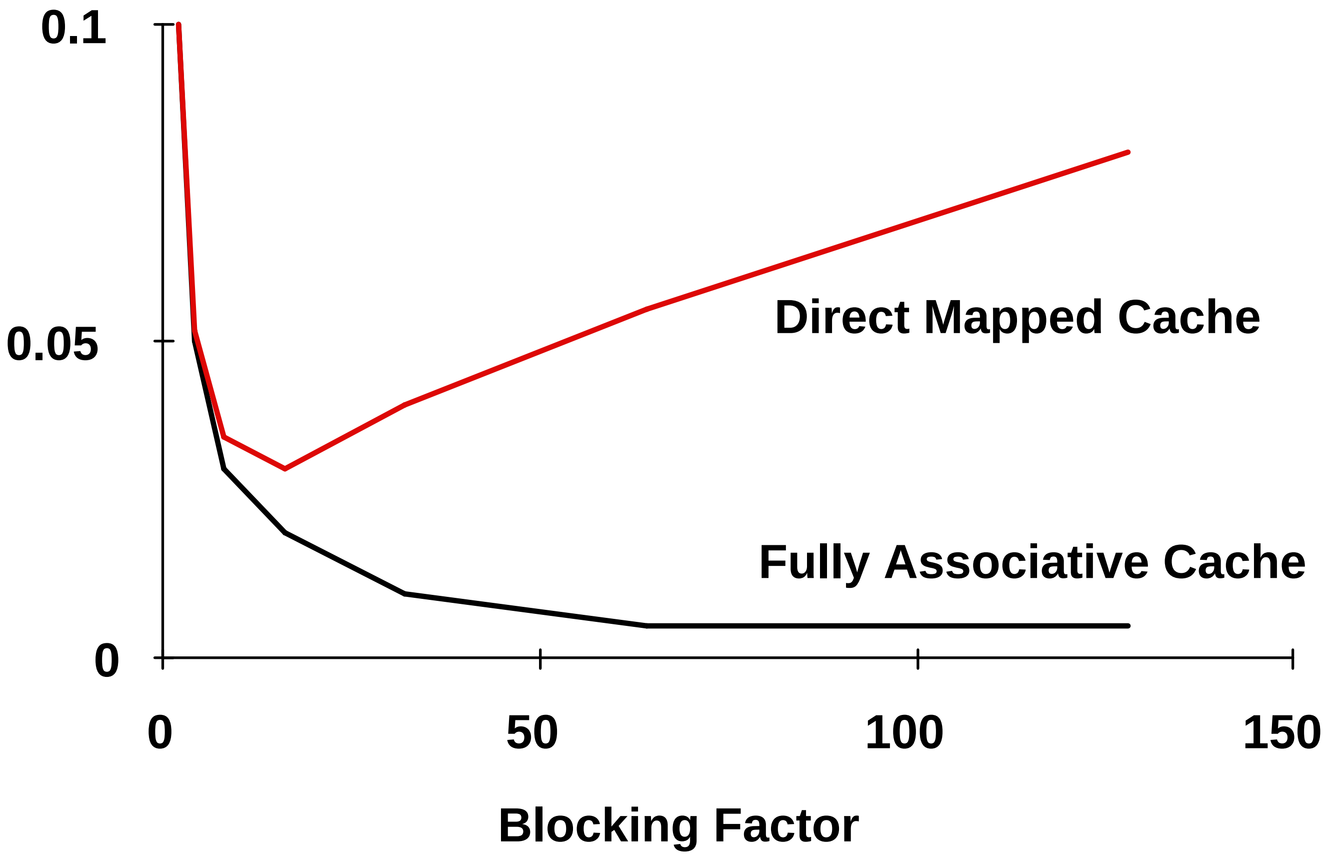
Blocking Example

```
/* After */
for (jj = 0; jj < N; jj = jj+B)
for (kk = 0; kk < N; kk = kk+B)
for (i = 0; i < N; i = i+1)
   for (j = jj; j < min(jj+B-1,N); j = j+1)
     {r = 0;
      for (k = kk; k < min(kk+B-1,N); k = k+1) {
        r = r + y[i][k]*z[k][j];};
      x[i][j] = x[i][j] + r;
     };
```

- B called *Blocking Factor*
- Capacity Misses from $2N^3 + N^2$ to $N^3/B+2N^2$
- Conflict Misses Too?

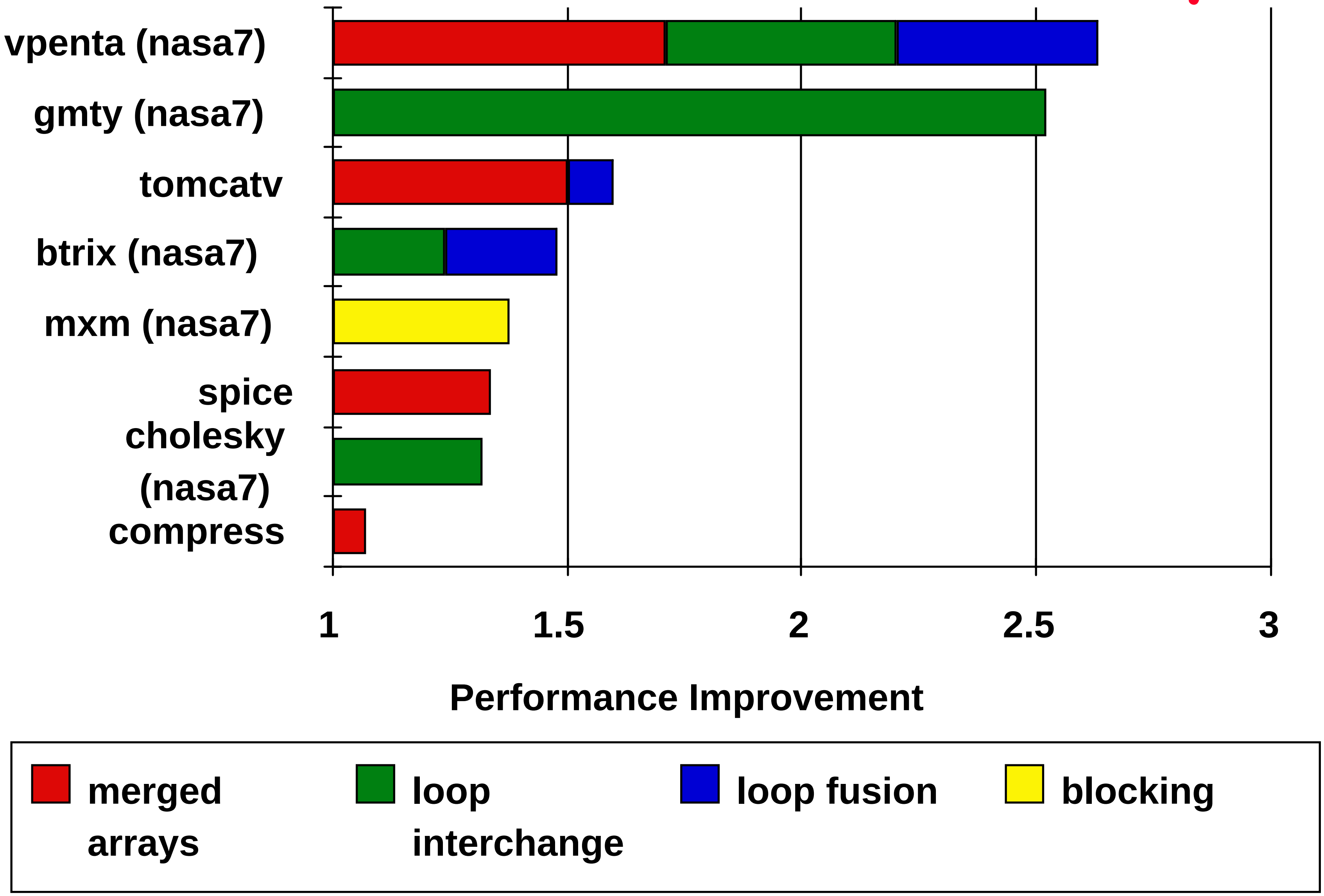
(We return to this example and this technique in Chapter 5)

Reducing Conflict Misses by Blocking



- Conflict misses in caches not FA vs. Blocking size
  - Lam et al [1991] a blocking factor of 24 had a fifth the misses vs. 48 despite both fit in cache

Summary of Compiler Optimizations to Reduce Cache Misses (by hand)



Summary: Miss Rate Reduction

$$CPUtime = IC \times \left(CPI_{Execution} + \frac{Memory\quad accesses}{Instruction} \times \textbf{Miss rate} \times Miss\quad penalty \right) \times Clock\quad cycle\quad time$$

- 3 Cs: Compulsory, Capacity, Conflict
  1. Reduce Misses via Larger Block Size
  2. Reduce Misses via Higher Associativity
  3. Reducing Misses via Victim Cache
  4. Reducing Misses via Pseudo-Associativity
  5. Reducing Misses by HW Prefetching Instr, Data
  6. Reducing Misses by SW Prefetching Data
  7. Reducing Misses by Compiler Optimizations
- Prefetching comes in two flavors:
  - Binding prefetch: Requests load directly into register.
    - Must be correct address and register!
  - Non-Binding prefetch: Load into cache.
    - Can be incorrect. Frees HW/SW to guess!

Review: Improving Cache Performance

1. Reduce the miss rate,
2. ***Reduce the miss penalty,*** or
3. Reduce the time to hit in the cache.

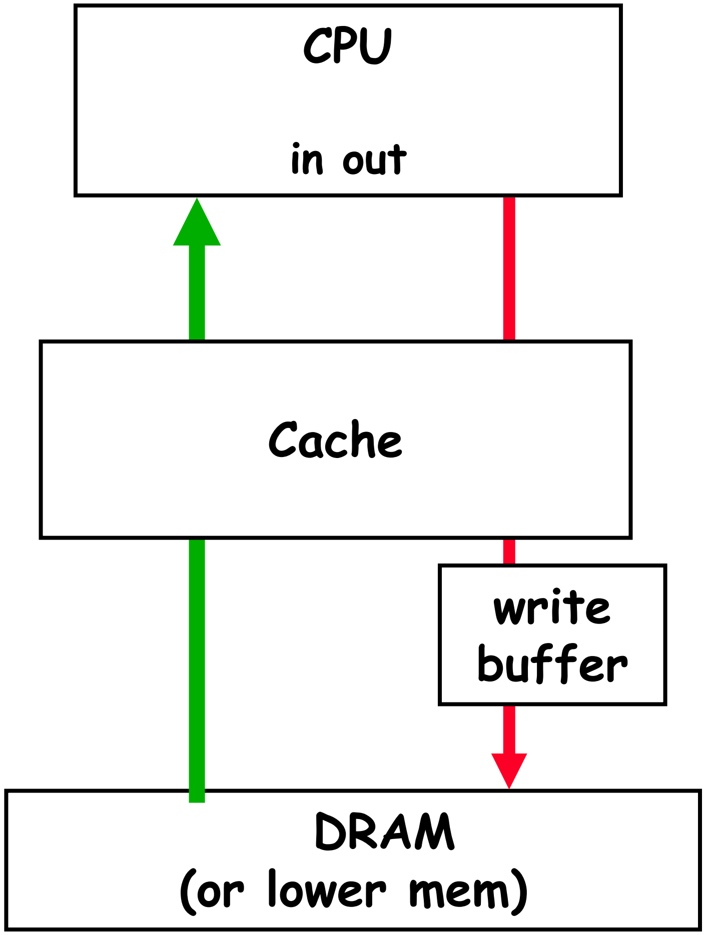
Write Policy: Write-Through vs Write-Back

- Write-through: all writes update cache and underlying memory/cache
  - Can always discard cached data - most up-to-date data is in memory
  - Cache control bit: only a *valid* bit
- Write-back: all writes simply update cache
  - Can't just discard cached data - may have to write it back to memory
  - Cache control bits: both *valid* and *dirty* bits
- Other Advantages:
  - Write-through:
    - memory (or other processors) always have latest data
    - Simpler management of cache
  - Write-back:
    - much lower bandwidth, since data often overwritten multiple times
    - Better tolerance to long-latency memory?

Write Policy 2: Write Allocate vs Non-Allocate (What happens on write-miss)

- Write allocate: allocate new cache line in cache
  - Usually means that you have to do a "read miss" to fill in rest of the cache-line!
  - Alternative: per/word valid bits
- Write non-allocate (or "write-around"):
  - Simply send write data through to underlying memory/cache - don't allocate new cache line!

1. Reducing Miss Penalty: Read Priority over Write on Miss



- Consider write-through with write buffers
  - RAW conflicts with main memory reads on cache misses
    - Could simply wait for write buffer to empty, before allowing read
    - Risks serious increase in read miss penalty (old MIPS 1000 by 50%)
    - Solution:
      - Check write buffer contents before read; if no conflicts, let the memory access continue
- Write-back also needs buffer to hold displaced blocks
  - Read miss replacing dirty block
  - Normal: Write dirty block to memory, and then do the read
  - Instead copy the dirty block to a write buffer, then do the read, and then do the write
  - CPU stall less since restarts as soon as do read

2. Reduce Miss Penalty: Early Restart and Critical Word First

- Don't wait for full block to be loaded before restarting CPU
  - *Early restart*—As soon as the requested word of the block ar□rives, send it to the CPU and let the CPU continue execution
  - *Critical Word First*—Request the missed word first from memory and send it to the CPU as soon as it arrives; let the CPU continue execution while filling the rest of the words in the block. Also called *wrapped fetch* and *requested word first*
- Generally useful only in large blocks,
- (Access to contiguous sequential words is very common – but doesn't benefit from either scheme – are they worthwhile?)

3. Reduce Miss Penalty: Non-blocking Caches to reduce stalls on misses

- *Non-blocking cache* or *lockup-free cache* allows data cache to continue to supply cache hits during a miss
  - requires full/empty bits on registers or out-of-order execution
  - requires multi-bank memories
- "*hit under miss*" reduces the effective miss penalty by working during miss instead of ignoring CPU requests
- "*hit under multiple miss*" or "*miss under miss*" may further lower the effective miss penalty by overlapping multiple misses
  - Significantly increases the complexity of the cache controller as there can be multiple outstanding memory accesses
  - Requires multiple memory banks (otherwise cannot support)
  - Pentium Pro allows 4 outstanding memory misses

Compare:
prefetching: overlap memory access with pre-miss instructions,
Non-blocking cache: overlap memory access with post-miss instructions

What happens on a Cache miss?

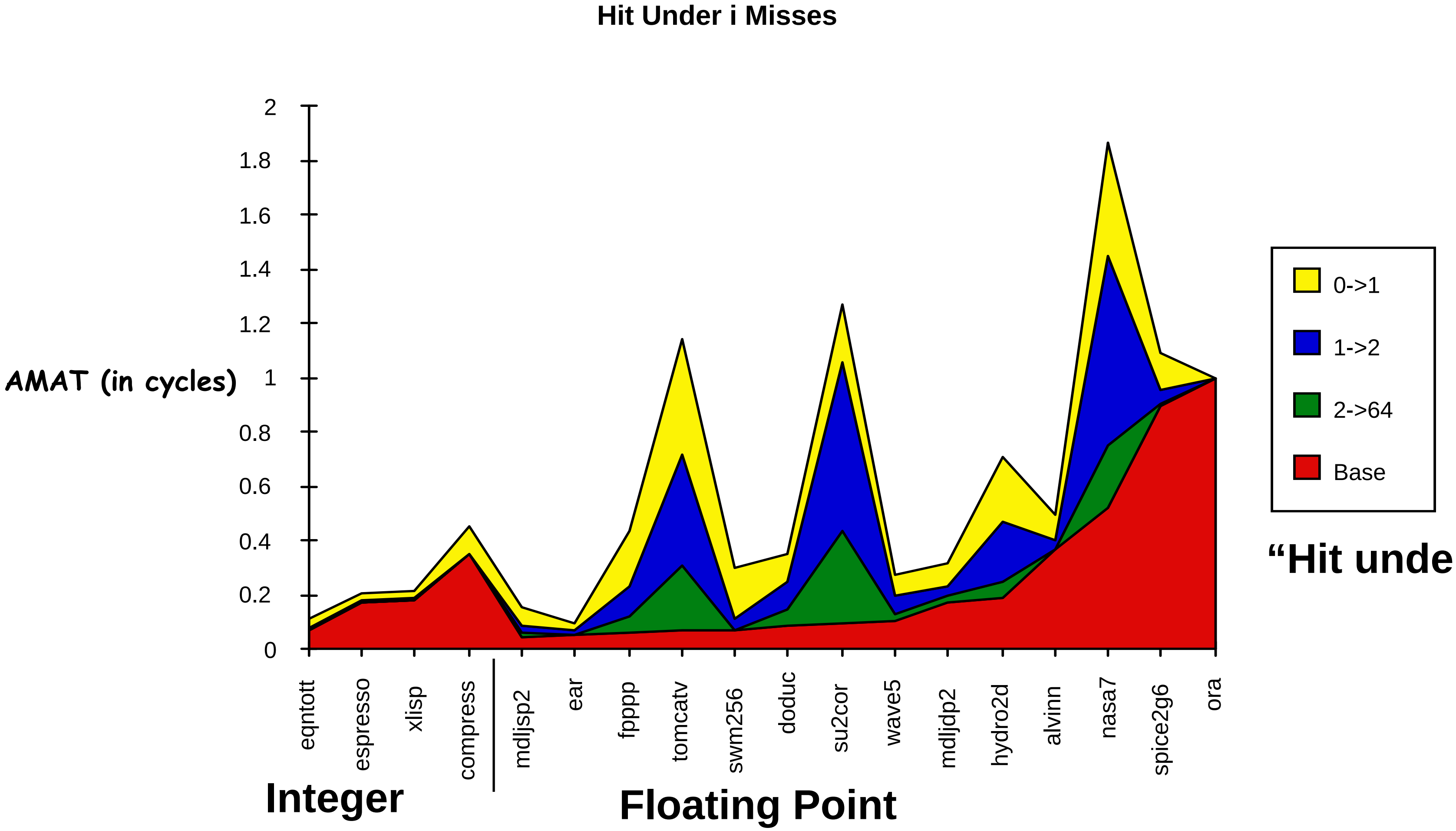
- For in-order pipeline, 2 options:
  - Freeze pipeline in Mem stage (popular early on: Sparc, R4000)

```
IF   ID   EX   Mem stall stall stall … stall Mem   Wr
     IF   ID   EX  stall stall stall … stall stall Ex Wr
```

  - Use Full/Empty bits in registers + MSHR queue
    - MSHR = "Miss Status/Handler Registers" (Kroft*)
      Each entry in this queue keeps track of status of outstanding memory requests to one complete memory line.
      - Per cache-line: keep info about memory address.
      - For each word: register (if any) that is waiting for result.
      - Used to "merge" multiple requests to one memory line
    - New load creates MSHR entry and sets destination register to "Empty". Load is "released" from pipeline.
    - Attempt to use register before result returns causes instruction to block in decode stage.
    - Limited "out-of-order" execution with respect to loads.
      Popular with in-order superscalar architectures.
- Out-of-order pipelines already have this functionality built in… (load queues, etc).

