

# Relational Model and Algebra

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## Example Relational Schema

branch		
<u>sortcode</u>	bname	cash
56	'Wimbledon'	94340.45
34	'Goodge St'	8900.67
67	'Strand'	34005.00

movement			
<u>mid</u>	no	amount	tdate
1000	100	2300.00	5/1/1999
1001	101	4000.00	5/1/1999
1002	100	-223.45	8/1/1999
1004	107	-100.00	11/1/1999
1005	103	145.50	12/1/1999
1006	100	10.23	15/1/1999
1007	107	345.56	15/1/1999
1008	101	1230.00	15/1/1999
1009	119	5600.00	18/1/1999

account				
<u>no</u>	type	cname	rate	sortcode
100	'current'	'McBrien, P.'	NULL	67
101	'deposit'	'McBrien, P.'	5.25	67
103	'current'	'Boyd, M.'	NULL	34
107	'current'	'Poulovassilis, A.'	NULL	56
119	'deposit'	'Poulovassilis, A.'	5.50	56
125	'current'	'Bailey, J.'	NULL	56

key branch(sortcode)

key branch(bname)

key movement(mid)

key account(no)

movement(no)  $\xRightarrow{fk}$  account(no)

account(sortcode)  $\xRightarrow{fk}$  branch(sortcode)

# Relational Algebra: A Query Language for the Relational Model

## Primitive operators of the Relational Algebra

Symbol	Name	Type
$\pi$	Project	Unary
$\sigma$	Select	Unary
$\times$	Cartesian Product	Binary
$\cup$	Union	Binary
$-$	Difference	Binary

- All operators take relations as input
- All operators produce one relation as their output
- Other (useful) operators may be defined in terms of the five primitive operators

Relational Algebra: Project  $\pi$ 

account				
<u>no</u>	type	cname	rate	sortcode
100	'current'	'McBrien, P.'	NULL	67
101	'deposit'	'McBrien, P.'	5.25	67
103	'current'	'Boyd, M.'	NULL	34
107	'current'	'Poulovassilis, A.'	NULL	56
119	'deposit'	'Poulovassilis, A.'	5.50	56
125	'current'	'Bailey, J.'	NULL	56

## Project Operator

$\pi_{no,type}$ account	
<u>no</u>	type
100	'current'
101	'deposit'
103	'current'
107	'current'
119	'deposit'
125	'current'

$\pi_{sortcode}$ account
sortcode
67
34
56

Relational Algebra: Select  $\sigma$ 

account				
<u>no</u>	type	cname	rate	sortcode
100	'current'	'McBrien, P.'	NULL	67
101	'deposit'	'McBrien, P.'	5.25	67
103	'current'	'Boyd, M.'	NULL	34
107	'current'	'Poulovassilis, A.'	NULL	56
119	'deposit'	'Poulovassilis, A.'	5.50	56
125	'current'	'Bailey, J.'	NULL	56

## Select Operator

$\sigma_{\text{rate}>0}\text{account}$				
<u>no</u>	type	cname	rate	sortcode
101	'deposit'	'McBrien, P.'	5.25	67
119	'deposit'	'Poulovassilis, A.'	5.50	56

Relational Algebra: Product  $\times$ 

branch		
sortcode	bname	cash
56	'Wimbledon'	94340.45
34	'Goodge St'	8900.67
67	'Strand'	34005.00

$\sigma_{\text{rate}>0}\text{account}$				
no	type	cname	rate	sortcode
101	'deposit'	'McBrien, P.'	5.25	67
119	'deposit'	'Poulovassilis, A.'	5.50	56

## Product Operator

branch $\times$ $\sigma_{\text{rate}>0}\text{account}$							
sortcode	bname	cash	no	type	cname	rate	sortcode
56	'Wimbledon'	94340.45	101	'deposit'	'McBrien, P.'	5.25	67
56	'Wimbledon'	94340.45	119	'deposit'	'Poulovassilis, A.'	5.50	56
34	'Goodge St'	8900.67	101	'deposit'	'McBrien, P.'	5.25	67
34	'Goodge St'	8900.67	119	'deposit'	'Poulovassilis, A.'	5.50	56
67	'Strand'	34005.00	101	'deposit'	'McBrien, P.'	5.25	67
67	'Strand'	34005.00	119	'deposit'	'Poulovassilis, A.'	5.50	56

## Quiz 1: RA Queries

branch		
sortcode	bname	cash
56	'Wimbledon'	94340.45
34	'Goodge St'	8900.67
67	'Strand'	34005.00

account				
no	type	cname	rate	sortcode
100	'current'	'McBrien, P.'	NULL	67
101	'deposit'	'McBrien, P.'	5.25	67
103	'current'	'Boyd, M.'	NULL	34
107	'current'	'Poulovassilis, A.'	NULL	56
119	'deposit'	'Poulovassilis, A.'	5.50	56
125	'current'	'Bailey, J.'	NULL	56

Which RA query lists the name of branches that have deposit accounts?

