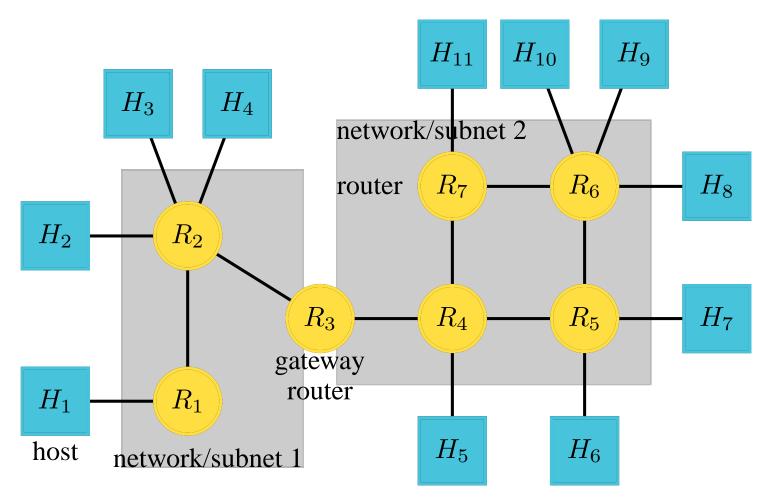
Routing and Internetworking

#### **A Network**



- **gateway** links networks together to form **internet**
- **routing** plans path through network

#### **Internets**

#### An internet can be

#### homogeneous internets

- all same type of network
- no need to convert frames

#### heterogeneous internets

- free mix of network technologies
- need to convert between frame formats
- convert between CO and CL?

#### **Switching the Data**

#### **circuit** switched networks

- physical communication channel to be set up between two hosts
- wasteful if bursty data
- e.g. PSTN

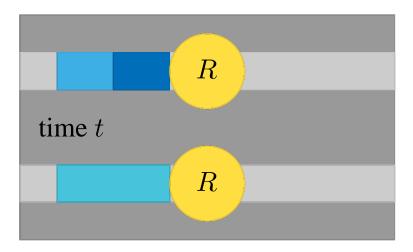
#### message switched networks

- sends data in **messages**, which can share a channel
- store-and-forward
- e.g. Post Office

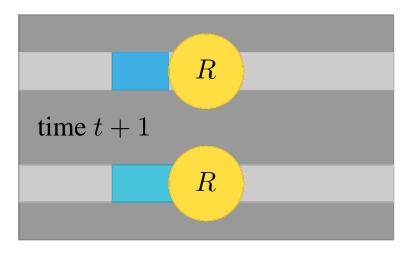
#### packet switched networks

- similar to message switching
- limit on size  $\rightarrow$  data sent in one or more **packets**

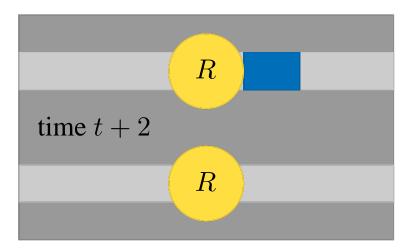
# Store-and-forward (1)



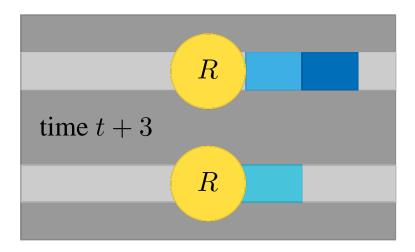
# Store-and-forward (2)



# **Store-and-forward (3)**



# **Store-and-forward (4)**



 $\blacksquare$  delay  $\propto$  packet size

#### **Example: Reducing Store-and-forward Delays**

A router operating on a PSN using store-and-forward has links into it operating at  $1Mbs^{-1}$ , and the router is found to delay packets by 0.1s. How many pieces must you fragment the packet into to reduce this delay to 0.02s, and how large are those packets?