* David Kroft, Lockup-free instruction fetch/prefetch cache organization, ICCA81
http://portal.acm.org/citation.cfm?id=801868

Value of Hit Under Miss for SPEC



0->1
1->2
2->64
Base

"Hit under n Misses"

- FP programs on average: AMAT= 0.68 -> 0.52 -> 0.34 -> 0.26
- Int programs on average: AMAT= 0.24 -> 0.20 -> 0.19 -> 0.19
- 8 KB Data Cache, Direct Mapped, 32B block, 16 cycle miss

4: Add a second-level cache

- L2 Equations

  $\text{AMAT} = \text{Hit Time}_{L1} + \text{Miss Rate}_{L1} \times \text{Miss Penalty}_{L1}$

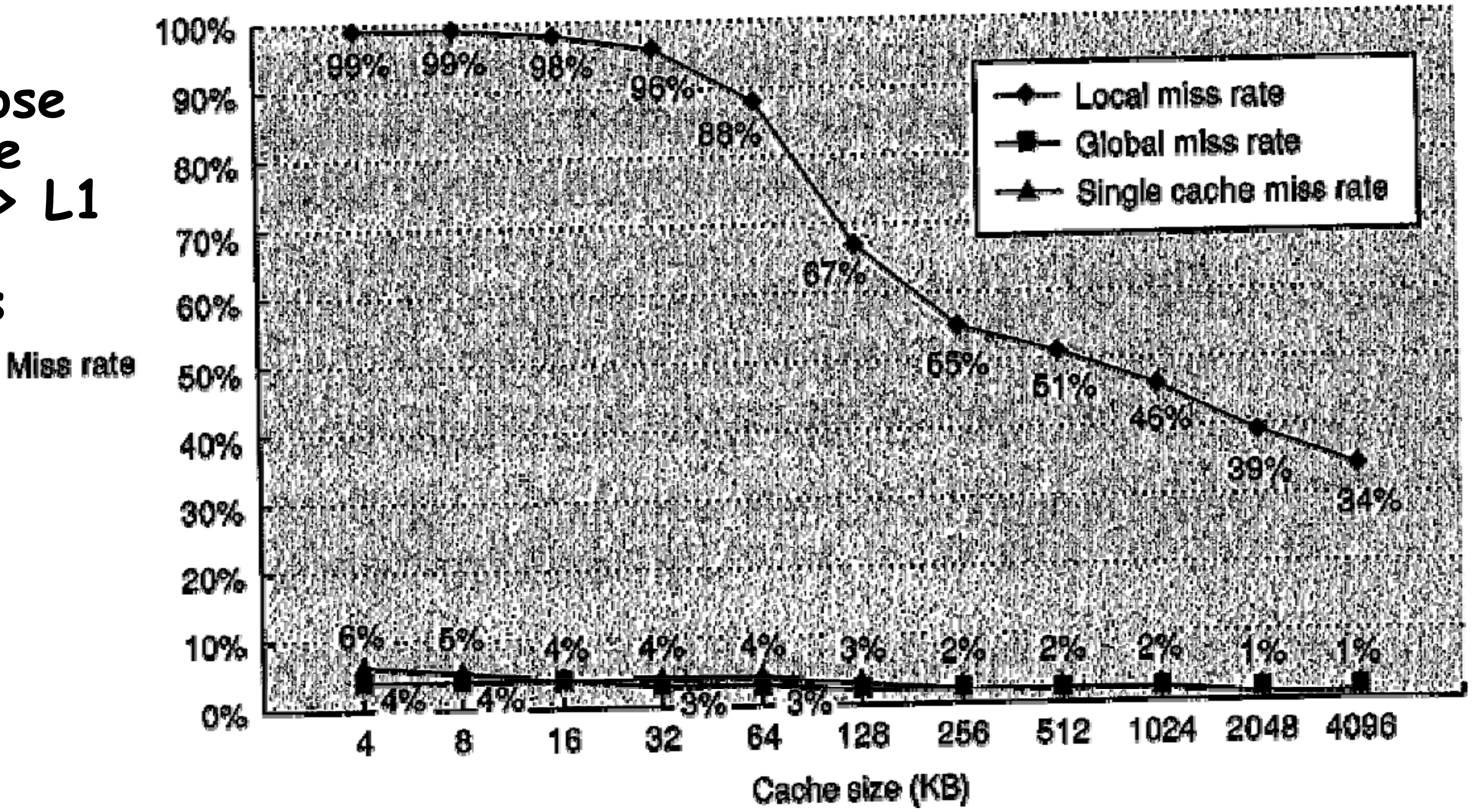
  $\text{Miss Penalty}_{L1} = \text{Hit Time}_{L2} + \text{Miss Rate}_{L2} \times \text{Miss Penalty}_{L2}$

  $$\text{AMAT} = \text{Hit Time}_{L1} + \underline{\text{Miss Rate}_{L1}} \times (\text{Hit Time}_{L2} + \underline{\text{Miss Rate}_{L2}} + \text{Miss Penalty}_{L2})$$

- Definitions:
  - *Local miss rate*— misses in this cache divided by the total number of memory accesses *to this cache* (Miss rate$_{L2}$)
  - *Global miss rate*—misses in this cache divided by the total number of memory accesses *generated by the CPU*
    (Miss Rate$_{L1}$ x Miss Rate$_{L2}$)
  - Global Miss Rate is what matters

Comparing Local and Global Miss Rates

- 32 KByte 1st level cache;
  Increasing 2nd level cache
- Global miss rate close to single level cache rate provided L2 >> L1
- Don't use local miss rate
- L2 not tied to CPU clock cycle!
- Cost & A.M.A.T.
- Generally Fast Hit Times and fewer misses
- Since hits are few, target miss reduction



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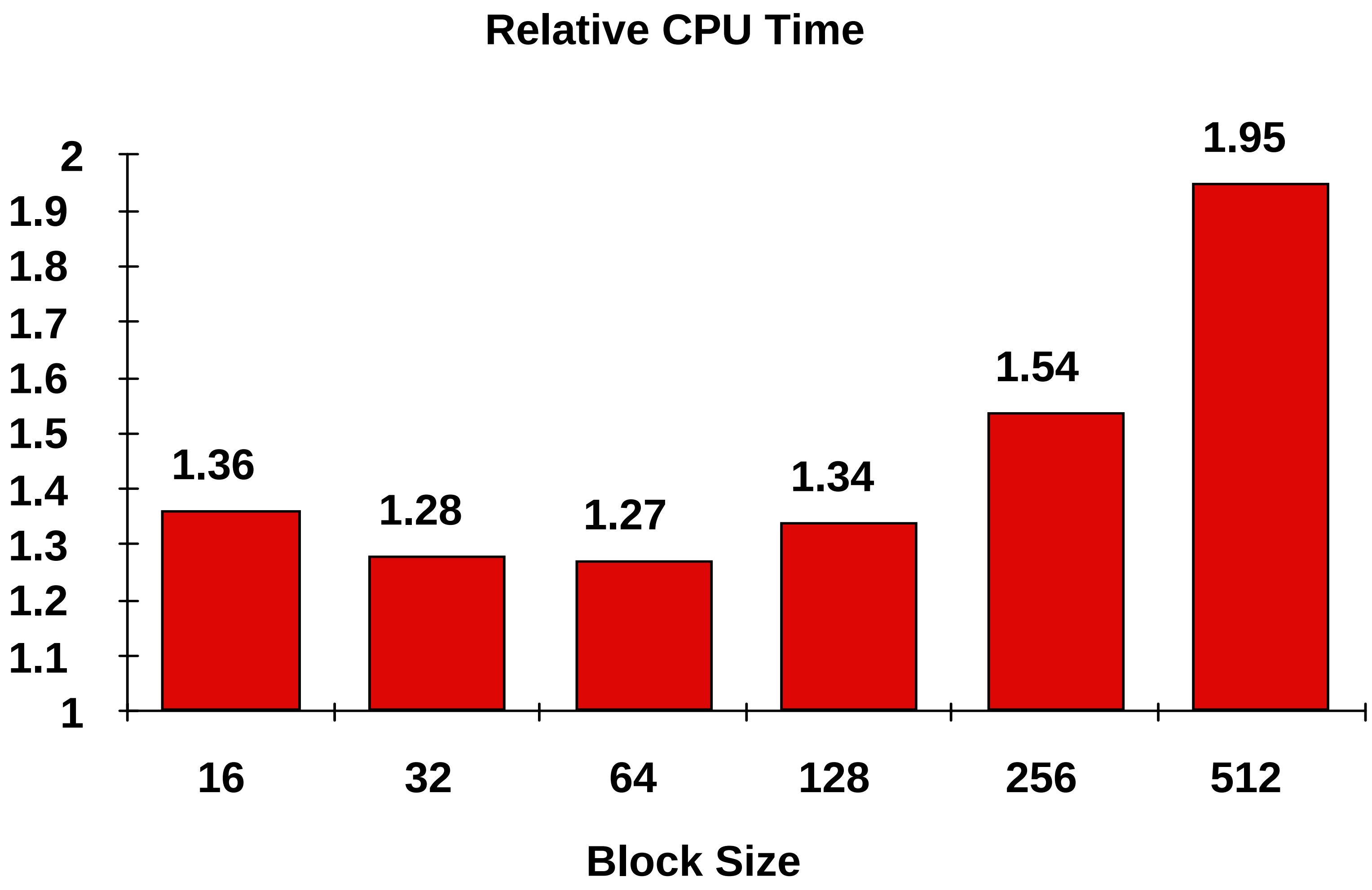
Reducing Misses: Which apply to L2 Cache?

- Reducing Miss Rate
  1. Reduce Misses via Larger Block Size
  2. Reduce Conflict Misses via Higher Associativity
  3. Reducing Conflict Misses via Victim Cache
  4. Reducing Conflict Misses via Pseudo-Associativity
  5. Reducing Misses by HW Prefetching Instr, Data
  6. Reducing Misses by SW Prefetching Data
  7. Reducing Capacity/Conf. Misses by Compiler Optimizations

L2 cache block size & A.M.A.T.



- 32KB L1, 8 byte path to memory

Reducing Miss Penalty Summary

$$CPUtime = IC \times \left(CPI_{Execution} + \frac{Memory\quad accesses}{Instruction} \times \mathbf{Miss\ rate} \times \boxed{Miss\quad penalty} \right) \times Clock\quad cycle\quad time$$

- Four techniques
  - Read priority over write on miss
  - Early Restart and Critical Word First on miss
  - Non-blocking Caches (Hit under Miss, Miss under Miss)
  - Second Level Cache
- Can be applied recursively to Multilevel Caches
  - Danger is that time to DRAM will grow with multiple levels in between
  - First attempts at L2 caches can make things worse, since increased worst case is worse

Average memory access time:

$$AMAT = HitTime + MissRate \times MissPenalt\ y$$

There are three ways to improve cache performance:

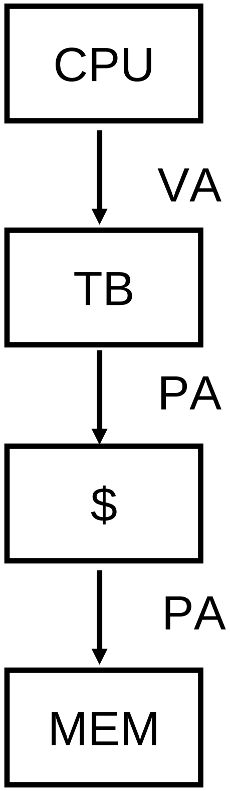
1. Reduce the miss rate,
2. Reduce the miss penalty, or
3. Reduce the time to hit in the cache.

Reducing the time to hit in the cache

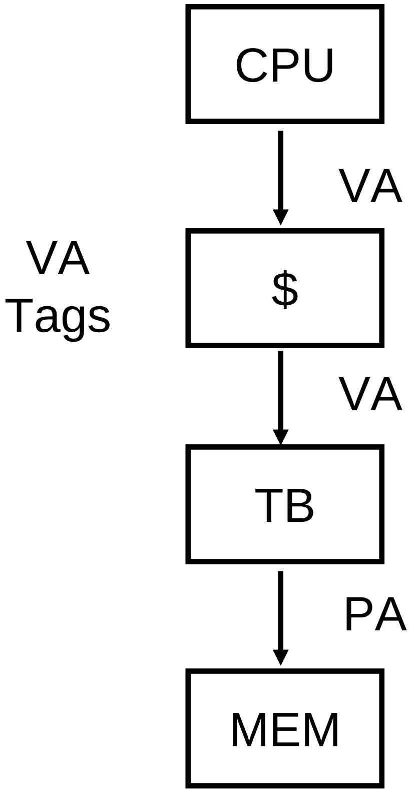
- Why does the Alpha 21164 have 8KB Instruction and 8KB data cache + 96KB second level cache, all on-chip?

1. Keep the cache small and simple
2. Keep address translation off the critical path
3. Pipeline the cache access

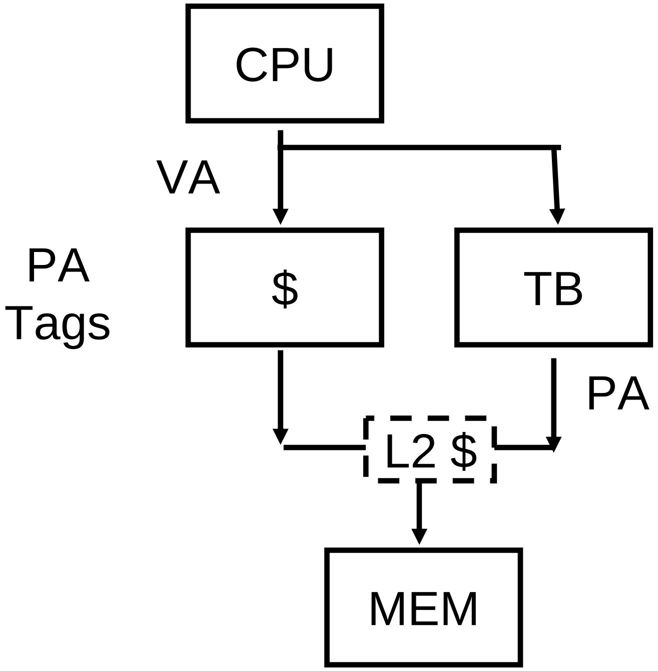
2. Fast hits by Avoiding Address Translation



Conventional
Organization



Virtually Addressed Cache
Translate only on miss
Synonym Problem

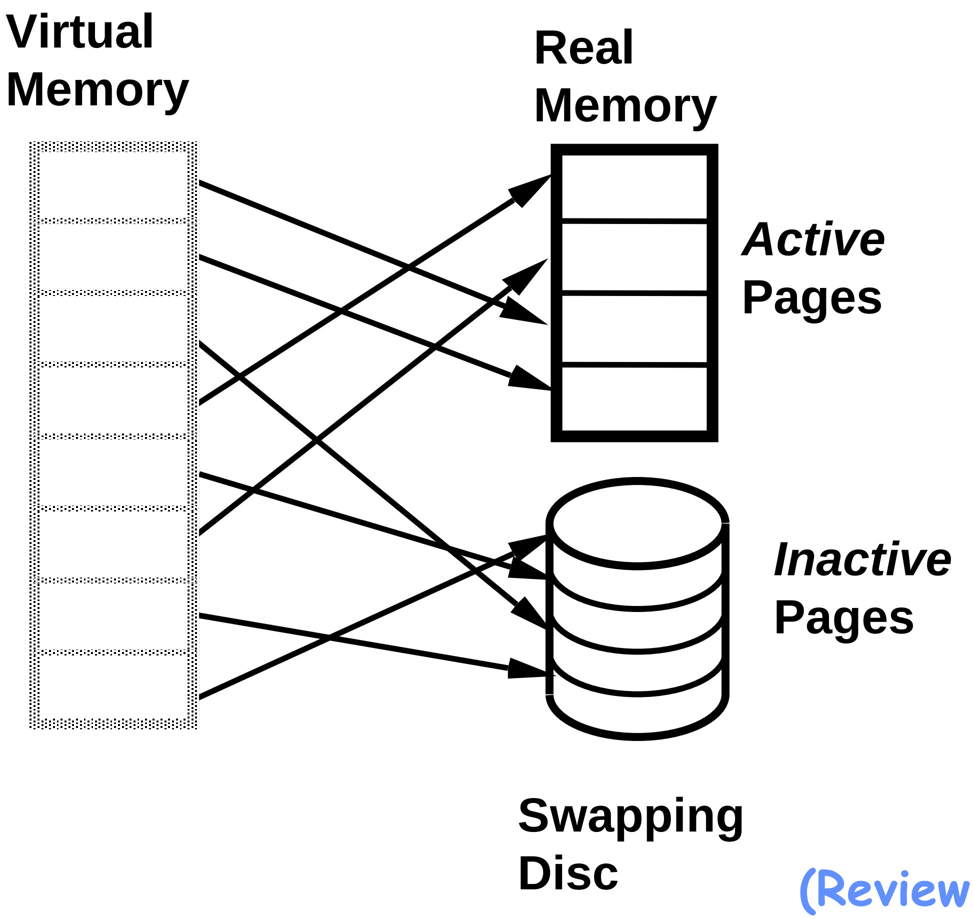


Overlap $ access
with VA translation:
requires $ index to
remain invariant
across translation

Paging

Virtual address space is divided into ***pages*** of equal size.

Main Memory is divided into ***page frames*** the same size.