A

 $\pi_{\text{sortcode}} \sigma_{\text{type}='deposit'} \text{account}$ 

B

 $\pi_{\text{bname}}$ 
 $\sigma_{\text{account.sortcode}=\text{branch.sortcode} \wedge \text{type}='deposit'}$   
 $(\text{account} \times \text{branch})$ 

C

 $\pi_{\text{bname}} (\text{branch} \times \sigma_{\text{type}='deposit'} \text{account})$ 

D

 $\pi_{\text{bname}} \sigma_{\text{type}='deposit'} (\text{account} \times \text{branch})$

# SPJ Queries

## Select Project Join (SPJ) queries

If a product of tables is formed, where a selection is then done that compares the attributes of those tables, we say that a **join** has been performed.

Normally not all columns of the product are returned, and therefore a project is also required.

## Branches with current accounts

$\pi_{\text{bname, no}} \sigma_{\text{branch.sortcode}=\text{account.sortcode} \wedge \text{account.type}=\text{'current'}} (\text{branch} \times \text{account})$

bname	no
'Goodge St'	103
'Wimbledon'	107
'Wimbledon'	125
'Strand'	100

Relational Algebra: Union  $\cup$ 

$\pi_{\text{sortcode as id}} \text{account}$	$\pi_{\text{no as id}} \text{account}$
id	id
67	100
34	101
56	103
	107
	119
	125

Union Operator	
$\pi_{\text{sortcode as id}} \text{account} \cup \pi_{\text{no as id}} \text{account}$	id
	67
	34
	56
	100
	101
	103
	107
	119
	125

- relations must be **union compatible**

## Relational Algebra: Difference –

$\pi_{\text{no}}\text{account}$	$\pi_{\text{no}}\text{movement}$
no	no
100	100
101	101
103	103
107	107
119	119
125	

## Difference Operator

 $\pi_{\text{no}}\text{account} - \pi_{\text{no}}\text{movement}$ 

<u>no</u>
125

- relations must be **union compatible**

## Rules for Combining Operators

Since all operators produce a relation as output, *any* operator may produce one of the inputs to any other operator.

### well formed RA query

- the output of the nested operator must contain the attributes required by an outer  $\pi$  or  $\sigma$
- the two inputs to a  $\cup$  or  $-$  must contain the same number of attributes

Derived Relational Algebra: Natural Join  $\bowtie$ 

## Natural Join

$$R \bowtie S = \sigma_{R.A_1=S.A_1 \wedge \dots \wedge R.A_m=S.A_m} R \times S$$

## Natural Join

$$\text{branch} \bowtie \text{account} = \sigma_{\text{branch.sortcode}=\text{account.sortcode}} \text{branch} \times \text{account}$$

branch $\bowtie$ account						
sortcode	bname	cash	no	type	cname	rate
34	'Goodge St'	8900.67	103	'current'	'Boyd, M.'	NULL
56	'Wimbledon'	94340.45	107	'current'	'Poulovassilis, A.'	NULL
56	'Wimbledon'	94340.45	119	'deposit'	'Poulovassilis, A.'	5.50
56	'Wimbledon'	94340.45	125	'current'	'Bailey, J.'	NULL
67	'Strand'	34005.00	100	'current'	'McBrien, P.'	NULL
67	'Strand'	34005.00	101	'deposit'	'McBrien, P.'	5.25

## Quiz 2: Natural Join

What is the result of  $\pi_{no}(\text{account} \bowtie \text{movement})$ ?

A

 $\pi_{no}(\text{account} \bowtie \text{movement})$ 

no
100
101
103
107
119
125

B

 $\pi_{no}(\text{account} \bowtie \text{movement})$ 

no
100
101
103
107
119

C

 $\pi_{no}(\text{account} \bowtie \text{movement})$ 

no
125

D

 $\pi_{no}(\text{account} \bowtie \text{movement})$ 

no
----

Derived Relational Algebra: Semi Join  $\ltimes$ 

## Semi Join

$$R \ltimes S = R \bowtie \pi_{Attr(R) \cap Attr(S)}(S)$$

## Semi Join

$$\text{account} \ltimes \text{movement} = \text{account} \bowtie \pi_{no}(\text{movement})$$

account $\ltimes$ movement				
<u>no</u>	type	cname	rate	sortcode
100	'current'	'McBrien, P.'	NULL	67
101	'deposit'	'McBrien, P.'	5.25	67
103	'current'	'Boyd, M.'	NULL	34
107	'current'	'Poulovassilis, A.'	NULL	56
119	'deposit'	'Poulovassilis, A.'	5.50	56

