

ER to Relational Mapping

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Designing a Relational Database Schema

How do you design a relational database schema for a particular UoD?

- 1 Need some way to model the semantics of the UoD as a conceptual schema
 - ER (many variants exist)
 - UML class diagrams
- 2 Need to map the ER/UML schema into a relational schema
- 3 Need to ensure that the relational schema is a good design
 - Normalisation

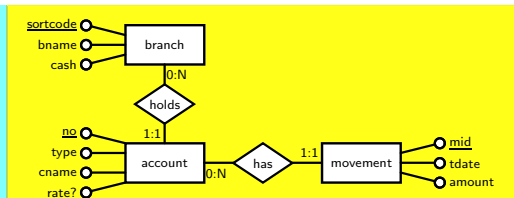
Semantic Modelling: ER Schemas

```
CREATE TABLE branch
(
  sortcode INTEGER NOT NULL,
  bname VARCHAR(20) NOT NULL,
  cash DECIMAL(10,2) NOT NULL,
  CONSTRAINT branch_pk PRIMARY KEY (sortcode)
)

CREATE TABLE account
(
  no INTEGER NOT NULL,
  type CHAR(8) NOT NULL,
  cname VARCHAR(20) NOT NULL,
  rate DECIMAL(4,2) NULL,
  sortcode INTEGER NOT NULL,
  CONSTRAINT account_pk PRIMARY KEY (no),
  CONSTRAINT account_fk FOREIGN KEY (sortcode) REFERENCES branch
)

CREATE INDEX account_type ON account (type)

CREATE TABLE movement
(
  mid INTEGER NOT NULL,
  no INTEGER NOT NULL,
  amount DECIMAL(10,2) NOT NULL,
  tdate DATETIME NOT NULL,
  CONSTRAINT movement_pk PRIMARY KEY (mid),
  CONSTRAINT movement_fk FOREIGN KEY (no) REFERENCES account
)
```



Mapping an ER Schema to a Relational Schema

Borrowing the terminology of **object relational mapping (ORM)**

- **table per type (TPT)**

Each ER entity is mapped to a corresponding single relation, with all directly connected attributes and relationships of that entity.

- **table per hierarchy (TPH)**

Each ER entity that is not a subclass of another ER entity is mapped to a single relation with all the data of the subclasses in the same table.

- **table per concrete type (TPC)**

Each ER entity is mapped to a corresponding single relation, with all attributes and relationships of any superclasses included in that relation.

- **table per modelling construct (TPM)**

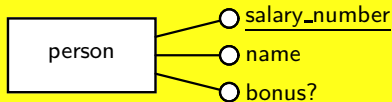
Store all objects in the ER Schema in a fixed set of tables in a relational model for the various ER modelling constructs.

Mapping ER^{KLOS} to a relational model: entities and attributes

Taking a TPT approach, there is a simple mapping of entities and attributes to tables and columns:

- 1 Each entity E maps to a table R_E
- 2 Each attribute A maps to a column C_A of R_E
- 3 If A is an optional attribute, then C_A is nullable, otherwise C_A is not nullable
- 4 If \vec{K} are key attribute(s), then \vec{C}_K are a key of R_E

Tables generated from entities



person(salary_number, name, bonus?)
department(dname)

Mapping ER^{KLOS} to a relational model: relationships

Taking a TPT approach, for each relationship R between E_1, E_2 , entities E_1, E_2 map to R_1, R_2 as before, and

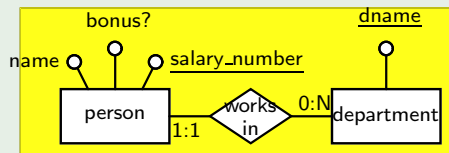
1 If R is a many-many relationship then it maps to

- 1 a table $R_R_1_R_2(\vec{K}_1, \vec{K}_2)$
- 2 a foreign key $R_R_1_R_2(\vec{K}_1) \rightarrow R_1(\vec{K}_1)$
- 3 a foreign key $R_E_1_E_2(\vec{K}_2) \rightarrow E_2(\vec{K}_2)$

2 If R is a one-many relationship then it maps to

- 1 a column \vec{K}_2 in R_1
- 2 a foreign key $R_1(\vec{K}_2) \rightarrow R_2(\vec{K}_2)$
- 3 if the participation of E_1 in R is optional, then \vec{K}_2 is an optional column of R_1