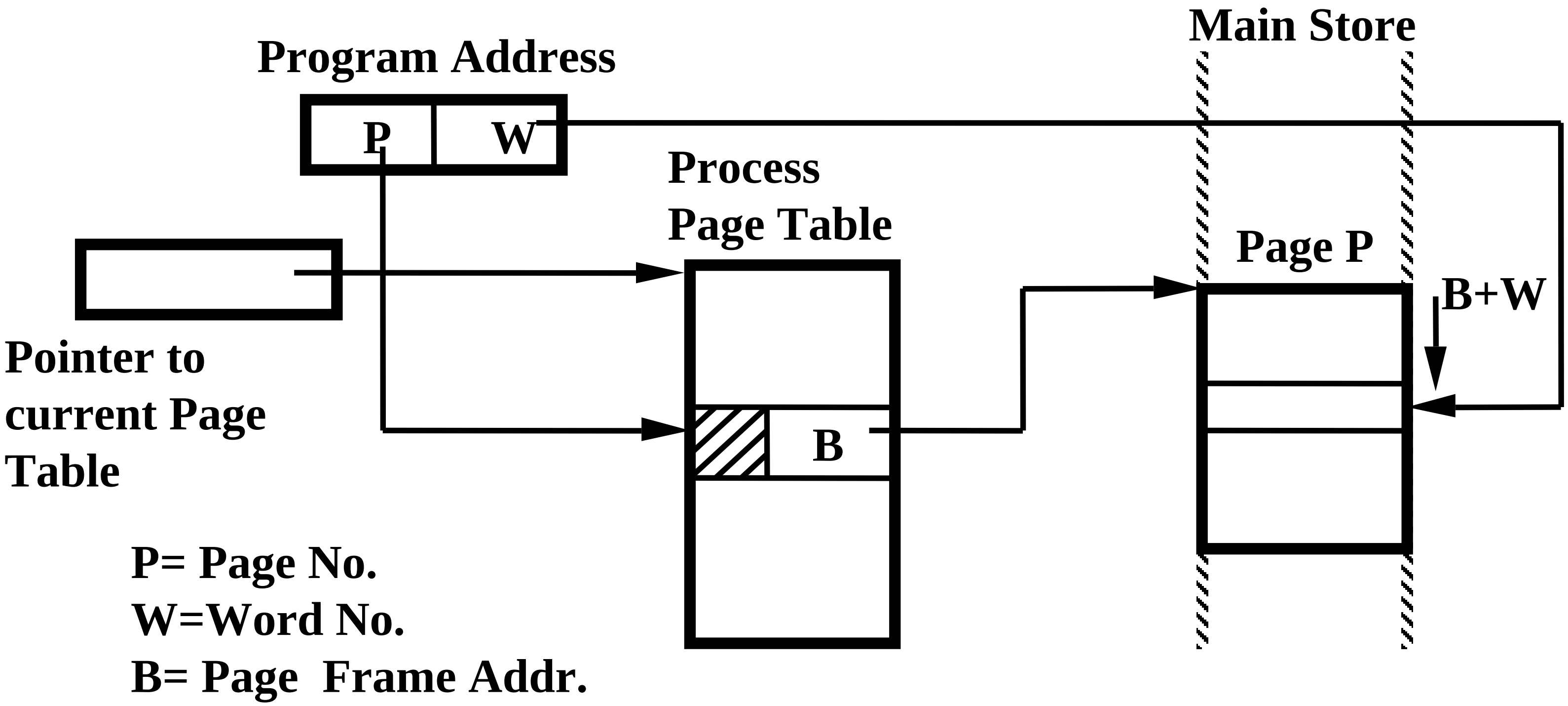
- **Running or ready process**
  - **some pages in main memory**
- **Waiting process**
  - **all pages can be on disk**
- **Paging is transparent to programmer**

Paging Mechanism

(1) Address Mapping
(2) Page Transfer

(Review introductory operating systems material for students lacking CS background)

Paging - Address Mapping



P= Page No.
W=Word No.
B= Page Frame Addr.

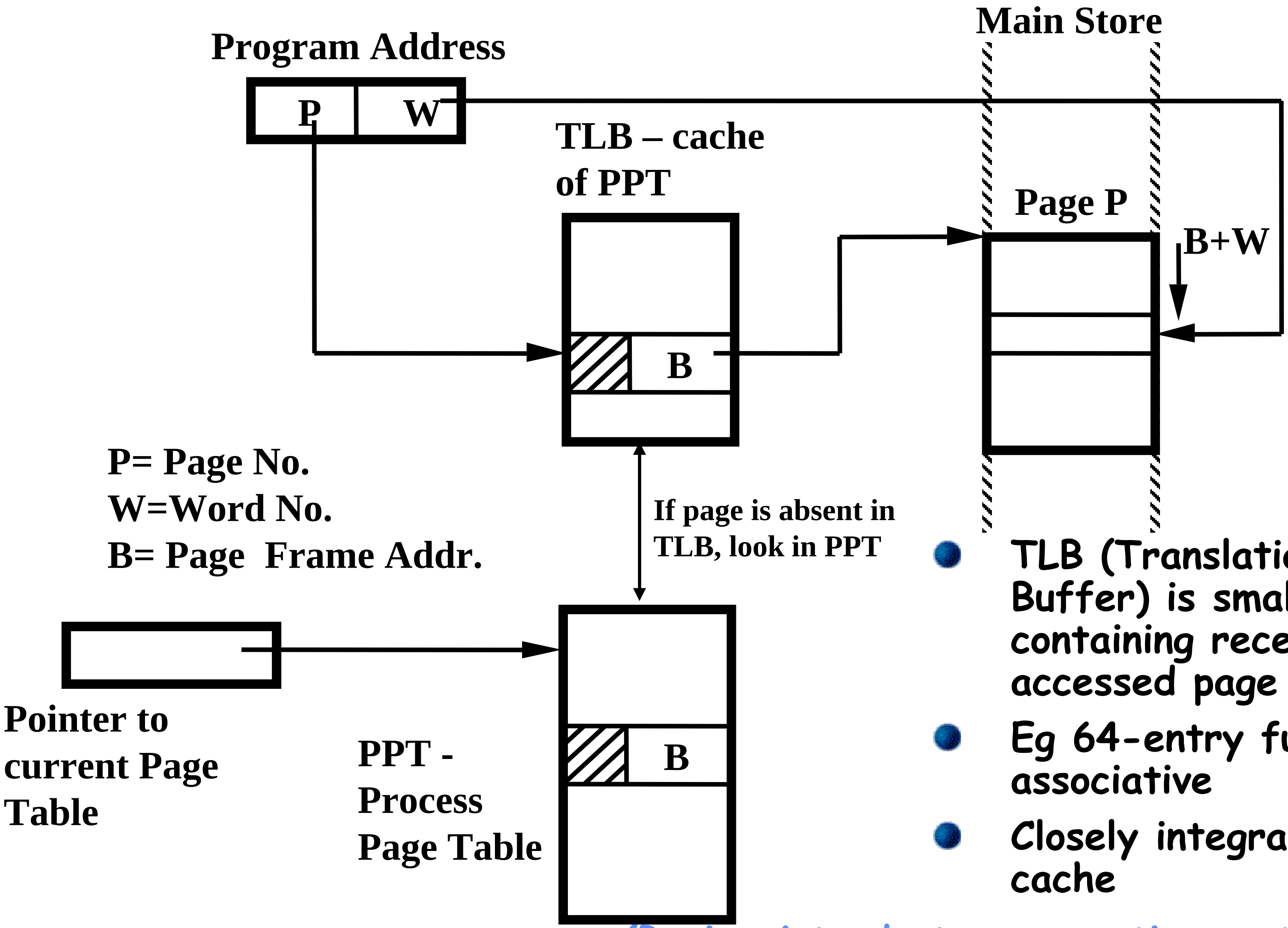
Example: Word addressed machine, W = 8 bits, page size = 256

$$Amap(P,W) := PPT[P] * 256 + W$$

Note: The Process Page Table (PPT) itself can be paged

(Review introductory operating systems material
for students lacking CS background)

Paging - Address Mapping



Program Address

P W

TLB – cache of PPT

B

Main Store

Page P

B+W

P= Page No.
W=Word No.
B= Page Frame Addr.

If page is absent in TLB, look in PPT

Pointer to current Page Table

PPT - Process Page Table

B

- TLB (Translation Lookaside Buffer) is small cache containing recently-accessed page table values
- Eg 64-entry fully-associative
- Closely integrated with L1 cache

(Review introductory operating systems material for students lacking CS background)

Paging - Page Transfer

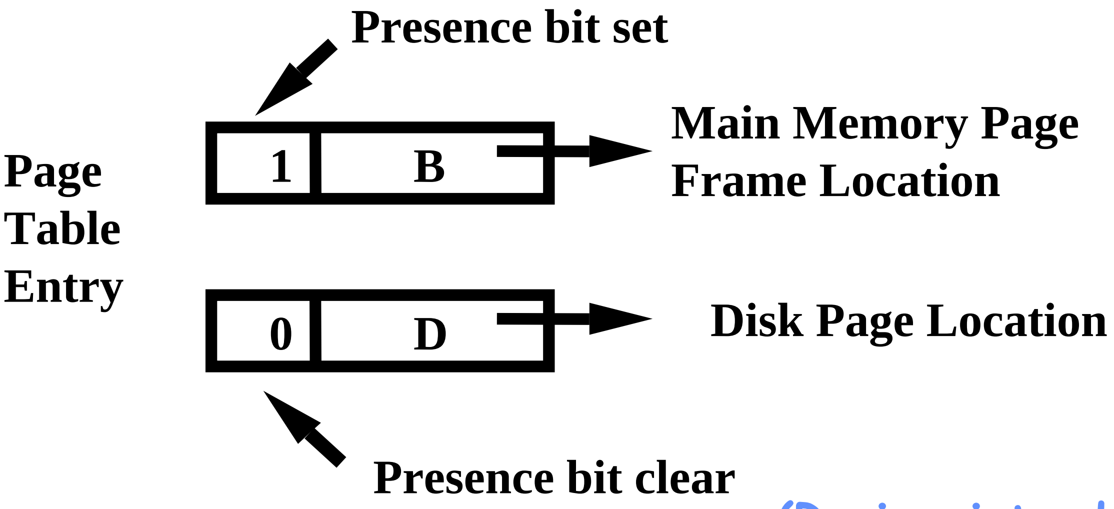
What happens when we access a page which is currently not in main memory (i.e. the page table entry is empty)?

✉ **Page Fault**

- → Suspend running process
- → Get page from disk
- → Update page table
- → Resume process (re-execute instruction)

? Can one instruction cause more than one page fault?

The location of a page on disk can be recorded in a separate table or in the page table itself using a *presence bit.*



Presence bit set

Page Table Entry

1 B

Main Memory Page Frame Location

0 D

Disk Page Location

Presence bit clear

Note: We can run another *ready* process while the page fault is being serviced.

(Review introductory operating systems material for students lacking CS background)

Synonyms and homonyms in address translation

- Homonyms *(same sound different meaning)*
  - same virtual address points to two different physical addresses in different processes
  - If you have a virtually-indexed cache, flush it between context switches - or include PID in cache tag
- Synonyms *(different sound same meaning)*
  - different virtual addresses (from the same or different processes) point to the same physical address
  - in a virtually addressed cache
    - a virtual address could be cached twice under different physical addresses
    - updates to one cached copy would not be reflected in the other cached copy
    - solution: make sure synonyms can't co-exist in the cache, *e.g., OS can forces synonyms to have the same index bits in a direct mapped cache* (sometimes called page colouring)

(a nice explanation in more detail can be found at http://www.ece.cmu.edu/~jhoe/course/ece447/handouts/L22.pdf)

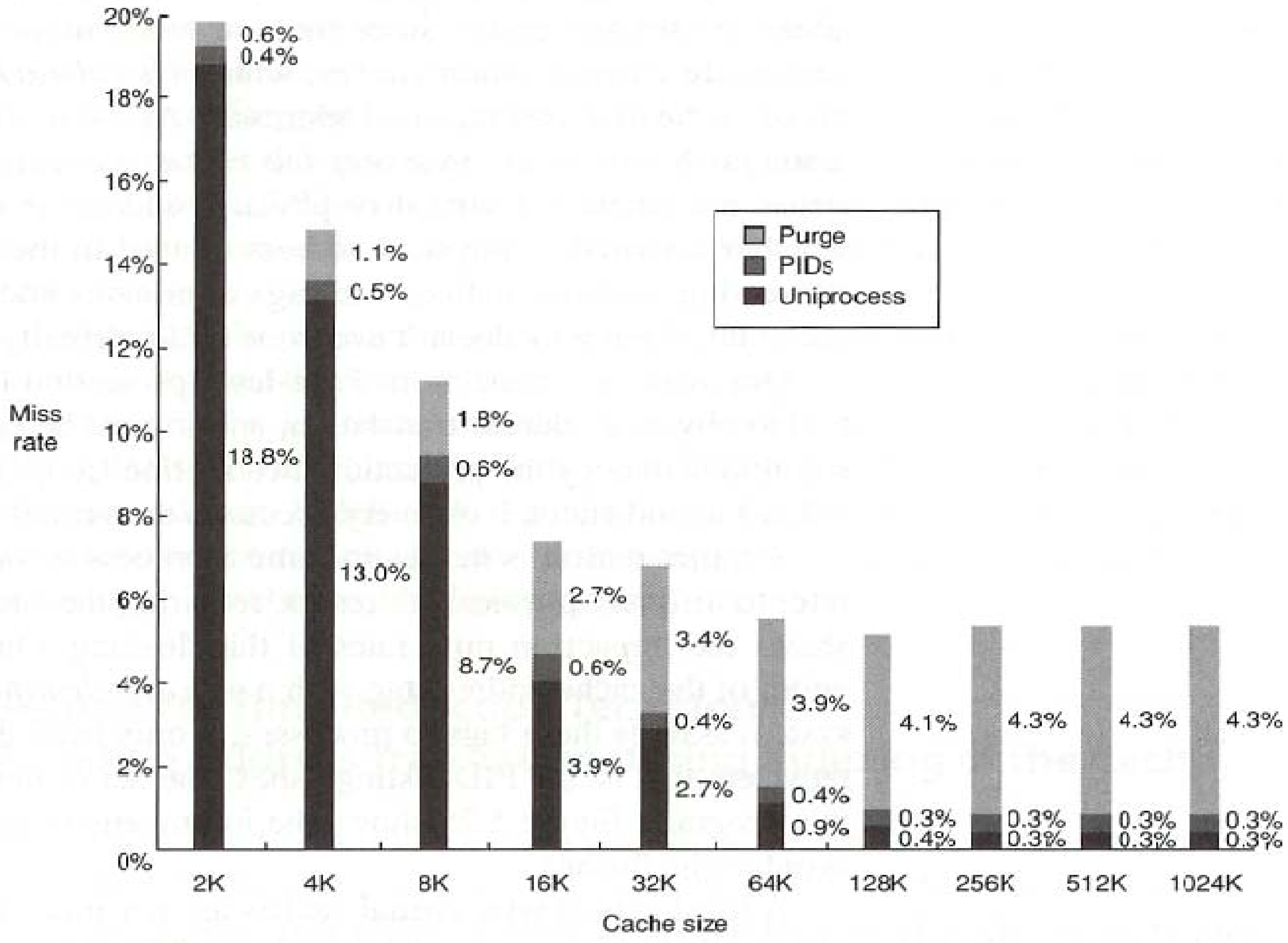
2. Fast hits by Avoiding Address Translation

- Send virtual address to cache?
  - Called *Virtually Addressed Cache* or just *Virtual Cache* vs. *Physical Cache*
  - Every time process is switched logically must flush the cache; otherwise get false hits
    - Cost is time to flush + "compulsory" misses from empty cache
  - Dealing with *aliases* (sometimes called *synonyms/homonyms*); Two different virtual addresses map to same physical address,

    Two different physical addresses mapped to by the same virtual address in different contexts
  - I/O must interact with cache, so need virtual address
- Solution to aliases
  - HW guarantees covers index field & direct mapped, they must be unique; called *page coloring*
- Solution to cache flush
  - Add *process identifier tag* that identifies process as well as address within process: can't get a hit if wrong process

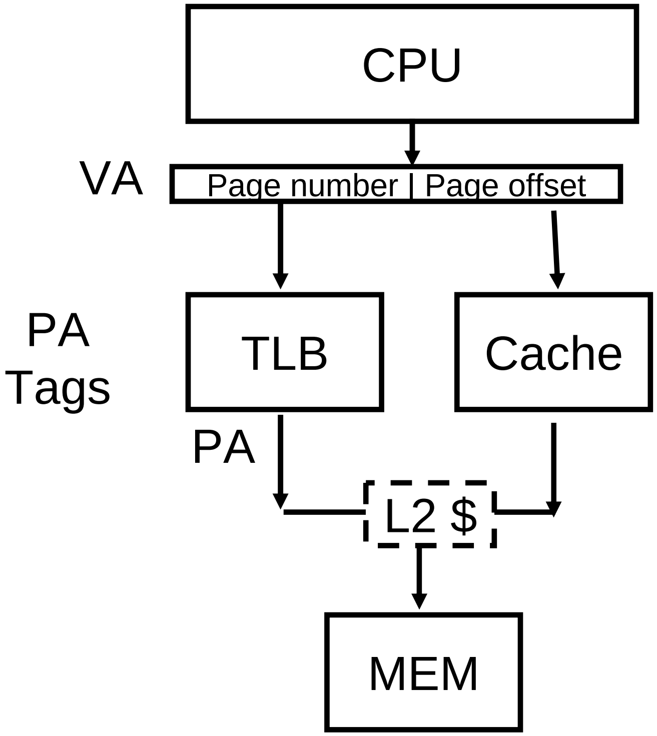
2. Fast Cache Hits by Avoiding Translation: Process ID impact

- Black is uniprocess
- Light Gray is multiprocess when flush cache
- Dark Gray is multiprocess when use Process ID tag
- Y axis: Miss Rates up to 20%
- X axis: Cache size from 2 KB to 1024 KB



2. Fast Cache Hits by Avoiding Translation: Index with Physical Portion of Address

- If index is physical part of address, can start tag access in parallel with translation so that can compare to physical tag
- Limits cache to page size: what if want bigger caches and uses same trick?
  - Higher associativity
  - Page coloring
    - A cache conflict occurs if two cache blocks that have the same tag (physical address) are mapped to two different virtual addresses
    - Make sure OS never creates a page table mapping with this property

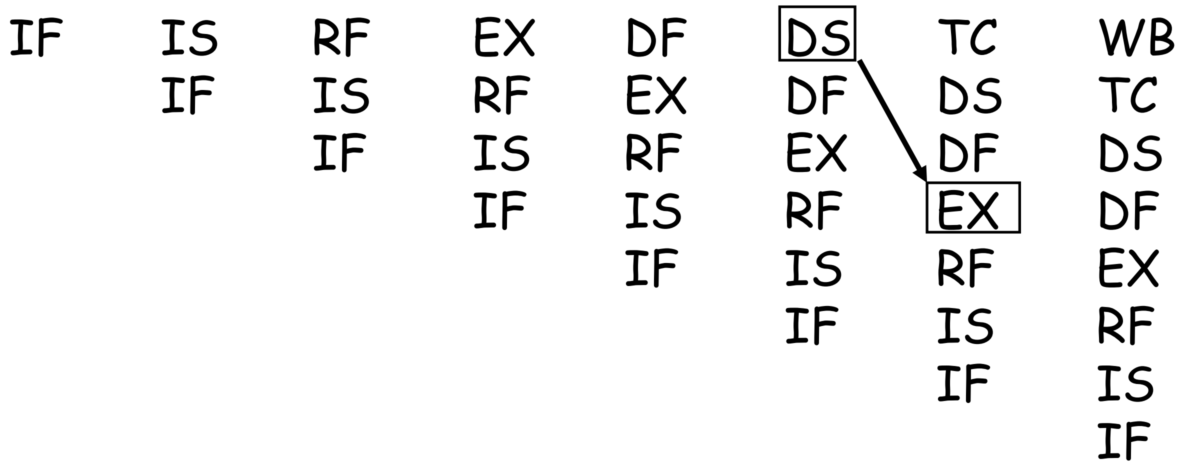


3: Fast Hits by pipelining Cache Case Study: MIPS R4000

- 8 Stage Pipeline:
  - IF–first half of fetching of instruction; PC selection happens here as well as initiation of instruction cache access.
  - IS–second half of access to instruction cache.
  - RF–instruction decode and register fetch, hazard checking and also instruction cache hit detection.
  - EX–execution, which includes effective address calculation, ALU operation, and branch target computation and condition evaluation.
  - DF–data fetch, first half of access to data cache.
  - DS–second half of access to data cache.
  - TC–tag check, determine whether the data cache access hit.
  - WB–write back for loads and register-register operations.
- What is impact on Load delay?
  - Need 2 instructions between a load and its use!