# Derived Relational Algebra: Joins

## Natural Join

$$R \bowtie S = \sigma_{R.A_1=S.A_1 \wedge \dots \wedge R.A_m=S.A_m} R \times S$$

## Equi Join

$$R \stackrel{A=B}{\bowtie} S = \sigma_{R.A=S.B} R \times S$$

## Semi Join

$$R \ltimes S = R \bowtie \pi_{Attr(R) \cap Attr(S)}(S)$$

## Theta Join

$$R \stackrel{\theta}{\bowtie} S = \sigma_{\theta} R \times S$$

## Quiz 3: Understanding join operators

branch		
sortcode	bname	cash
56	'Wimbledon'	94340.45
34	'Goodge St'	8900.67
67	'Strand'	34005.00

account				
no	type	cname	rate	sortcode
100	'current'	'McBrien, P.'	NULL	67
101	'deposit'	'McBrien, P.'	5.25	67
103	'current'	'Boyd, M.'	NULL	34
107	'current'	'Poulovassilis, A.'	NULL	56
119	'deposit'	'Poulovassilis, A.'	5.50	56
125	'current'	'Bailey, J.'	NULL	56

Which RA query produces the most tuples?

A

branch  $\bowtie_{\text{branch.sortcode} < \text{account.sortcode}}$  account

B

branch  $\bowtie$  account

C

branch  $\bowtie$  account

D

branch  $\bowtie_{\text{branch.sortcode} = \text{account.sortcode}}$  account

## Quiz 4: Foreign Keys and Natural Joins (1)

Suppose  $R$  and  $S$  only share attribute  $A$ , and there is a foreign key  $S(A) \xrightarrow{fk} R(A)$ .

If  $|R| = 100$  and  $|S| = 1,000$ , what is  $|R \bowtie S|$ ?

A

100

B

1,000

C

100,000

D

900

Note that  $|R|$  returns the number of tuples in the current extent of  $R$

## Quiz 5: Foreign Keys and Natural Joins (2)

Suppose  $R$  and  $S$  only share attribute  $A$ , and there is a foreign key  $R(A) \xrightarrow{fk} S(A)$ .

If  $|R| = 100$  and  $|S| = 1,000$ , what is  $|R \bowtie S|$ ?

A  
100

B  
1,000

C  
100,000

D  
900

Derived Relational Algebra: Intersection  $\cap$ 

## Intersection

$$R \cap S = R - (R - S)$$

 $\pi_{\text{no}}\text{account} \cap \pi_{\text{no}}\text{movement}$ 

$\pi_{\text{no}}\text{account}$
no
100
101
103
107
119
125

$\pi_{\text{no}}\text{account} - \pi_{\text{no}}\text{movement}$
<u>no</u>
125

$\pi_{\text{no}}\text{account} \cap \pi_{\text{no}}\text{movement}$
no
100
101
103
107
119

## Quiz 6: Intersection

name	email
'McBrien, P.'	p.mcbrien@imperial.ac.uk
'Poulovassilis, A.'	ap@dcs.bbk.ac.uk
'Pietzuch, P.'	prp@doc.ic.ac.uk

What is the result of  $\pi_{\text{name}} \text{account} \cap \pi_{\text{name}} \text{email}$ ?

A

cname

'McBrien, P.'

'Boyd, M.'

'Poulovassilis, A.'

'Bailey, J.'

'Pietzuch, P.'

B

cname

'McBrien, P.'

'Boyd, M.'

'Poulovassilis, A.'

'Bailey, J.'

C

cname

'McBrien, P.'

'Poulovassilis, A.'

'Pietzuch, P.'

D

cname

'McBrien, P.'

'Poulovassilis, A.'

## Worksheet: Relational Algebra Operators

branch		
sortcode	bname	cash
56	'Wimbledon'	94340.45
34	'Goodge St'	8900.67
67	'Strand'	34005.00

movement			
mid	no	amount	tdate
1000	100	2300.00	5/1/1999
1001	101	4000.00	5/1/1999
1002	100	-223.45	8/1/1999
1004	107	-100.00	11/1/1999
1005	103	145.50	12/1/1999
1006	100	10.23	15/1/1999
1007	107	345.56	15/1/1999
1008	101	1230.00	15/1/1999
1009	119	5600.00	18/1/1999

account				
no	type	cname	rate	sortcode
100	'current'	'McBrien, P.'	NULL	67
101	'deposit'	'McBrien, P.'	5.25	67
103	'current'	'Boyd, M.'	NULL	34
107	'current'	'Poulovassilis, A.'	NULL	56
119	'deposit'	'Poulovassilis, A.'	5.50	56
125	'current'	'Bailey, J.'	NULL	56

key branch(sortcode)

key branch(bname)

key movement(mid)

key account(no)

movement(no)  $\xRightarrow{fk}$  account(no)

account(sortcode)  $\xRightarrow{fk}$  branch(sortcode)