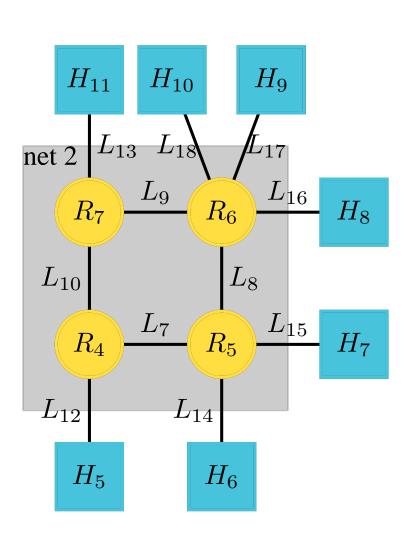
 $delay \propto packet size$ 

$$\frac{delay}{packet \ size} = constant$$

 $0.1/1 = 0.02/s \rightarrow s = 0.2$ , hence 5 packets are required for each of the old packets

packet size= delay = 
$$\frac{0.02 \times 1 \times 10^6}{8}$$
 = 2500 bytes

#### **Virtual Circuits**



 $R_7$  to/from from/to

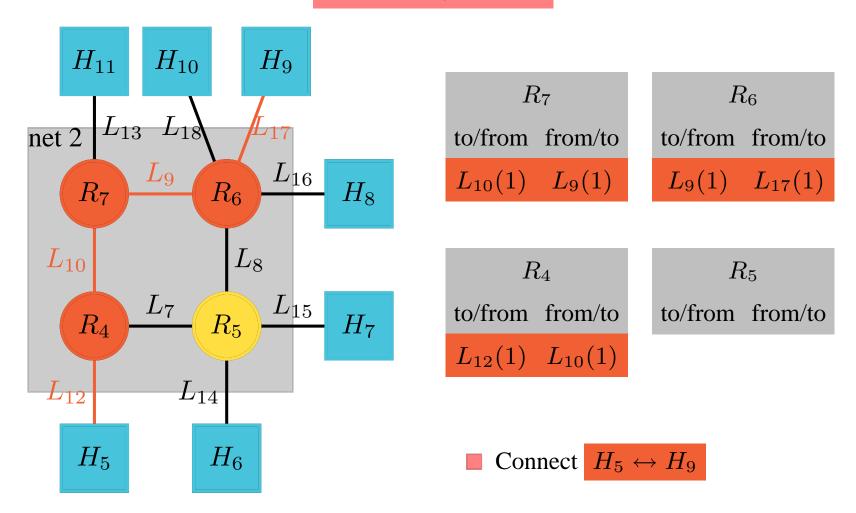
 $$R_{6}$$  to/from from/to

 $R_4$  to/from from/to

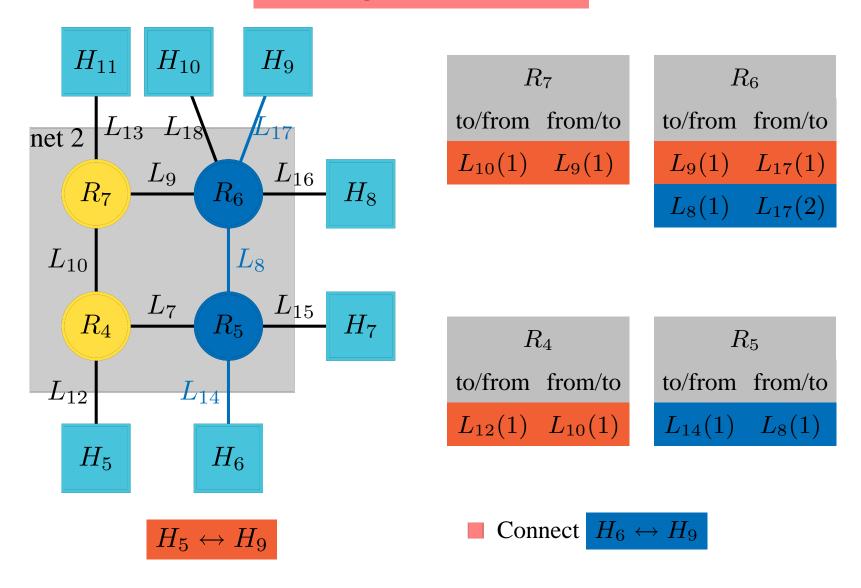
 $R_5$  to/from from/to

- Virtual Circuits allow some features of CO to be present in PSN
- Divide links by virtual circuit identifiers (VCI)  $L_x(1), L_x(2), L_x(3), \dots$
- Connect links up VCIs