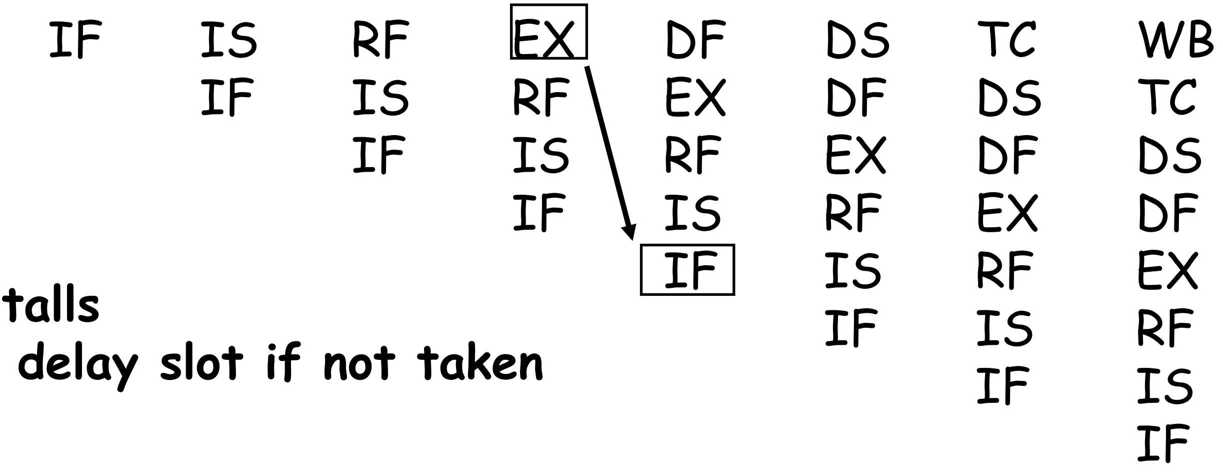
Case Study: MIPS R4000

TWO Cycle
Load Latency



IF	IS	RF	EX	DF	DS	TC	WB
	IF	IS	RF	EX	DF	DS	TC
		IF	IS	RF	EX	DF	DS
			IF	IS	RF	EX	DF
				IF	IS	RF	EX
					IF	IS	RF
						IF	IS
							IF

THREE Cycle
Branch Latency
(conditions evaluated
during EX phase)

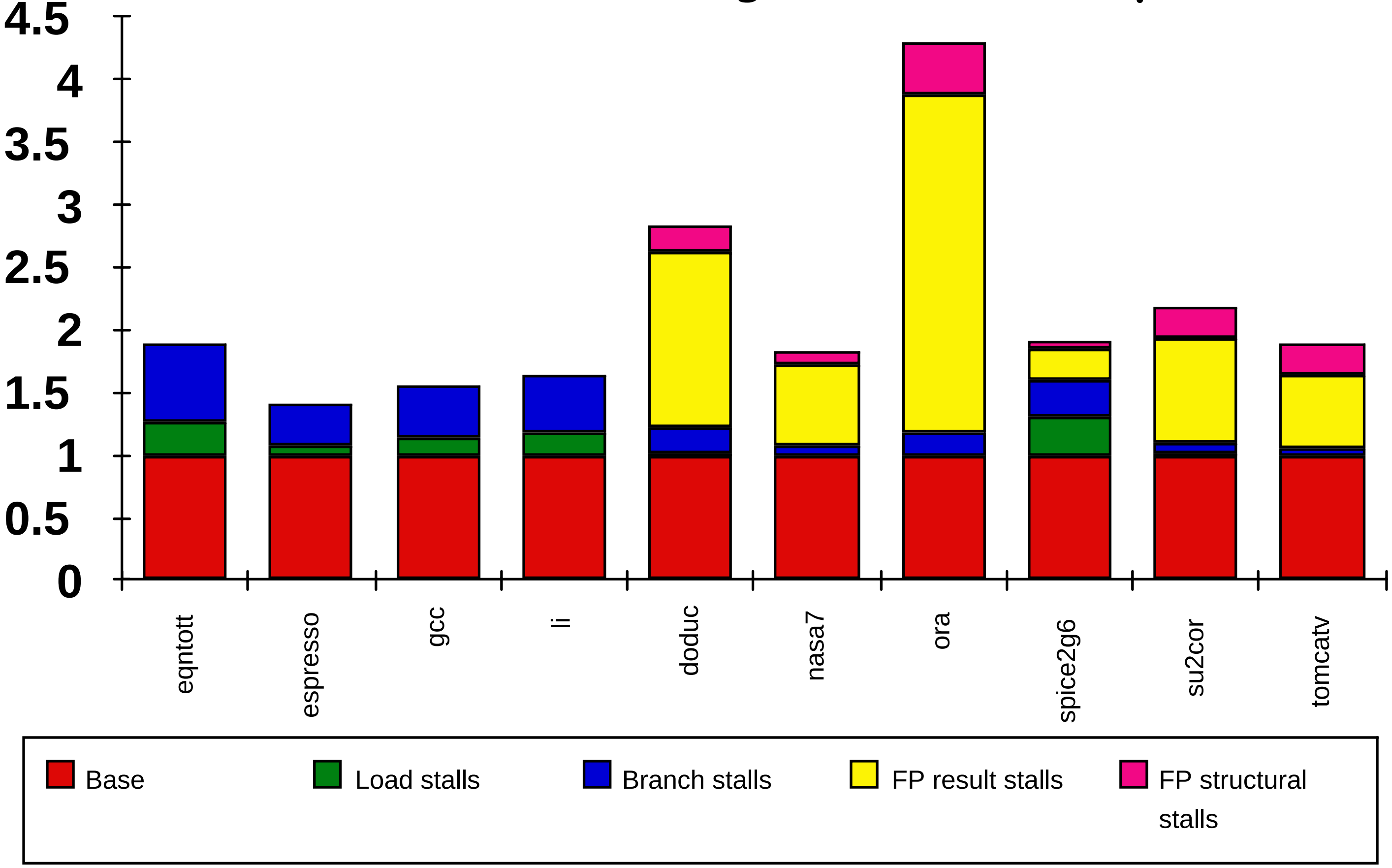


IF	IS	RF	EX	DF	DS	TC	WB
	IF	IS	RF	EX	DF	DS	TC
		IF	IS	RF	EX	DF	DS
			IF	IS	RF	EX	DF
				IF	IS	RF	EX
					IF	IS	RF
						IF	IS
							IF

Delay slot plus two stalls
Branch likely cancels delay slot if not taken

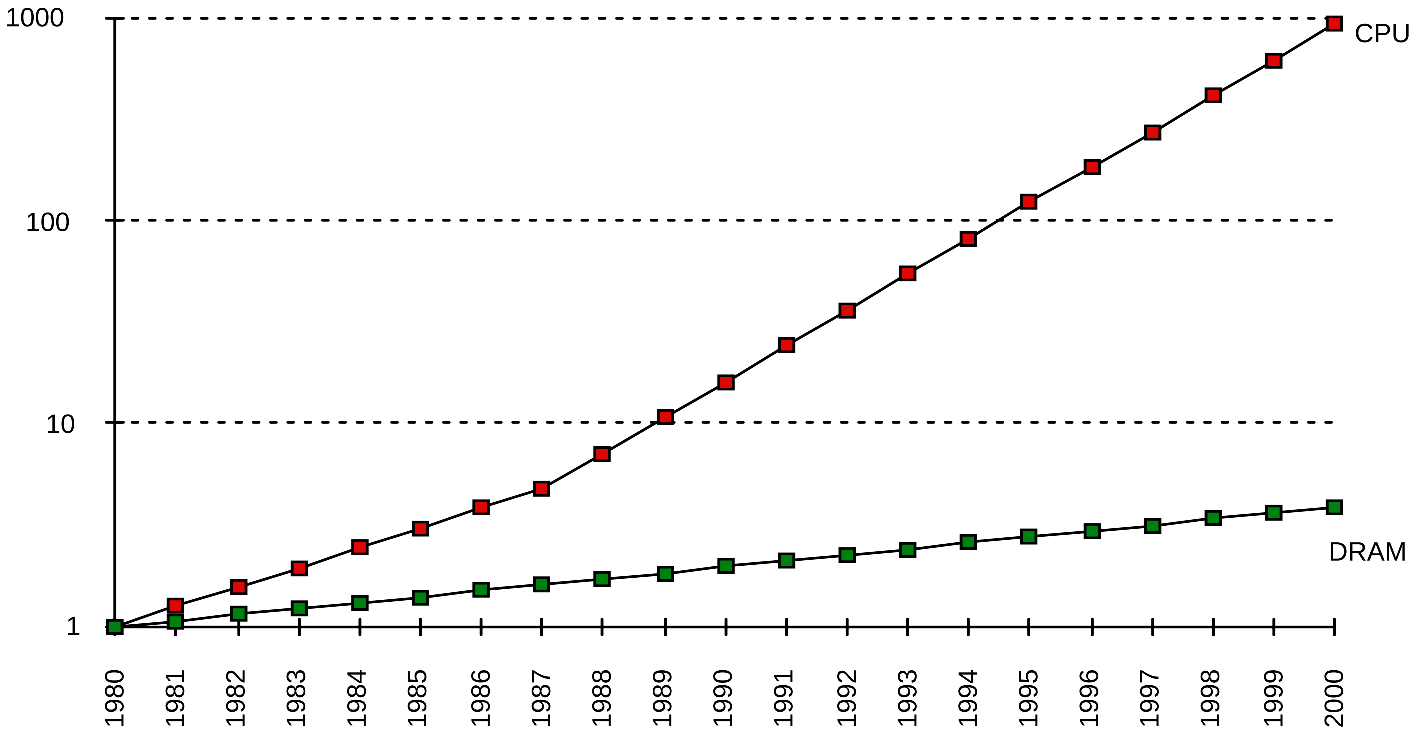
R4000 Performance

- Not ideal CPI of 1:
  - Load stalls (1 or 2 clock cycles)
  - Branch stalls (2 cycles + unfilled slots)
  - FP result stalls: RAW data hazard (latency)
  - FP structural stalls: Not enough FP hardware (parallelism)



What is the Impact of What You've Learned About Caches?

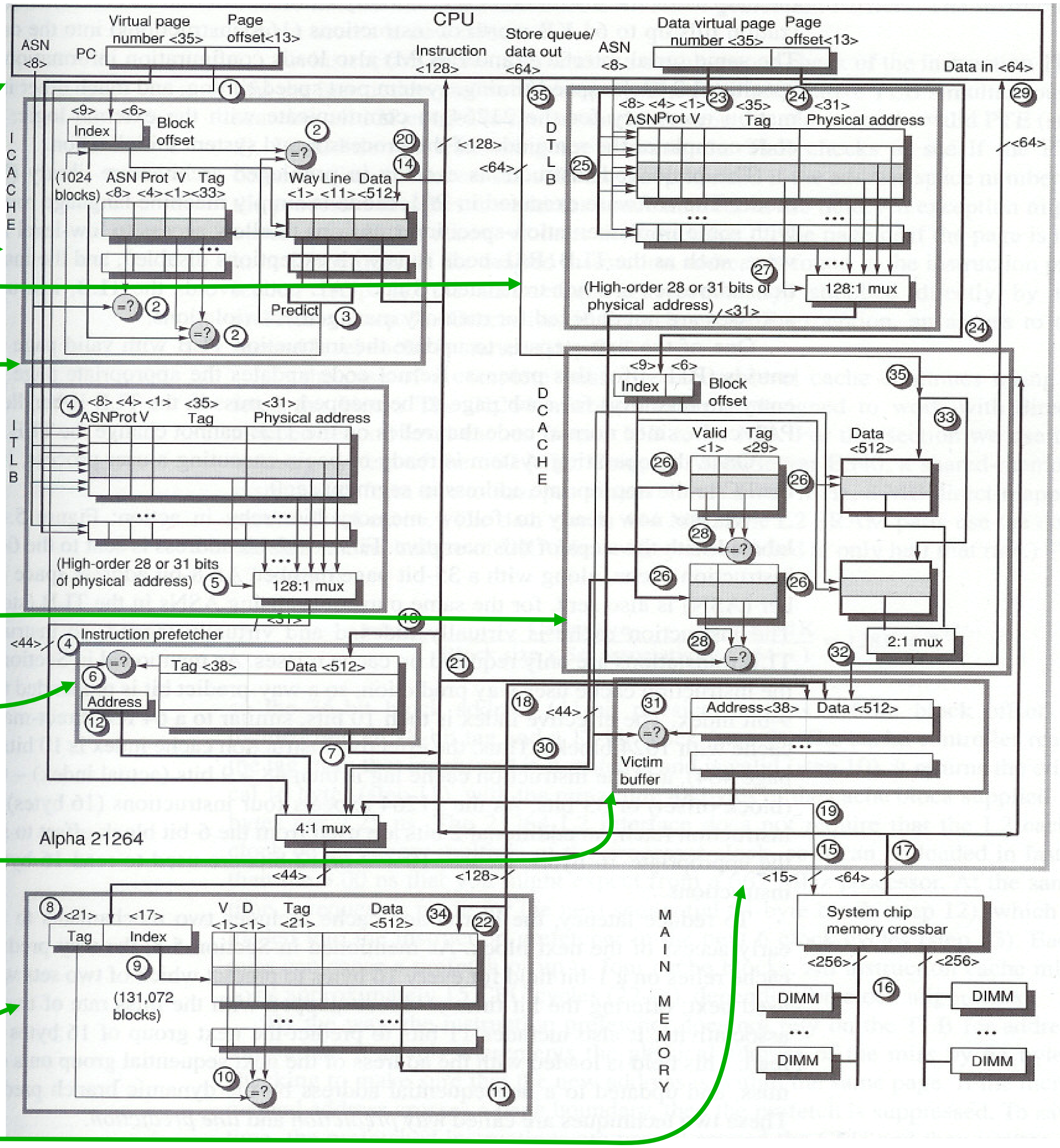
- 1960-1985: Speed = ƒ(no. operations)
- 1990
  - Pipelined Execution & Fast Clock Rate
  - Out-of-Order execution
  - Superscalar Instruction Issue



- 1998: Speed = ƒ(non-cached memory accesses)
- Superscalar, Out-of-Order machines hide L1 data cache miss (-5 clocks) but not L2 cache miss (-50 clocks)?