## Equivalences Involving Project

### Project and Project

$$\pi_{\vec{X}} \pi_{\vec{Y}} R \equiv \pi_{\vec{X}} R$$

You can eliminate any inner project (note that to be well formed  $\vec{X} \subseteq \vec{Y}$ )

### Project and Select

$$\pi_{\vec{X}} \sigma_{P(\vec{Y})} R \equiv \sigma_{P(\vec{Y})} \pi_{\vec{X}} R$$

You can move a project of attributes  $\vec{X}$  inside a select, provided the select predicate can be answered from those attributes, *i.e.*  $\vec{Y} \subseteq \vec{X}$

### Project and Product

$$\pi_{\vec{X}}(R \times S) \equiv \pi_{\vec{X} \cap \text{Atts}(R)} R \times \pi_{\vec{X} \cap \text{Atts}(S)} S$$

### Project and Union

$$\pi_{\vec{X}}(R \cup S) \equiv \pi_{\vec{X}} R \cup \pi_{\vec{X}} S$$

### Project and Difference

$$\pi_{\vec{X}}(R - S) \supseteq \pi_{\vec{X}} R - \pi_{\vec{X}} S$$

## Equivalences Involving Select

### Select and Project

$$\sigma_{P(\vec{X})} \pi_{\vec{X}} R \equiv \pi_{\vec{X}} \sigma_{P(\vec{X})} R$$

### Select and Select

$$\sigma_{P_x(\vec{X})} \sigma_{P_y(\vec{Y})} R \equiv \sigma_{P_x(\vec{X}) \wedge P_y(\vec{Y})} R$$

### Select and Product

$$\sigma_{P(\vec{X})} (R \times S) \equiv \sigma_{P(\vec{X})} R \times S \iff \vec{X} \subseteq \text{Atts}(R)$$

You can move a select predicate  $P(\vec{X})$  onto one of the relations inside a product provided  $\vec{X} \subseteq \text{Atts}(R)$ .

### Select and Union

$$\sigma_{P(\vec{X})} (R \cup S) \equiv \sigma_{P(\vec{X})} R \cup \sigma_{P(\vec{X})} S$$

### Select and Difference

$$\sigma_{P(\vec{X})} (R - S) \equiv \sigma_{P(\vec{X})} R - S$$

## Quiz 7: Equivalent RA Expressions (Unary Operators)

Which RA expression is not equivalent to the other three?

A

 $\pi_{no} \sigma_{type='current'} \mathbf{account}$ 

B

 $\pi_{no} \sigma_{type='current'} \pi_{no,type,cname} \mathbf{account}$ 

C

 $\pi_{no} \sigma_{type <> 'deposit'} \pi_{no,type,cname} \mathbf{account}$ 

D

 $\pi_{no} \sigma_{type='current'} \sigma_{type <> 'deposit'} \mathbf{account}$

# Equivalences Involving Binary Operators

## Product and Union

$$R \times (S \cup T) \equiv (R \times S) \cup (R \times T)$$

## Product and Difference

$$R \times (S - T) \equiv (R \times S) - (R \times T)$$

## Union and Product

$R \cup (S \times T)$  unable to move  $\cup$  inside  $\times$

## Union and Difference

$R \cup (S - T)$  unable to move  $\cup$  inside  $-$

## Difference and Product

$R - (S \times T)$  unable to move  $-$  inside  $\times$

## Difference and Union

$$R - (S \cup T) \equiv (R - S) - T$$

## Quiz 8: Equivalent RA Expressions (Binary Operators)

Which equivalence does not hold?

A

$$(R \times S) \times T \equiv R \times (S \times T)$$

B

$$(R - S) - T \equiv R - (S - T)$$

C

$$(R \cup S) \cup T \equiv R \cup (S \cup T)$$

D

$$(R \cap S) \cap T \equiv R \cap (S \cap T)$$

## Worksheet: Equivalences Between RA Expressions

- 1  $\pi_{no,type} \sigma_{sortcode=56} \pi_{no,type,sortcode} \sigma_{type='deposit'} \mathbf{account}$
- 2  $\sigma_{account.no=movement.no} (\pi_{no,cname} \mathbf{account} \times \pi_{mid,no} \sigma_{amount>1000} \mathbf{movement})$
- 3  $\sigma_{account.no=movement.no} (\pi_{no,cname,rate} \mathbf{account} \times$   
 $(\sigma_{amount>1000} \pi_{mid,no} \mathbf{movement} \cup \sigma_{amount<100} \pi_{mid,no} \mathbf{movement}))$
- 4  $\pi_{no,cname,tdate} \sigma_{amount<0 \wedge account.no=movement.no} \mathbf{account} \times \mathbf{movement}$