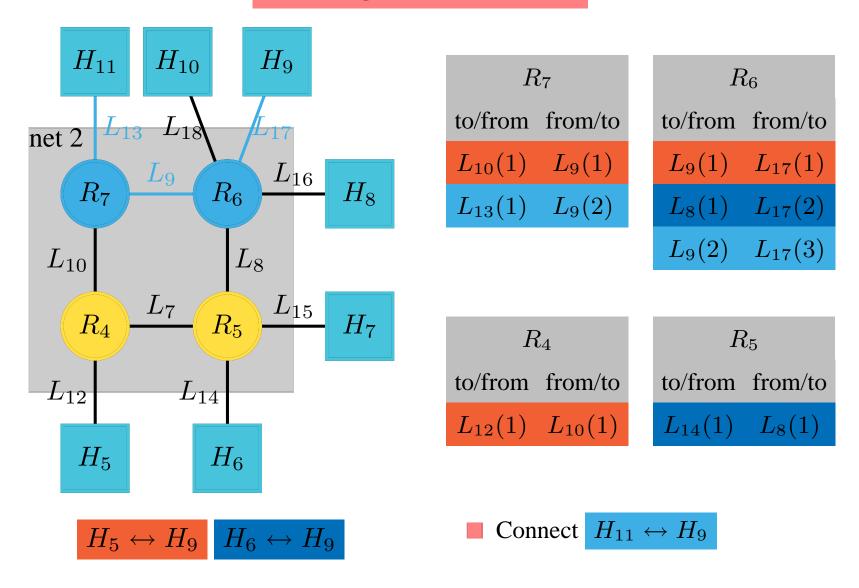
#### **Virtual Circuits**



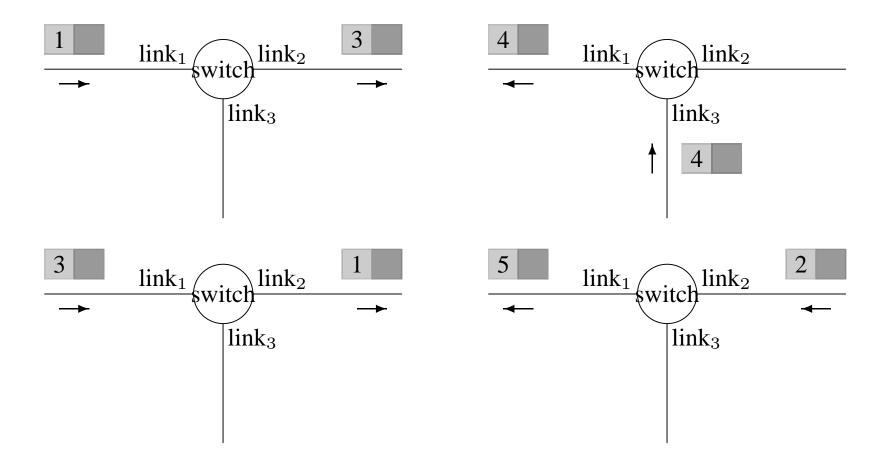
### **Routing Tables for VC**



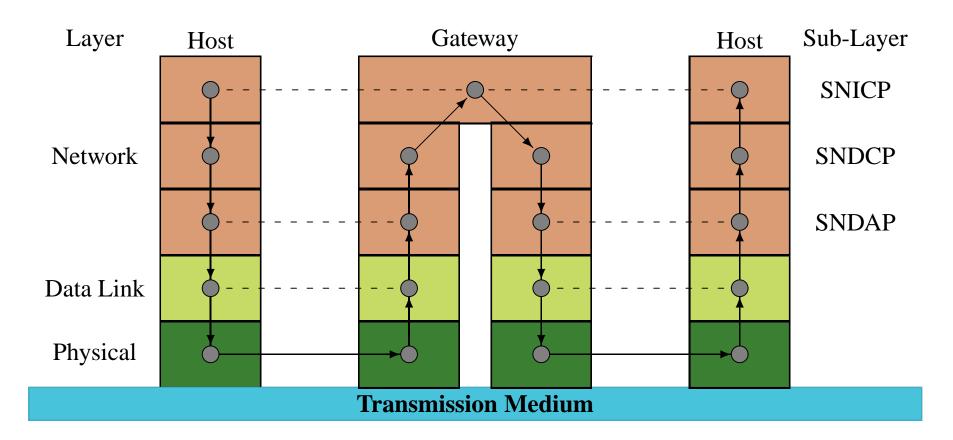
### **Routing Tables for VC**



#### **Worksheet: Virtual Circuit Identifiers**

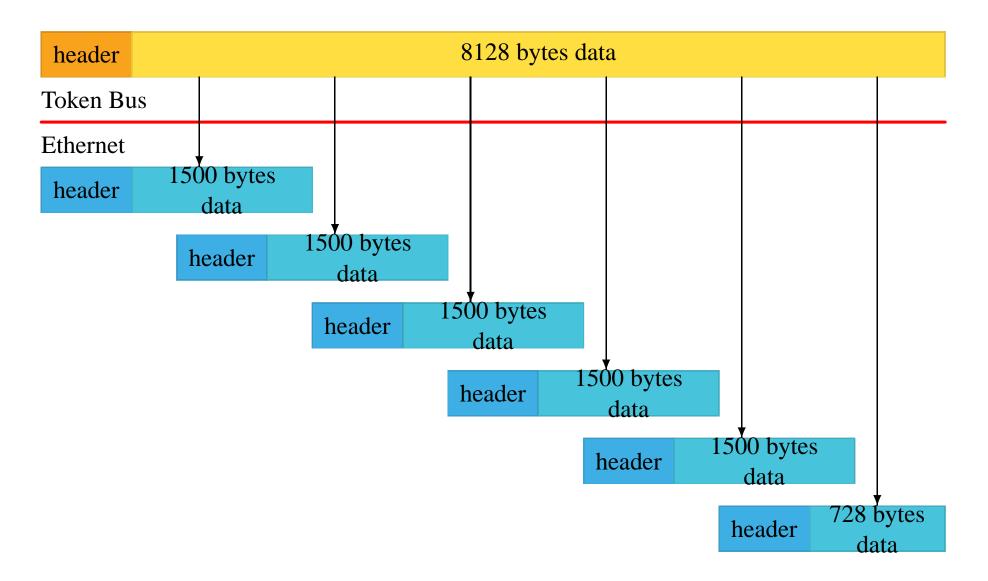