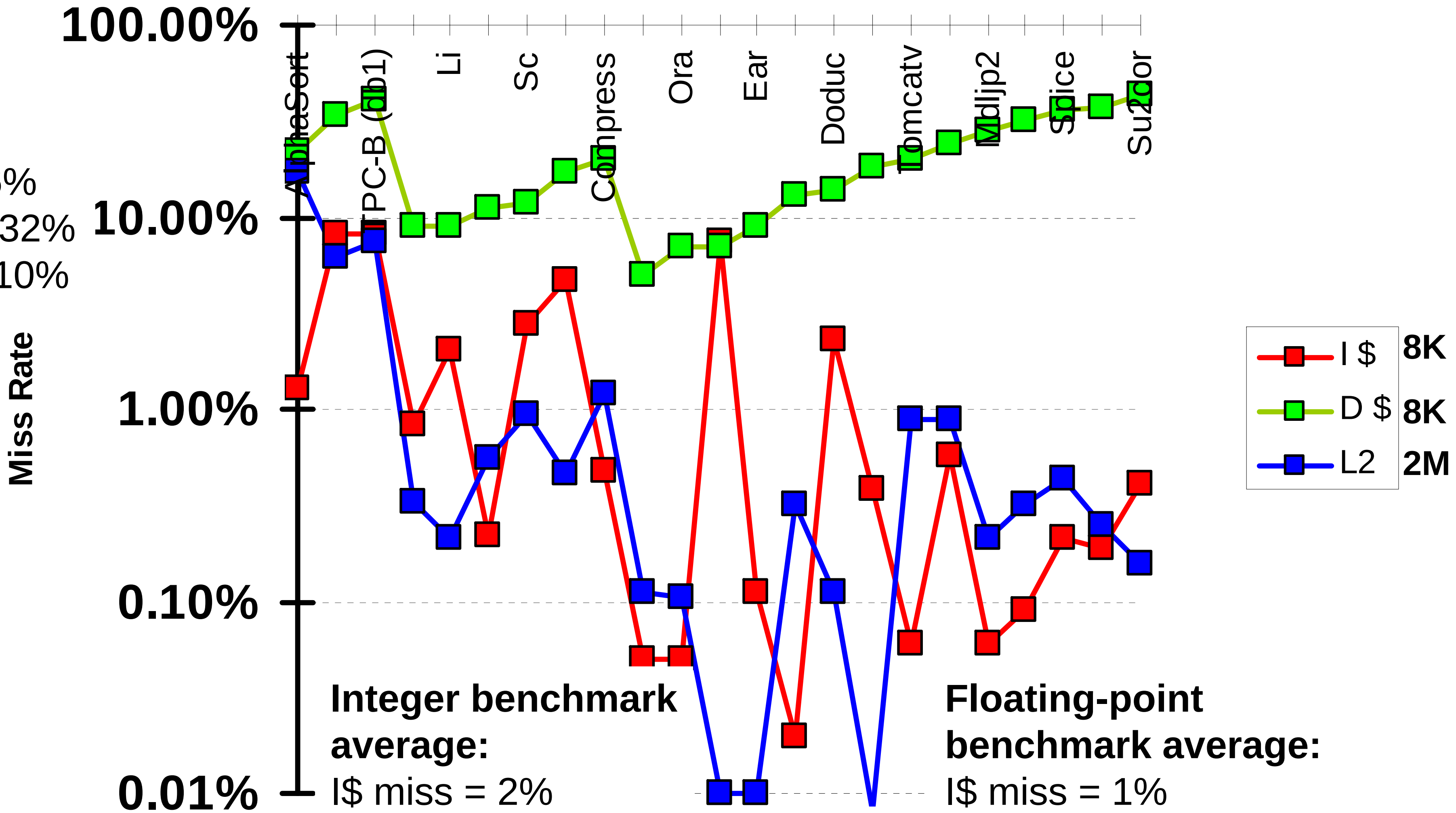
Alpha 21064

- Processor issues 48-bit virtual addresses
- Separate Instr & Data TLB & Caches
- TLBs fully associative
- TLB updates in SW ("Priv Arch Libr")
- Caches 8KB direct mapped, write thru, virtually-indexed, physically tagged
- Critical 8 bytes first
- Prefetch instr. stream buffer
- 4 entry write buffer between D$ & L2$ incorporates victim buffer: to give read priority over write
- 2 MB L2 cache, direct mapped, WB (off-chip)
- 256 bit path to main memory, 4 x 64-bit modules



Alpha Memory Performance: Miss Rates of SPEC92

Overall average:
I$ miss = 6%
D$ miss = 32%
L2 miss = 10%

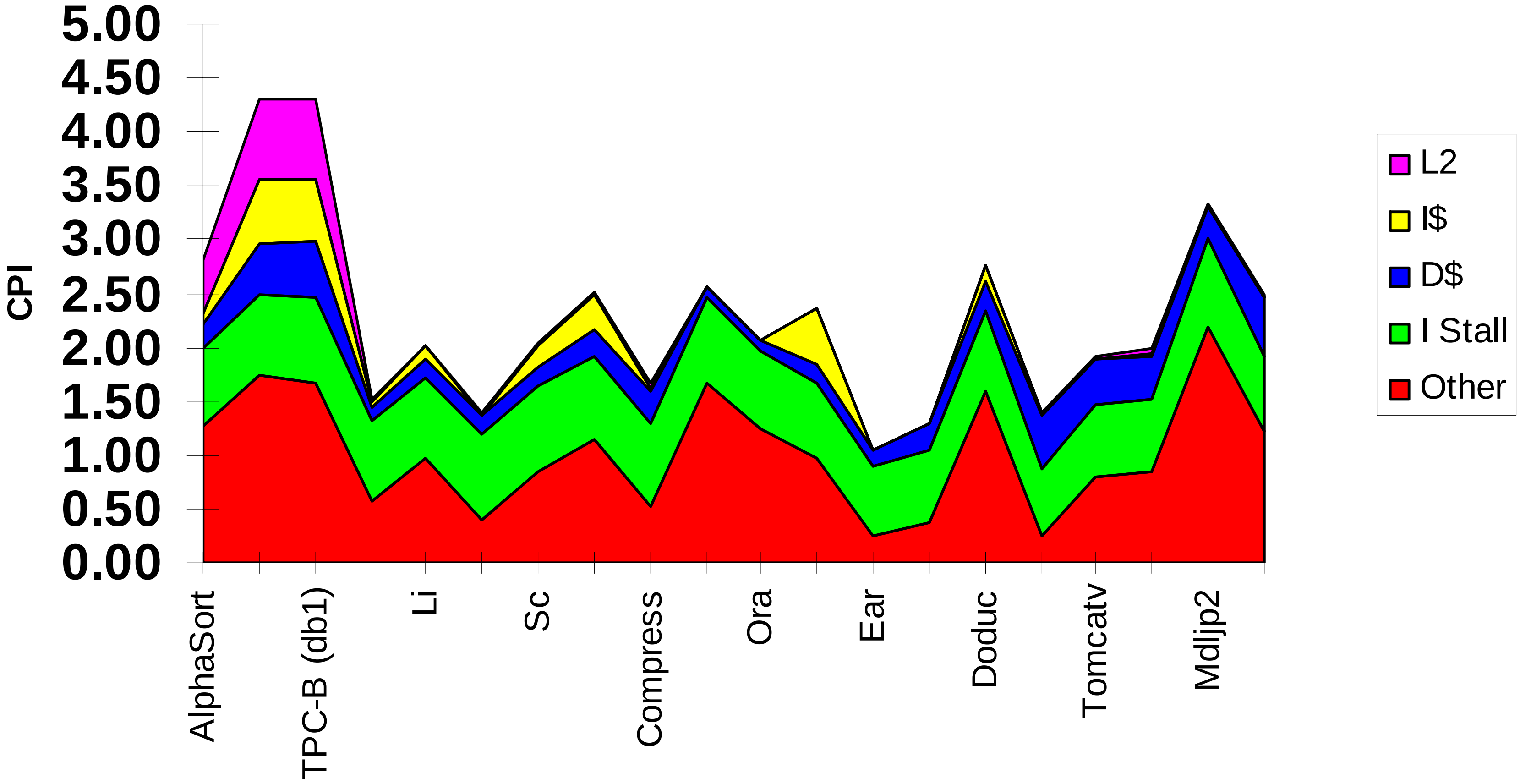


Integer benchmark average:
I$ miss = 2%
D$ miss = 13%
L2 miss = 0.6%

Floating-point benchmark average:
I$ miss = 1%
D$ miss = 21%
L2 miss = 0.3%

Alpha CPI Components

- Instruction stall: branch mispredict (green);
- Data cache (blue); Instruction cache (yellow); L2$ (pink)
  Other: compute + reg conflicts, structural conflicts



Cache Optimization Summary

	Technique	*MR*	*MP*	*HT*	*Complexity*
miss rate	Larger Block Size	+	–		0
	Higher Associativity	+		–	1
	Victim Caches	+			2
	Pseudo-Associative Caches	+			2
	HW Prefetching of Instr/Data	+			2
	Compiler Controlled Prefetching	+			3
	Compiler Reduce Misses	+			0
miss penalty	Priority to Read Misses		+		1
	Early Restart & Critical Word 1st		+		2
	Non-Blocking Caches		+		3
	Second Level Caches		+		2

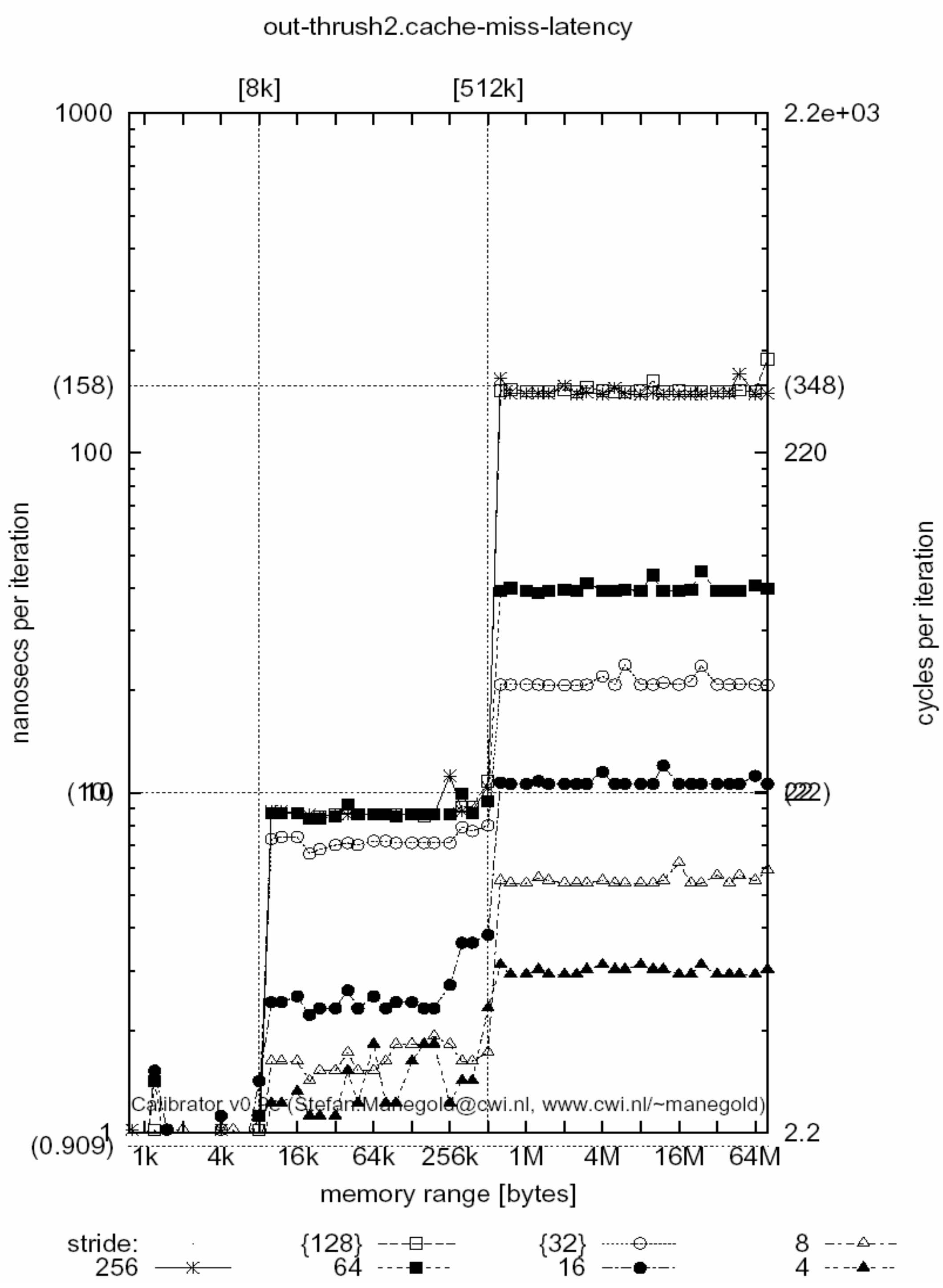
Practical exercise: explore memory hierarchy on your favourite computer

- Download Stefan Manegold's "cache and TLB calibrator":
  - http://www.cwi.nl/~manegold/Calibrator/calibrator.shtml

  (or find installed copy in ~phjk/ToyPrograms/C/ManegoldCalibrator)

- This program consists of a loop which runs over an array repeatedly
  - The size of the array is varied to evaluate cache size
  - The stride is varied to explore block size

Memory hierarchy of a 2.2GHz Intel Pentium 4 Xeon



- Memory access latency is close to 1ns when loop reuses array smaller than 8KB level-1 cache
- While array is smaller than 512KB, access time is 2-8ns, depending on stride
- When array exceeds 512KB, accesses miss both level-1 and level-2 caches
- Worst case (large stride) suffers 158ns access latency
- Q:
  - How many instructions could be executed in 158ns?
  - what is the level-1 cache block size?
  - What is the level-2 cache block size?

Instructions for running the Manegold calibrator

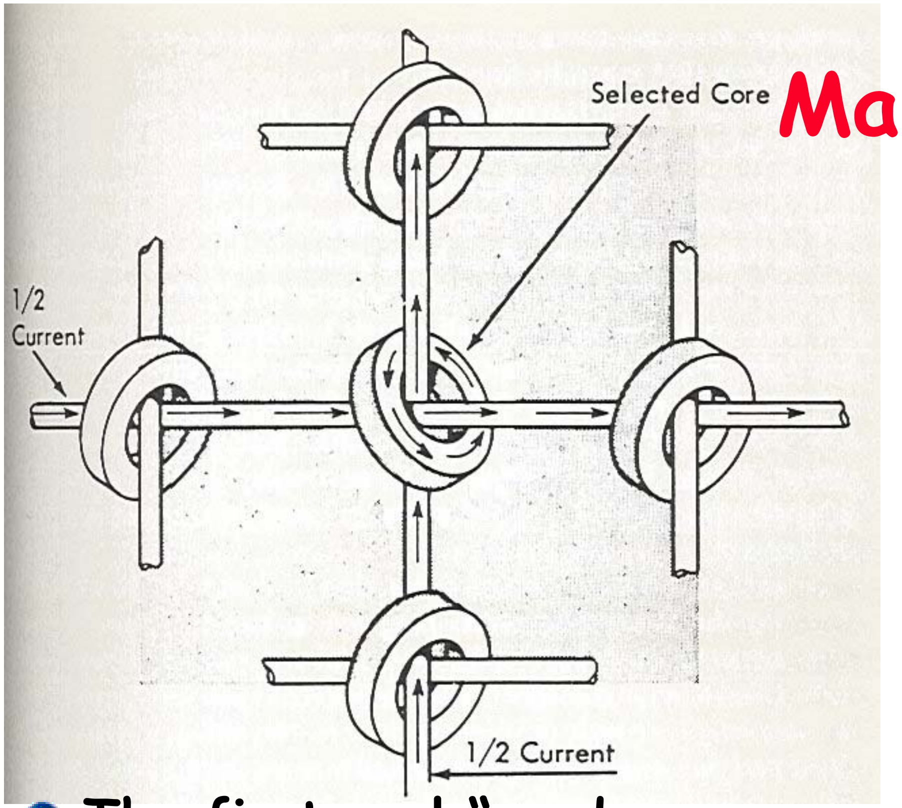
- Get a copy:
  - cp /homes/phjk/ToyPrograms/C/ManegoldCalibrator/calibrator.c ./
- Compile it:
  - gcc –O3 –o calibrator calibrator.s
- Find out CPU MHz
  - cat /proc/cpuinfo
- Run it; ./calibrator <CPUMHz> <size> <filename>
- Eg on media03:
  - ./calibrator 3000 64M media03
  - Output is delivered to a set of files “media03.*”
- Plot postscript graphs using generated gnuplot scripts:
  - gnuplot media03.cache-miss-latency.gp
  - gnuplot media03.cache-replace-time.gp
  - gnuplot media03.TLB-miss-latency.gp
- View the generated postscript files:
  - gv media03.cache-miss-latency.ps &

Extra material for interest

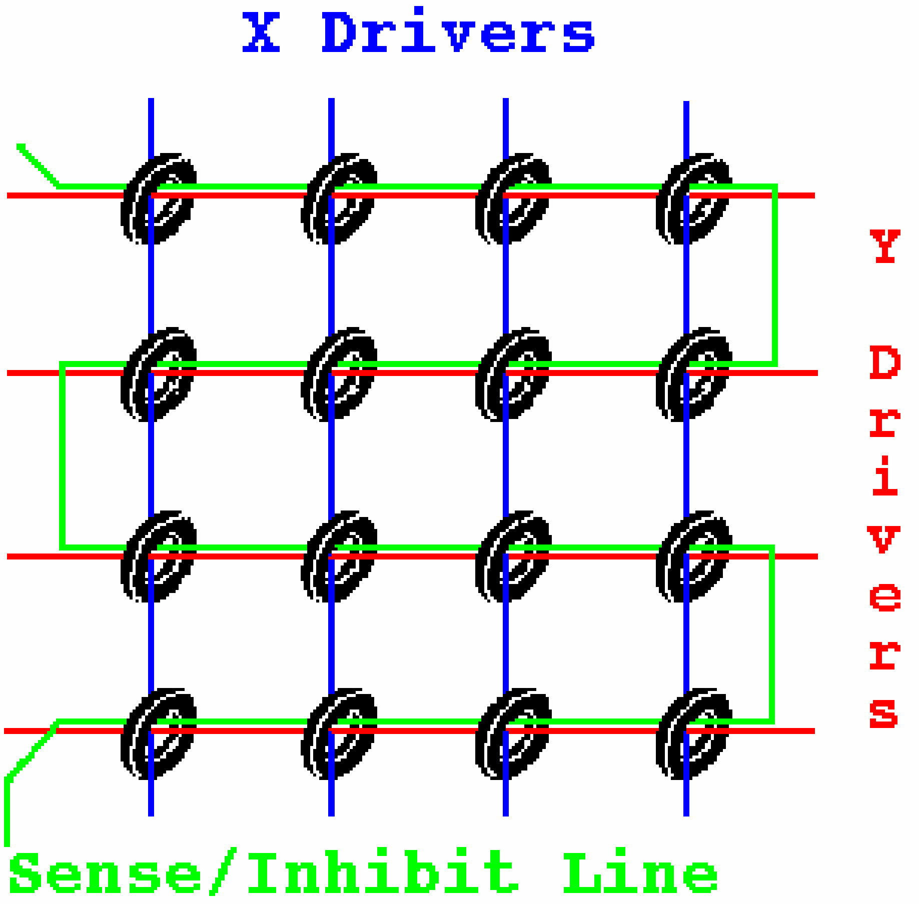
Main Memory Background

- Performance of Main Memory:
  - Latency: Cache Miss Penalty
    - *Access Time*: time between request and word arrives
    - *Cycle Time*: time between requests
  - Bandwidth: I/O & Large Block Miss Penalty (L2)
- Main Memory is *DRAM*: Dynamic Random Access Memory
  - Dynamic since needs to be refreshed periodically (8 ms, 1% time)
  - Addresses divided into 2 halves (Memory as a 2D matrix):
    - *RAS* or *Row Access Strobe*
    - *CAS* or *Column Access Strobe*
- Cache uses *SRAM*: Static Random Access Memory
  - No refresh (6 transistors/bit vs. 1 transistor
    Size: DRAM/SRAM - *4-8*,
    Cost/Cycle time: SRAM/DRAM - *8-16*