Tables generated from relationships



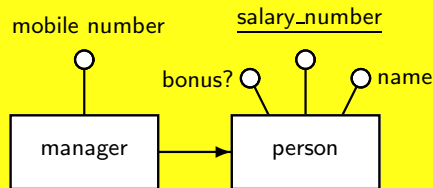
```
person(salary_number, name, bonus?, dname)
department(dname)
person(dname) → department(dname)
```

Mapping ER^{KLOS} to a relational model: subsets

Taking a TPT approach, for each subset E_s of E , entities E_s, E map to tables R_s, R as before and:

- 1 a key \vec{K} in R_s (where \vec{K} is the key of R)
- 2 a foreign key $R_s(\vec{K}) \rightarrow R(\vec{K})$

Tables generated from subsets



```
person(salary_number, name, bonus?)  
manager(salary_number, mobile_phone)  
manager(salary_number) →  
    person(salary_number)
```

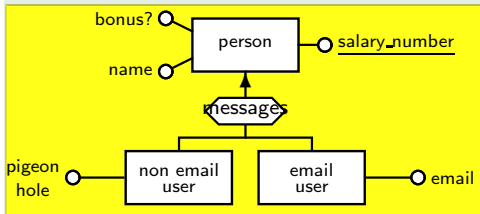
Mapping ER^{KLOS} to a relational model

Mapping ER^D to a relational model

Taking a TPT approach, if E is a generalisation of E_1, \dots, E_n , then entities E_1, \dots, E_n, E map to tables R_1, \dots, R_n, R as before and:

- 1 treat each $E_x \in E_1, \dots, E_n$ as a subset of E
- 2 no implementation of disjointness using just PKs and FKs

Tables generated from generalisations

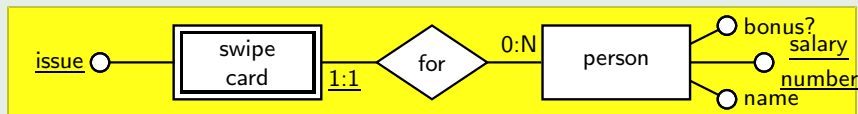


```
person(salary_number, name, bonus?)
non_email_user(salary_number, pigeon_hole)
non_email_user(salary_number) →
    person(salary_number)
email_user(salary_number, email)
email_user(salary_number) →
    person(salary_number)
```

Mapping ER^W to a relational model

- If E_W is a weak entity that maps to a relation R_W , the foreign key R_K due to the participation in a relationship is also used in the key of R_K

Tables generated from weak entities



person(salary_number, name, bonus?)

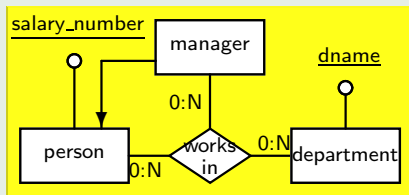
swipe_card(salary_number, issue)

swipe_card(salary_number) → person(salary_number)

Mapping $ER^{\mathcal{H}}$ to a relational model

Rules for binary relationship R between E_1, E_2 generalise to rules for R between E_1, \dots, E_n

Tables generated from n -ary entities



person(salary_number)

manager(salary_number)

manager(salary_number) \rightarrow person(salary_number)

department(dname)

works_in(person_salary_number, manager_salary_number, dname)

works_in(person_salary_number) \rightarrow person(salary_number)

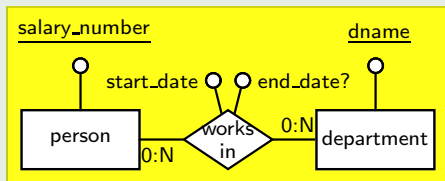
works_in(manager_salary_number) \rightarrow manager(salary_number)

works_in(dname) \rightarrow department(dname)

Mapping ER^A to a relational model

Attributes of a relationship go on the same table as that which implements the relationship

Tables generated from attributes of relationships



person(salary_number)

department(dname)

works_in(salary_number, dname, start_date, end_date?)

works_in(salary_number) → person(salary_number)

works_in(dname) → department(dname)

Mapping ER^{ADHKLOS}W to a relational model