#### The OSI Internetworking Model

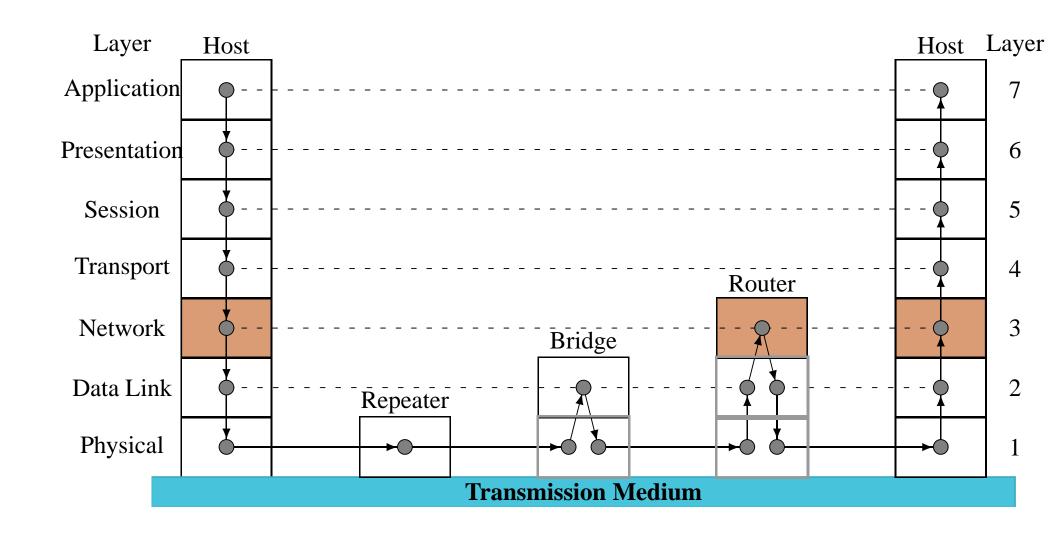


- Subnetwork Independent Convergence Protocol (SNICP)
- **Subnetwork Dependent Convergence Protocol (SNDCP)**
- Subnetwork Dependent Access Protocol (SNDAP)