Main Memory Deep Background

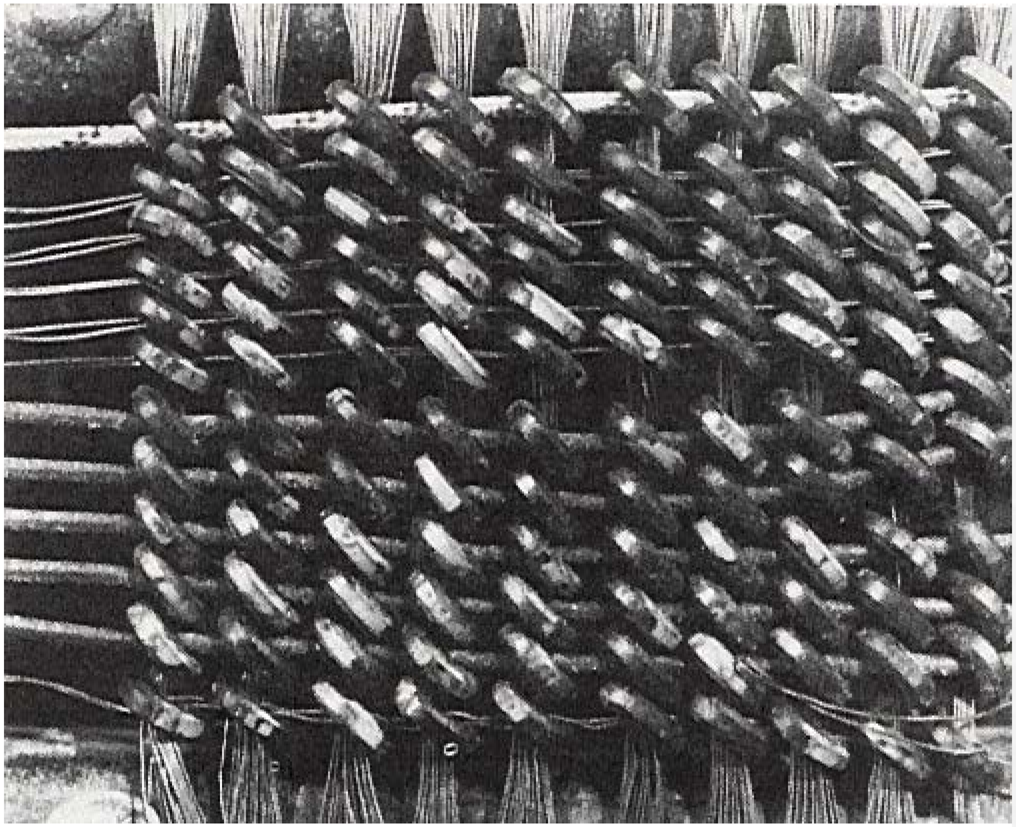


- The first real "random-access memory" technology was based on magnetic "cores" – tiny ferrite rings threaded with copper wires
- That's why people talk about "Out-of-Core", "In-Core," "Core Dump"
- Non-volatile, magnetic
- Lost out when 4 Kbit DRAM became available
- Access time 750 ns, cycle time 1500-3000 ns



Pulse on sense line if any core flips its magnetisation state

http://www.faqs.org/docs/electric/Digital/DIGI_15.html
http://www.psych.usyd.edu.au/pdp-11/core.html



- The first magnetic core memory, from the IBM 405 Alphabetical Accounting Machine. The photo shows the single drive lines through the cores in the long direction and fifty turns in the short direction. The cores are 150 mil inside diameter, 240 mil outside, 45 mil high. This experimental system was tested successfully in April 1952.

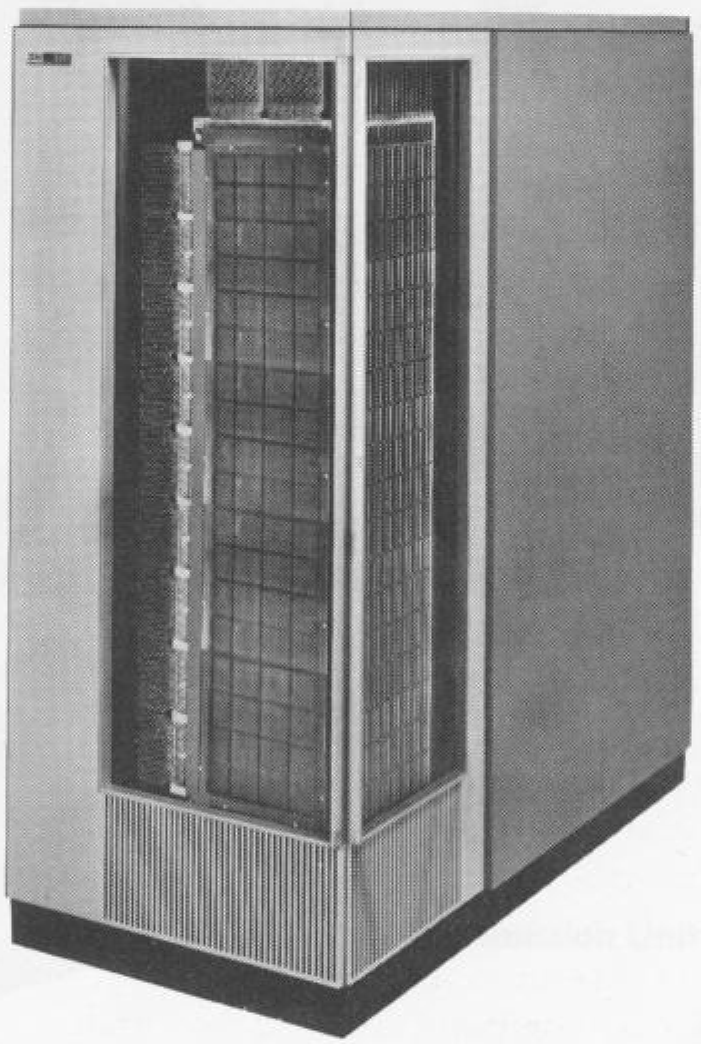


Figure 10. IBM 2361 Core Storage

- 524,000 36-bit words and a total cycle time of eight microseconds in each memory (1964 – for the IBM7094)

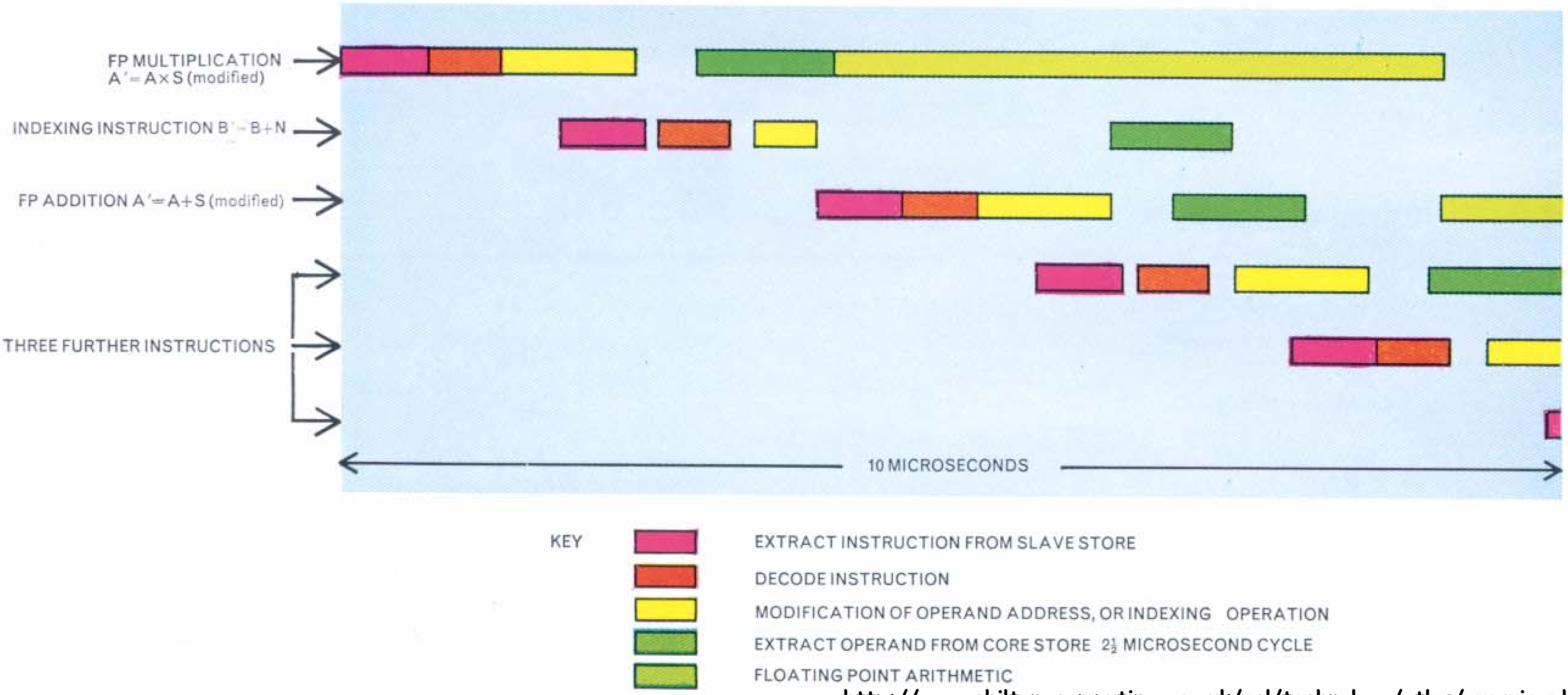
Sources: http://www-03.ibm.com/ibm/history/exhibits/space/space_2361.html
http://www.columbia.edu/acis/history/core.html

http://www.chilton-computing.org.uk/acl/technology/atlas/overview.htm

Atlas

- First was at University of Manchester
- University of London had the second one
- Commissioned May 1964
- Shut down Sept 1972

Pipelined instruction processing in Atlas


FP MULTIPLICATION
A' = A×S (modified)
INDEXING INSTRUCTION B' = B+N
FP ADDITION A' = A+S (modified)
THREE FURTHER INSTRUCTIONS
10 MICROSECONDS
KEY
EXTRACT INSTRUCTION FROM SLAVE STORE
DECODE INSTRUCTION
MODIFICATION OF OPERAND ADDRESS, OR INDEXING OPERATION
EXTRACT OPERAND FROM CORE STORE 2½ MICROSECOND CYCLE
FLOATING POINT ARITHMETIC

http://www.chilton-computing.org.uk/acl/technology/atlas/overview.htm

- **Atlas is most famous for pioneering virtual memory**
- **Also**
  - **Pipelined execution**
  - **Cache memory ("slave store") – 32 words**
  - **Floating point arithmetic hardware**

DRAM cell design

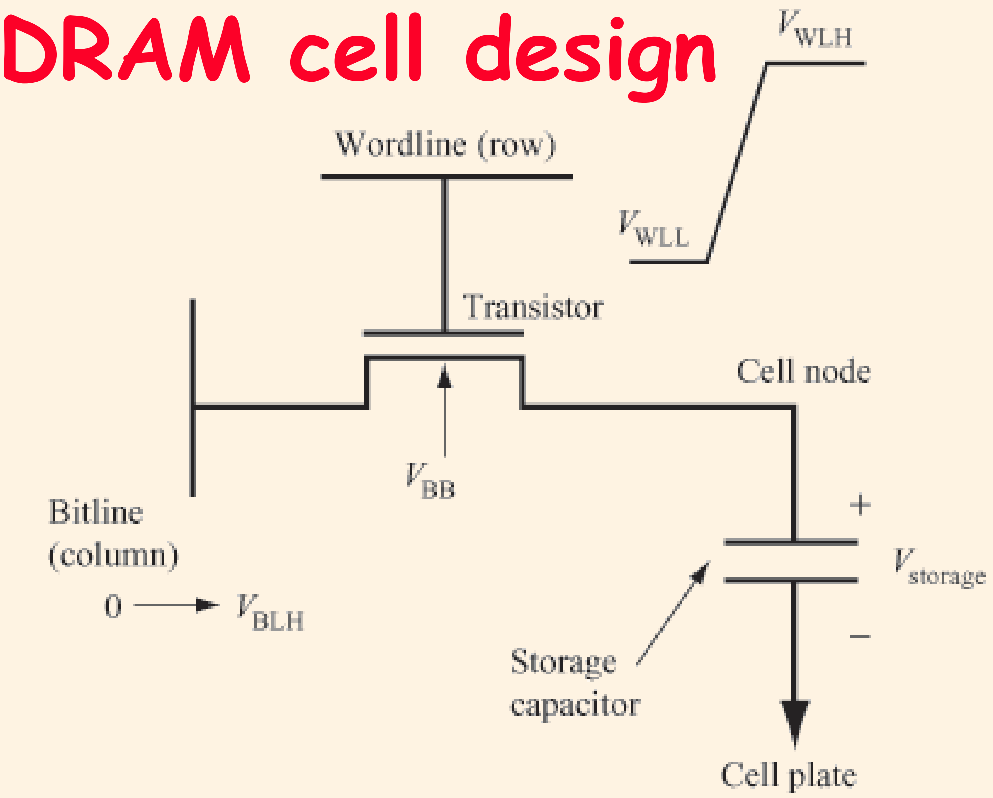


Figure 1

Schematic of a one-transistor DRAM cell [1]. The array device (transistor) is addressed by switching the wordline voltage from V_{WLL} (wordline-low) to V_{WLH} (wordline-high), enabling the bitline and the capacitor to exchange charge. In this example, a data state of either a "0" (0 V) or a "1" (V_{BLH}) is written from the bitline to the storage capacitor. V_{BB} is the electrical bias applied to the p-well.

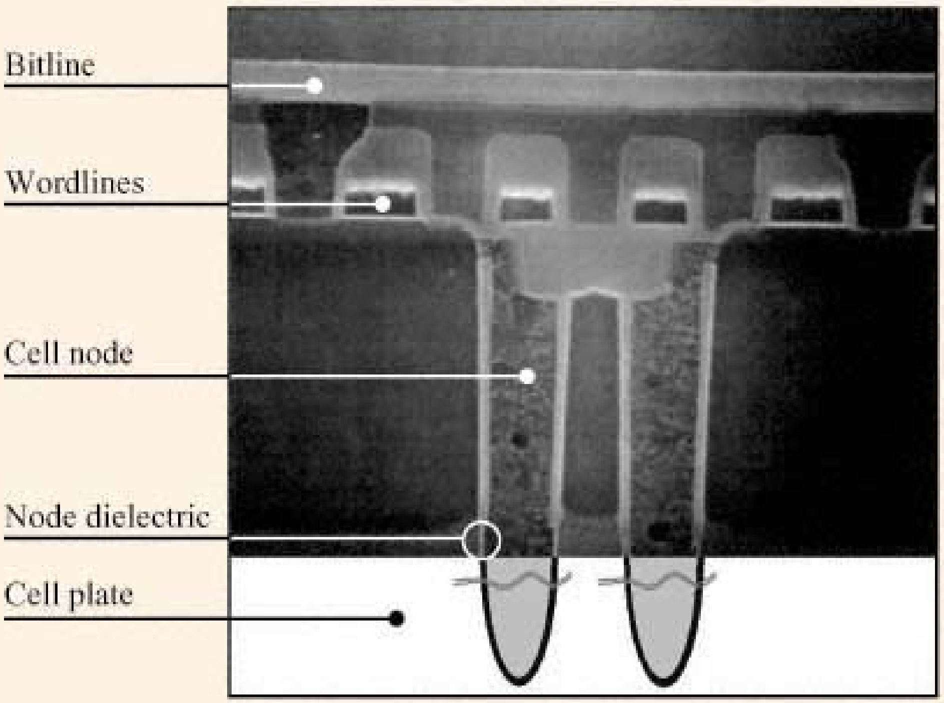


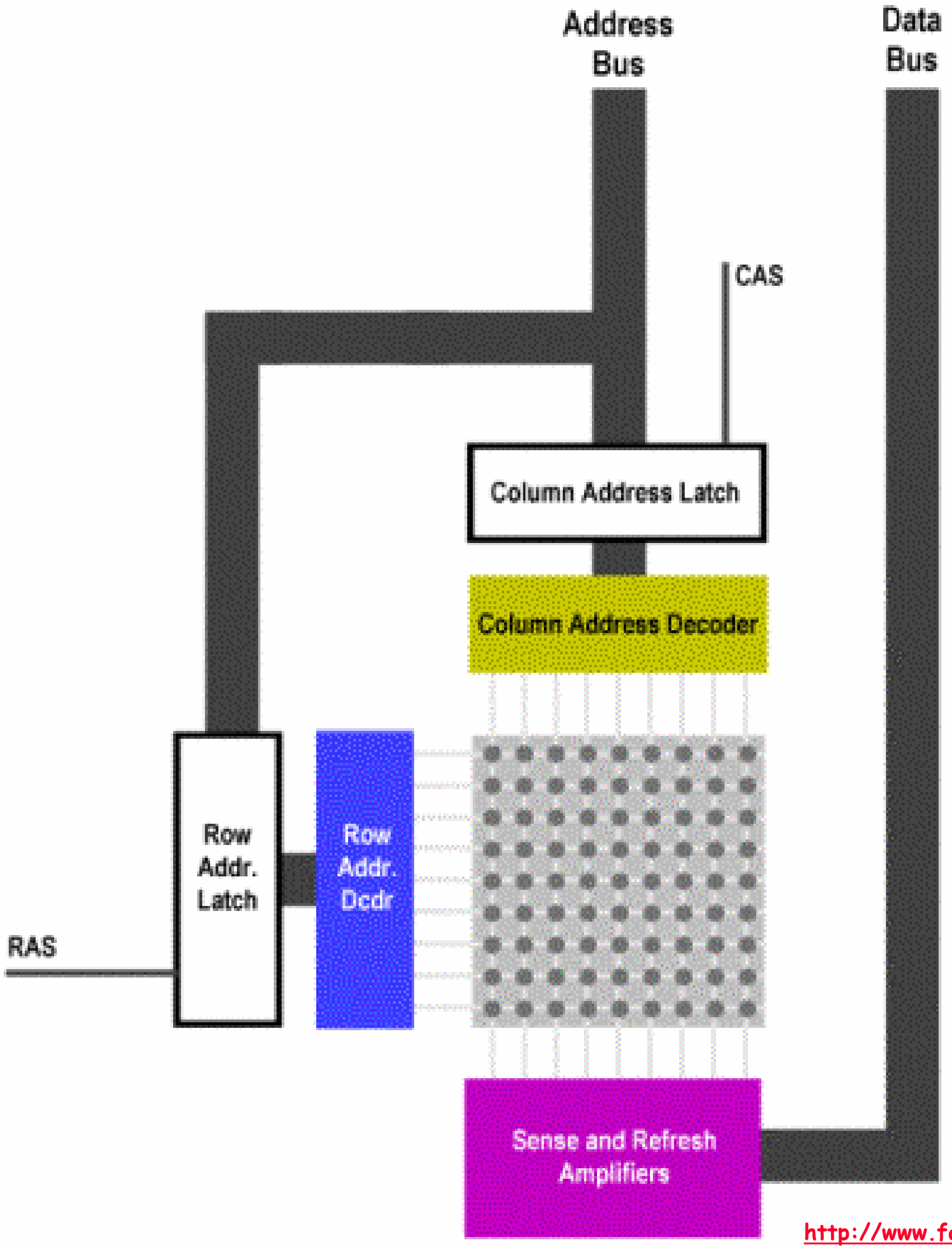
Figure 4

SEM photomicrograph of 0.25-μm trench DRAM cell suitable for scaling to 0.15μm and below. Reprinted with permission from [17]; © 1995 IEEE.

- Single transistor
- Capacitor stores charge
- Decays with time
- Destructive read-out

http://www.research.ibm.com/journal/rd/462/mandelman.html

DRAM array design



- Square array of cells
- Address split into Row address and Column Address bits
- Row address selects row of cells to be activated
- Cells discharge
- Cell state latched by per-column sense amplifiers
- Column address selects data for output
- Data must be written back to selected row

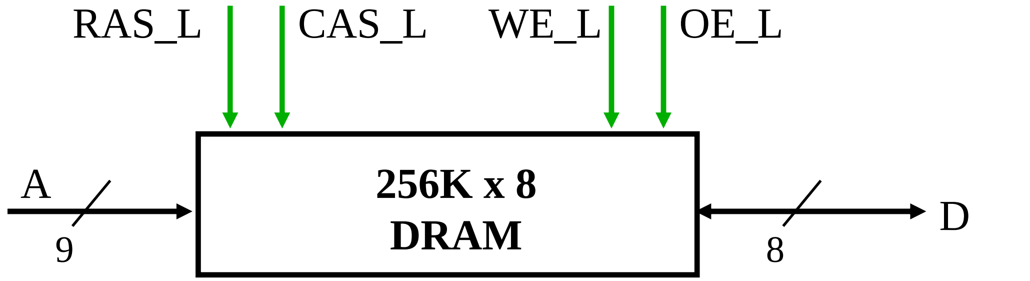
http://www.faculty.iu-bremen.de/birk/lectures/PC101-2003/08dram/Principles/DRAM02.htm

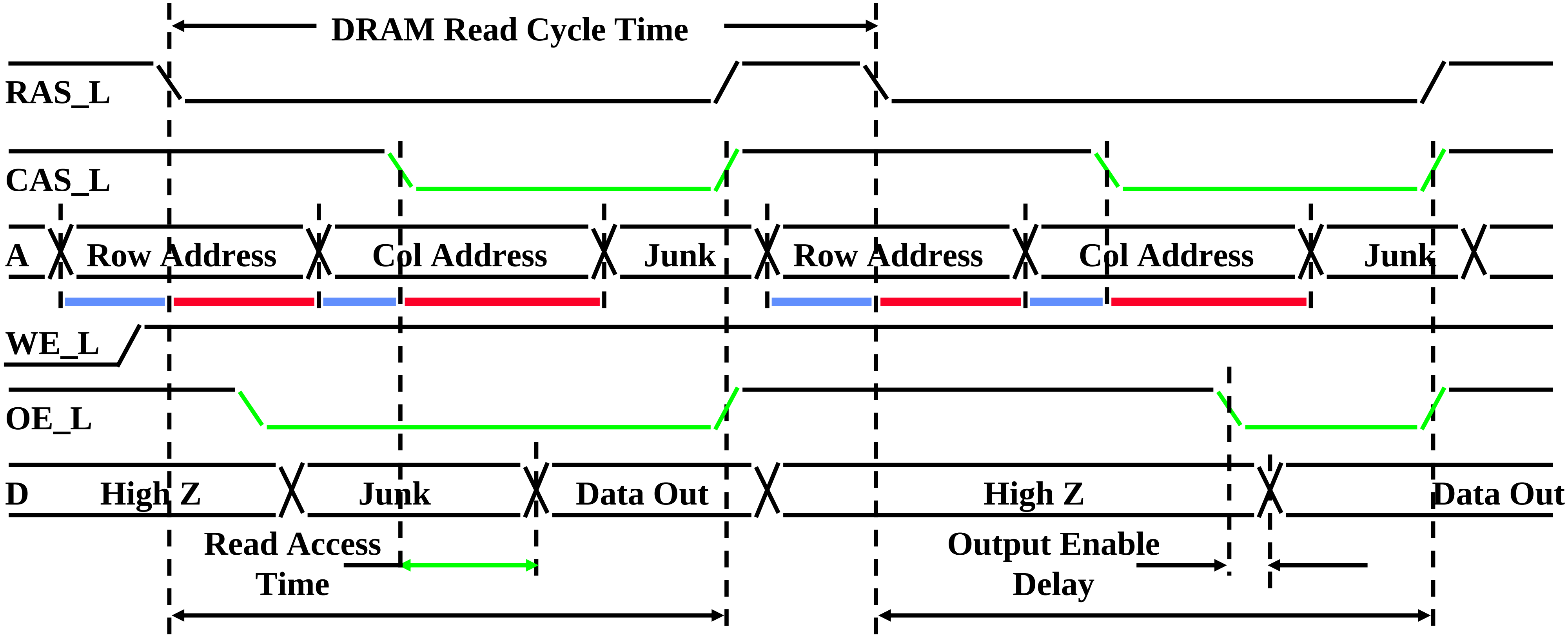
4 Key DRAM Timing Parameters

- t_{RAC}: minimum time from RAS line falling to the valid data output.
  - Quoted as the speed of a DRAM when buy
  - A typical 4Mb DRAM t_{RAC} = 60 ns
  - Speed of DRAM since on purchase sheet?
- t_{RC}: minimum time from the start of one row access to the start of the next.
  - t_{RC} = 110 ns for a 4Mbit DRAM with a t_{RAC} of 60 ns
- t_{CAC}: minimum time from CAS line falling to valid data output.
  - 15 ns for a 4Mbit DRAM with a t_{RAC} of 60 ns
- t_{PC}: minimum time from the start of one column access to the start of the next.
  - 35 ns for a 4Mbit DRAM with a t_{RAC} of 60 ns

DRAM Read Timing

- Every DRAM access begins at:
  - The assertion of the RAS_L
  - 2 ways to read: early or late v. CAS





Early Read Cycle: OE_L asserted before CAS_L

Late Read Cycle: OE_L asserted after CAS_L

DRAM Performance

- A 60 ns (t_{RAC}) DRAM can
  - perform a row access only every 110 ns (t_{RC})
  - perform column access (t_{CAC}) in 15 ns, but time between column accesses is at least 35 ns (t_{PC}).
    - In practice, external address delays and turning around buses make it 40 to 50 ns
- These times do not include the time to drive the addresses off the microprocessor nor the memory controller overhead!

DRAM History

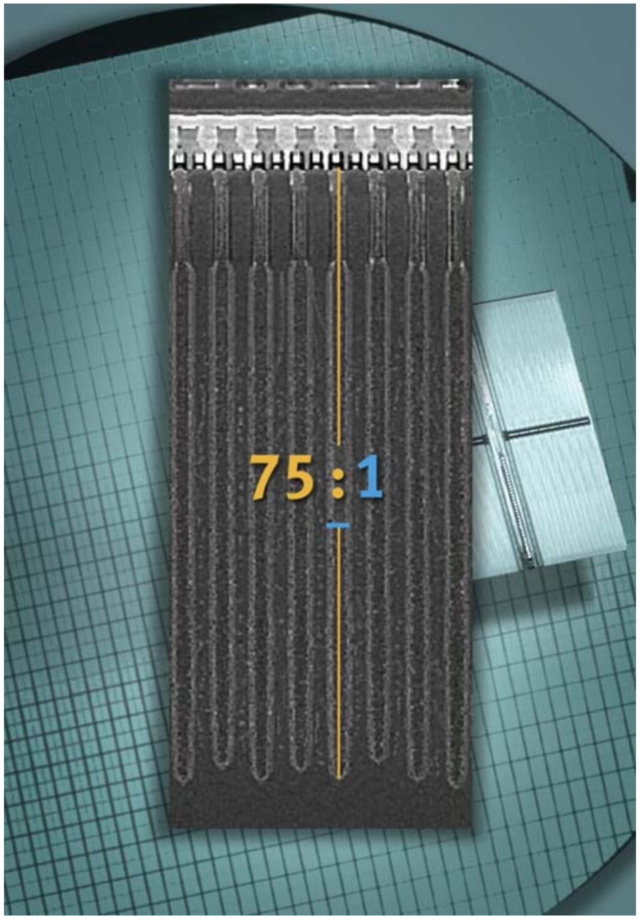
- DRAMs: capacity +60%/yr, cost –30%/yr
  - 2.5X cells/area, 1.5X die size in -3 years
- 2007 DRAM fab line costs $4.6B (2004 prices)
  - DRAM only: density, leakage v. speed
- Rely on increasing no. of computers & memory per computer (60% market)
  - SIMM or DIMM is replaceable unit
    => computers use any generation DRAM
- Commodity, second source industry
  => high volume, low profit, conservative
  - Little organization innovation in 20 years
- Order of importance: 1) Cost/bit 2) Capacity
  - First RAMBUS: 10X BW, +30% cost => little impact

"Elpida to Build $4.6B DRAM Fab in Japan" (*Electronic News*, 6/9/2004)
http://www.reed-electronics.com/electronicnews/article/CA424812.html

DRAM Today: 1 Gbit DRAM and more

	Infineon (Dresden)
Organisation	x4,x8,x16
Clock	133-200 MHz
Data Pins	68
Die Size	160 mm^2
Metal Layers	3
Technology	110nm

Video: http://registration.infineon.com/registration/video/video.asp



Future (in lab 2005):

- 0.7 micron
- Hi-k dielectric (Al_2O_3)
- 75:1 trench capacitor aspect ratio

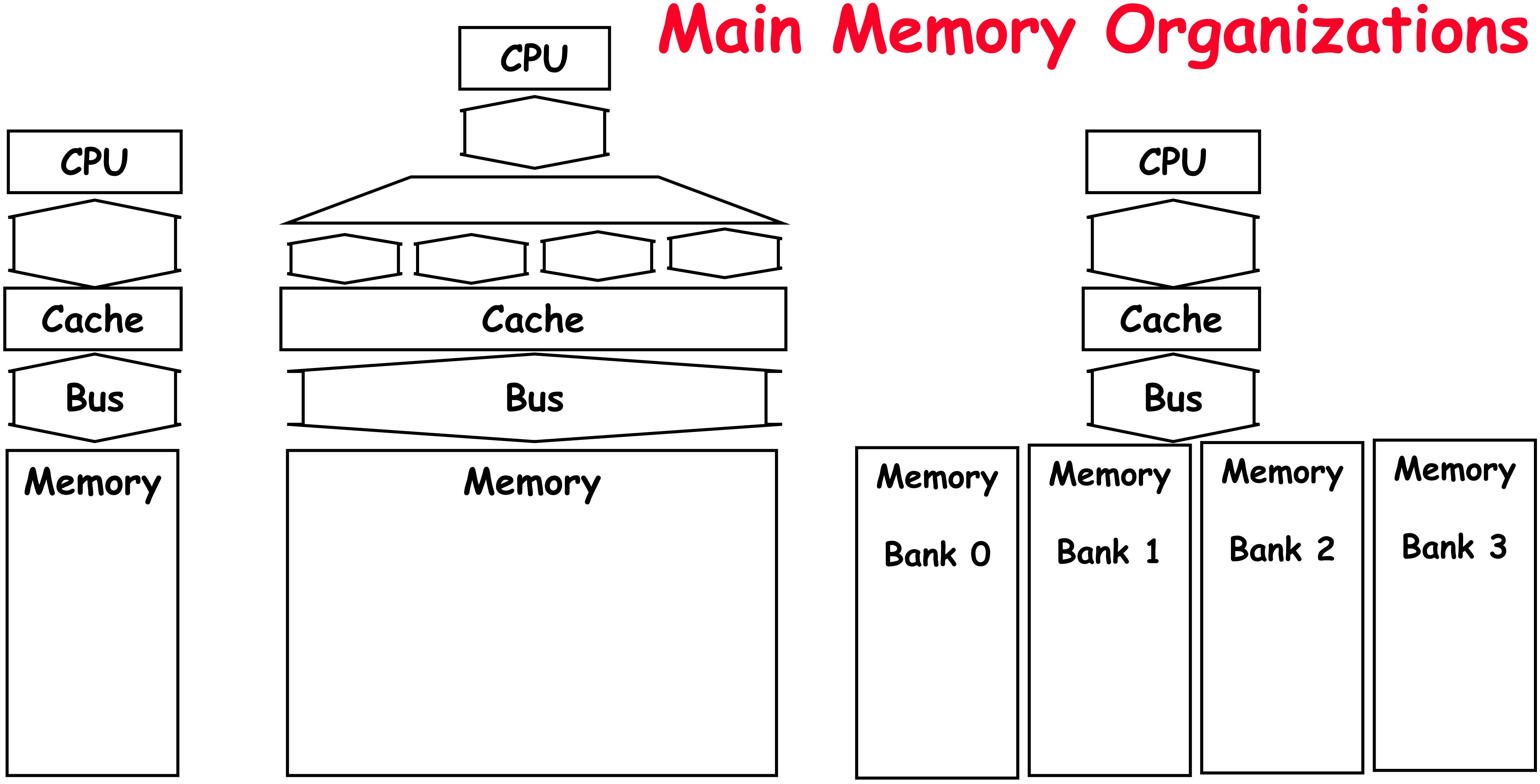
Infineon News 2004-12-14 http://www.infineon.com/cgi/ecrm.dll/jsp/showfrontend.do?lang=EN&news_nav_oid=-9979&content_type=NEWS&content_oid=117074

Infineon Samples 1Gbit DDR SDRAMs *Electronics News*, 8/26/2003 http://www.reed-electronics.com/electronicnews/article/CA318799.html

Fast Memory Systems: DRAM specific

- Multiple CAS accesses: several names (page mode)
  - *Extended Data Out (EDO)*: 30% faster in page mode
- New DRAMs to address gap;
  what will they cost, will they survive?
  - *RAMBUS*: "reinvent DRAM interface"
    - Each Chip a module vs. slice of memory
    - Short bus between CPU and chips
    - Does own refresh
    - Variable amount of data returned
    - Originally 1 byte / 2 ns (500 MB/s per chip)
    - Direct Rambus DRAM (DRDRAM) 16 bits at 400MHz, with a transfer on both clock edges, leading to 1.6GB/s
    - 20% increase in DRAM area
  - *Synchronous DRAM*: 2 banks on chip, a clock signal to DRAM, transfer synchronous to system clock (66 - 150 MHz). "Double Data Rate" DDR SDRAM also transfers on both clock edges
  - Intel claims RAMBUS Direct (16 b wide) is future PC memory?
- Niche memory or main memory?
  - e.g., Video RAM for frame buffers, DRAM + fast serial output

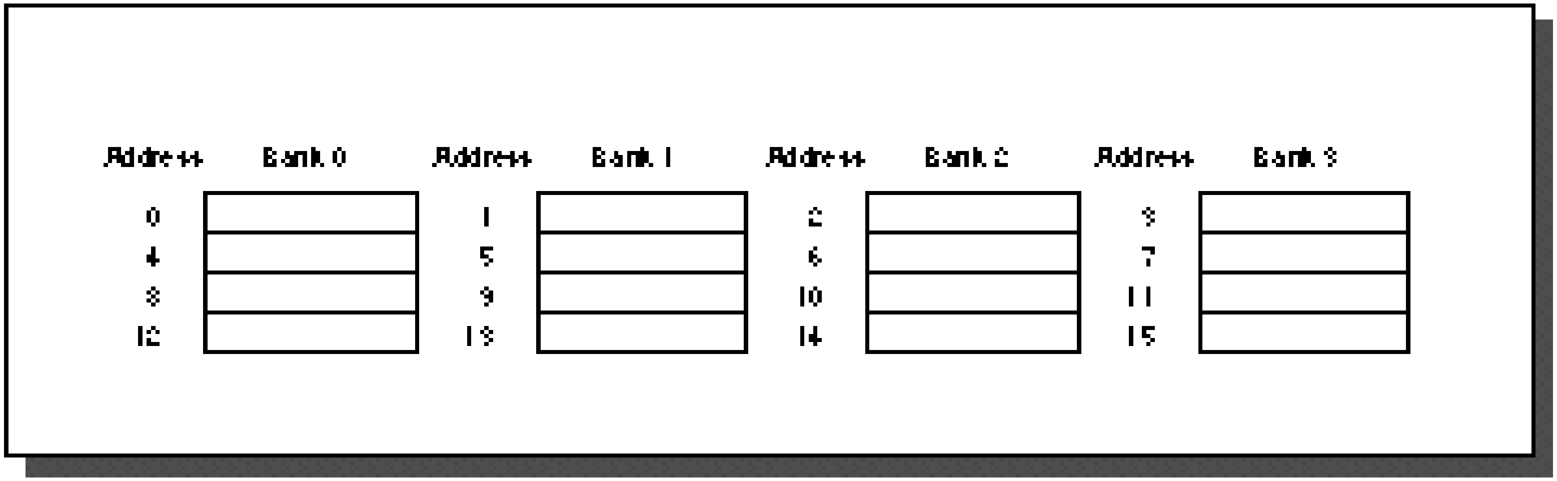
Main Memory Organizations



- ***Simple***:
  - CPU, Cache, Bus, Memory same width (32 or 64 bits)
- ***Wide***:
  - CPU/Mux 1 word; Mux/Cache, Bus, Memory N words (Alpha: 64 bits & 256 bits; UtraSPARC 512)
- ***Interleaved***:
  - CPU, Cache, Bus 1 word: Memory N Modules (4 Modules); example is *word interleaved*

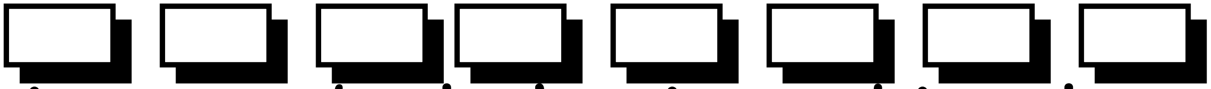
Main Memory Performance

- Timing model (word size is 32 bits)
  - 1 to send address,
  - 6 access time, 1 to send data
  - Cache Block is 4 words
- *Simple M.P.* = 4 x (1+6+1) = 32
- *Wide M.P.* = 1 + 6 + 1 = 8
- *Interleaved M.P.* = 1 + 6 + 4x1 = 11