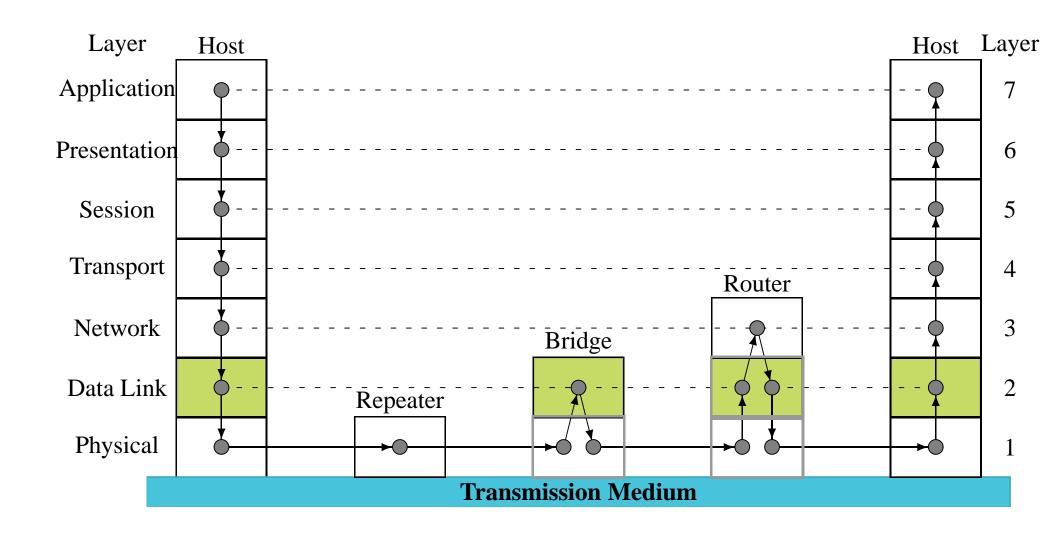
## **Fragmentation**



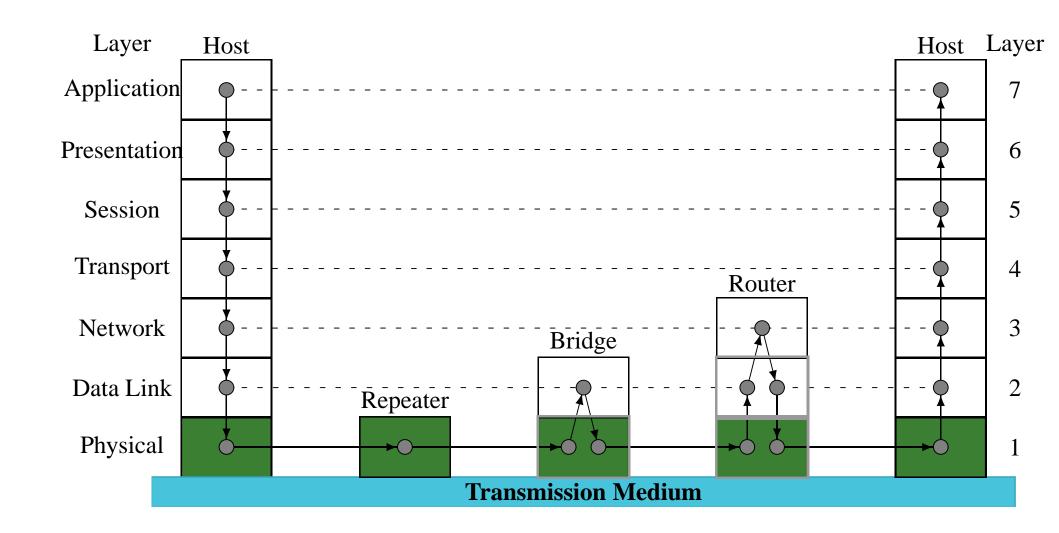
#### **Homogeneous Internetworking: Network Layer**