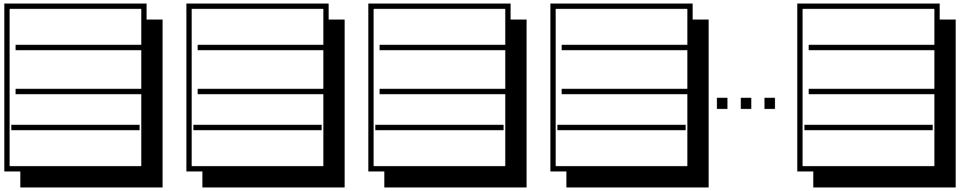


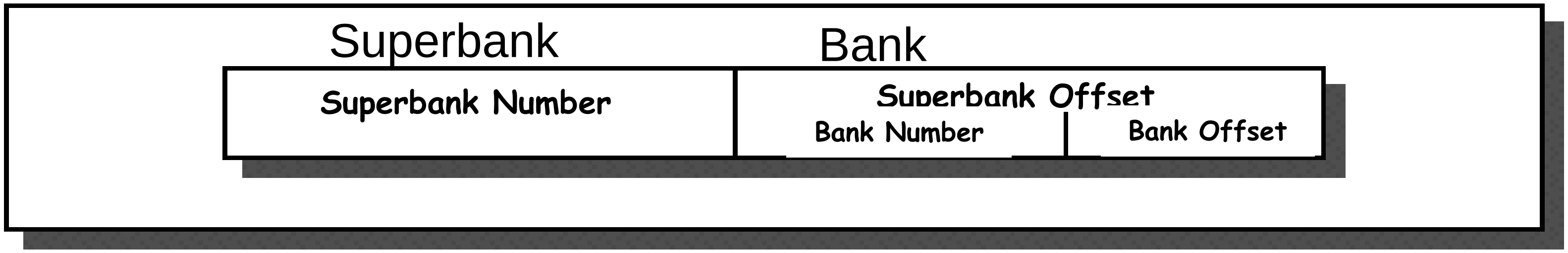
Independent Memory Banks

- Memory banks for independent accesses vs. faster sequential accesses
  - Multiprocessor
  - I/O
  - CPU with Hit under n Misses, Non-blocking Cache
- *Superbank*: all memory active on one block transfer (or *Bank*)



- *Bank*: portion within a superbank that is word interleaved (or *Subbank*)





Independent Memory Banks

- How many banks?

  number banks ≤ number clocks to access word in bank

  - For sequential accesses, otherwise will return to original bank before it has next word ready
  - (like in vector case)
- Increasing DRAM => fewer chips => harder to have banks

Avoiding Bank Conflicts

- Lots of banks

```
int x[256][512];
for (j = 0; j < 512; j = j+1)
      for (i = 0; i < 256; i = i+1)
            x[i][j] = 2 * x[i][j];
```

- Conflicts occur even with 128 banks, since 512 is multiple of 128, conflict on word accesses
- SW: loop interchange or declaring array not power of 2 (“array padding”)
- HW: Prime number of banks
  - bank number = address mod number of banks
  - address within bank = address / number of words in bank
  - modulo & divide per memory access with prime no. banks?
  - address within bank = address mod number words in bank
  - bank number? easy if 2^N words per bank

Fast Bank Number

- **Chinese Remainder Theorem**

  As long as two sets of integers ai and bi follow these rules

  $$b_i = x \bmod a_i, 0 \le b_i < a_i, 0 \le x < a_0 \times a_1 \times a_2 \times \ldots$$

  and that ai and aj are co-prime if i ≠ j, then the integer x has only one solution (unambiguous mapping):

  - bank number = b_0, number of banks = a_0 (= 3 in example)
  - address within bank = b_1, number of words in bank = a_1 (= 8 in example)
  - N word address 0 to N-1, prime no. banks, words power of 2

		Seq. Interleaved			Modulo Interleaved		
Bank Number:		0	1	2	0	1	2
	Address						within
Bank:	*0*	0	1	2	0	16	8
	1	3	4	5	9	1	17
	2	6	7	8	18	10	2
	3	9	10	11	3	19	11
	4	12	13	14	12	4	20
	5	15	16	17	21	13	5
	6	18	19	20	6	22	14
	7	21	22	23	15	7	23

DRAMs per PC over Time

DRAM Generation

Minimum Memory Size	‘86 1 Mb	‘89 4 Mb	‘92 16 Mb	‘96 64 Mb	‘99 256 Mb	‘02 1 Gb
4 MB	32 →	8				
8 MB		16 →	4			
16 MB			8 →	2		
32 MB				4 →	1	
64 MB				8 →	2	
128 MB					4 →	1
256 MB					8 →	2

Need for Error Correction!

- Motivation:
  - Failures/time *proportional* to number of bits!
  - As DRAM cells shrink, more vulnerable
- Went through period in which failure rate was low enough without error correction that people didn't do correction
  - DRAM banks too large now
  - Servers always corrected memory systems
- Basic idea: add redundancy through parity bits
  - Simple but wasteful version:
    - Keep three copies of everything, vote to find right value
    - 200% overhead, so not good!
  - Common configuration: Random error correction
    - SEC-DED (single error correct, double error detect)
    - One example: 64 data bits + 8 parity bits (11% overhead)
    - Papers up on reading list from last term tell you how to do these types of codes
  - Really want to handle failures of physical components as well
    - Organization is multiple DRAMs/SIMM, multiple SIMMs
    - Want to recover from failed DRAM and failed SIMM!
    - Requires more redundancy to do this
    - All major vendors thinking about this in high-end machines

Architecture in practice

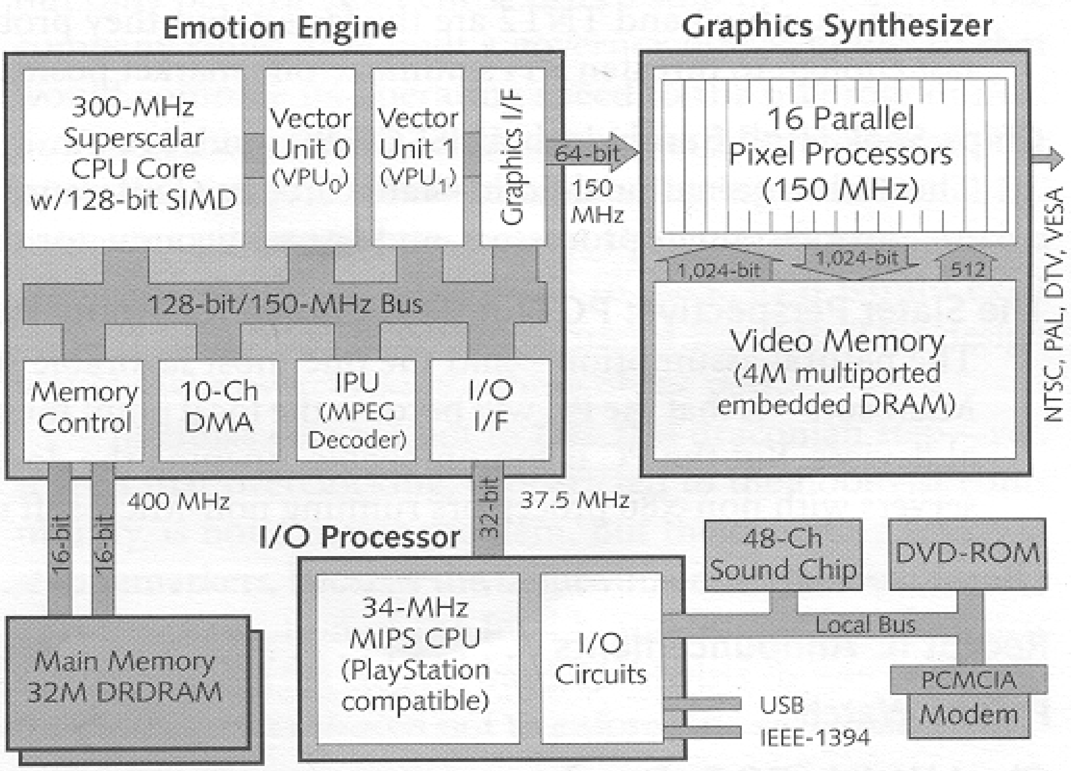


Figure 1. PlayStation 2000 employs an unprecedented level of parallelism to achieve workstation-class 3D performance.

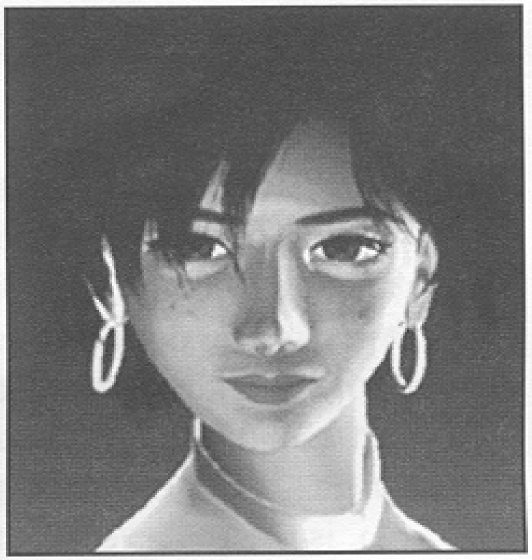


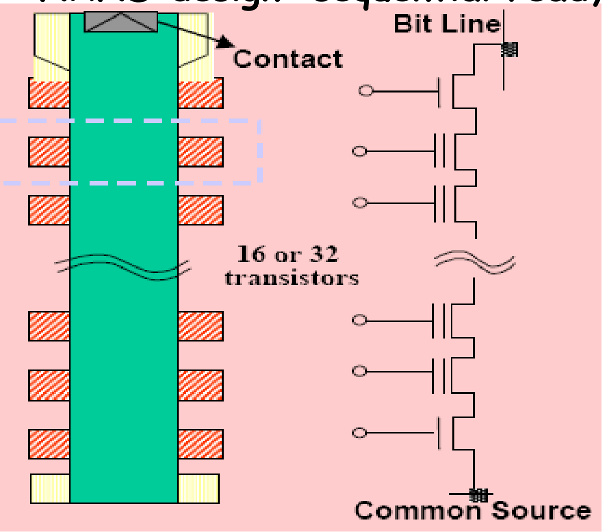
Figure 2. PlayStation 2000 screenshot. (Source: Namco)

- (as reported in Microprocessor Report, Vol 13, No. 5)
    - Emotion Engine: 6.2 GFLOPS, 75 million polygons per second
    - Graphics Synthesizer: 2.4 Billion pixels per second
    - Claim: *Toy Story* realism brought to games!

More esoteric Storage Technologies?

- FLASH
  - Mosfet cell with two gates
  - One "floating"
  - To program, charge tunnels via <7nm dielectric
  - Cells can only be erased (reset to 0) in blocks

NAND design: sequential read, high density

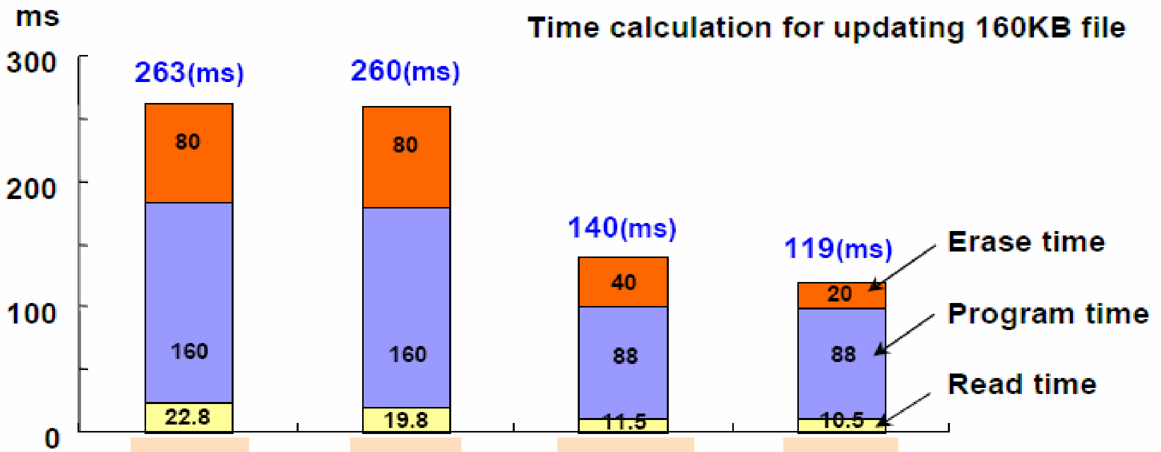


Professional CompactFlash Type II
LEXAR Media
40x
WA WRITE ACCELERATION TECHNOLOGY
8GB

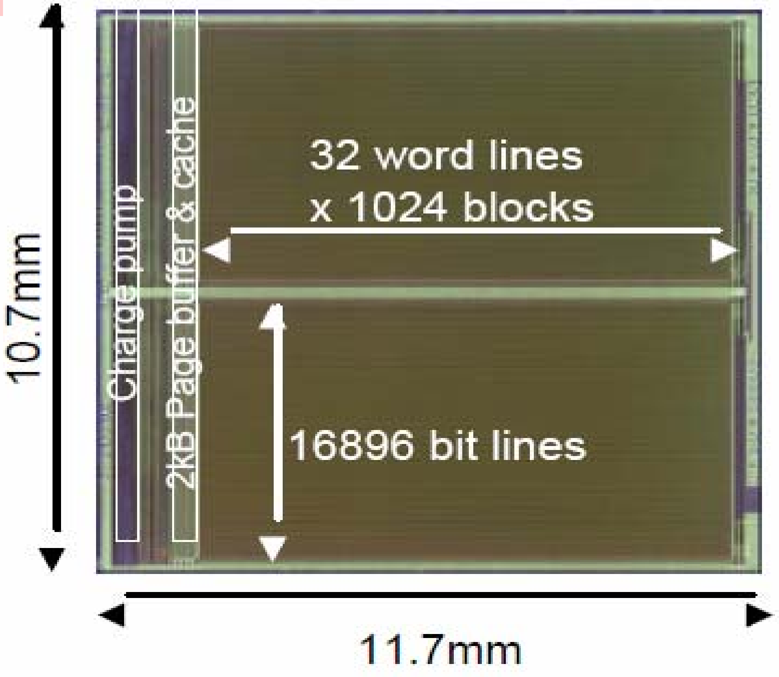
Jan 2005: $1000

16GB Q2'05

Time calculation for updating 160KB file



	8Mb	16Mb	32/64Mb	128Mb
Page size	256Byte	256Byte	512Byte	512Byte
Block size	4KByte	4KByte	8KByte	16KByte
Read/page	10us	10us	10/7us	10us
Program/page	250us	250us	250/200us	200us
Erase/block	2ms	2ms	2ms	2ms

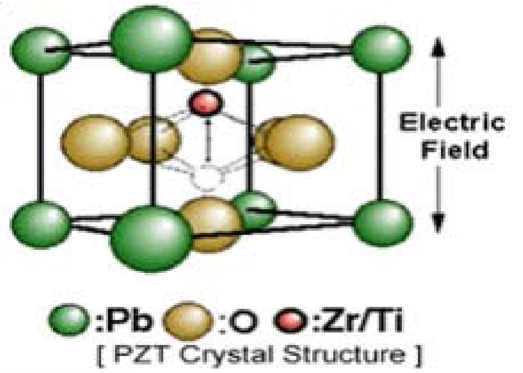


Gbit NAND Flash memory

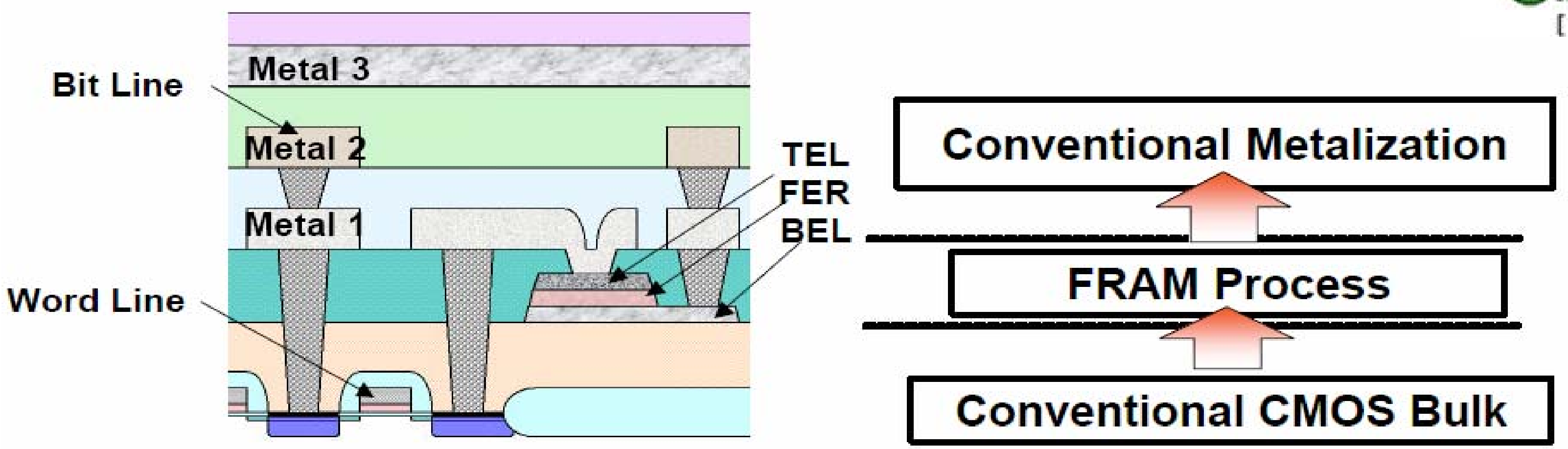
bwrc.eecs.berkeley.edu/Classes/ICDesign/ EE241_s02/Lectures/lecture28-Flash.pdf

More esoteric Storage Technologies?

- FRAM
  - Perovskite ferroelectric crystal forms dielectric in capactor, stores bit via phase change
  - 100ns read, 100ns write
  - Very low write energy (ca.1nJ)



[PZT Crystal Structure]



Conventional Metalization

FRAM Process

Conventional CMOS Bulk

Additional FRAM process
between conventional CMOS bulk and metalization

⇩

Compatible with conventional CMOS technology
and existing CMOS cell libraries

http://www.fma.fujitsu.com/fram/framDocs01.asp?grOut=Documentation&sec=Documentation

- Fully integrated with logic fab process
- Currently used in Smartcards/RFID
- Soon to overtake Flash?

Main Memory Summary

- Wider Memory
- Interleaved Memory: for sequential or independent accesses
- Avoiding bank conflicts: SW & HW
- DRAM specific optimizations: page mode & Specialty DRAM
- Need Error correction