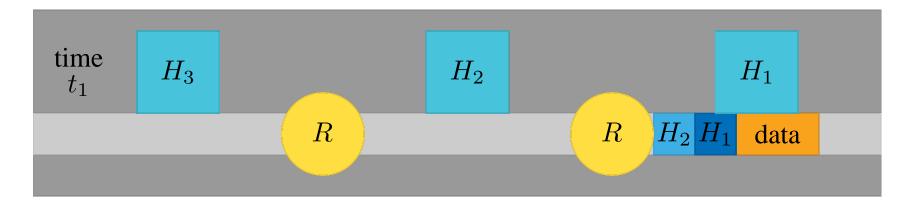
### Homogeneous Internetworking: Data Link Layer



#### **Homogeneous Internetworking: Physical Layer**

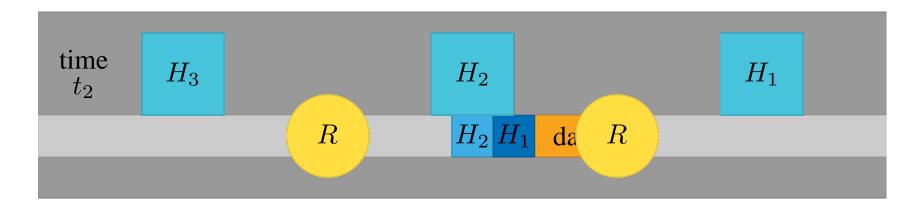


# Repeater (1)

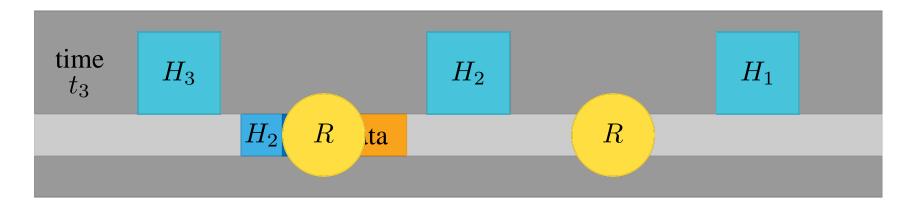


• 'repeats' frames bit by bits; no processing of frames or routing

# Repeater (2)

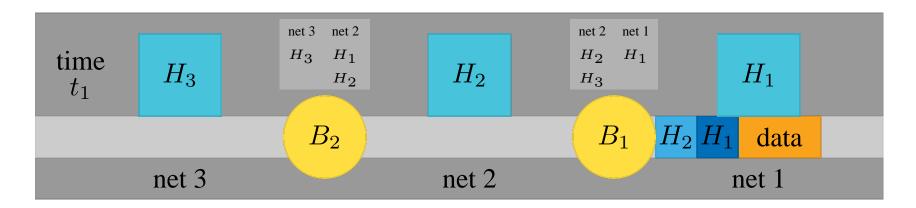


# Repeater (3)



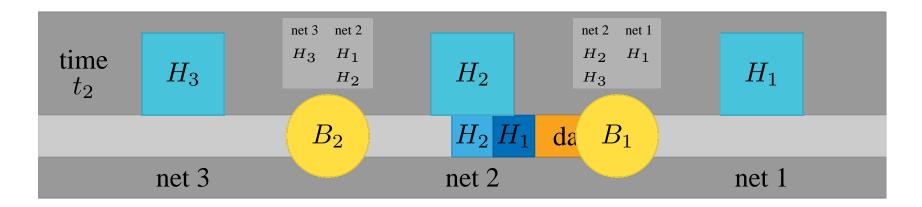
also called **hub** 

# Bridge (1)

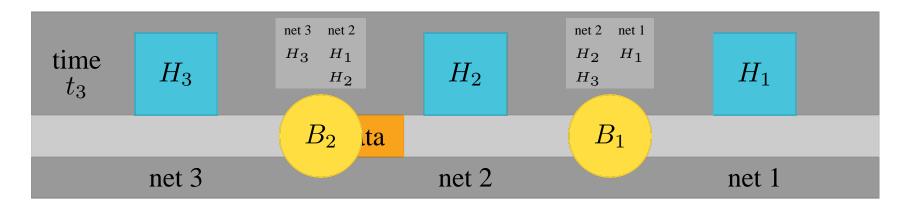


copies frames as required; store-and-forward

# Bridge (2)

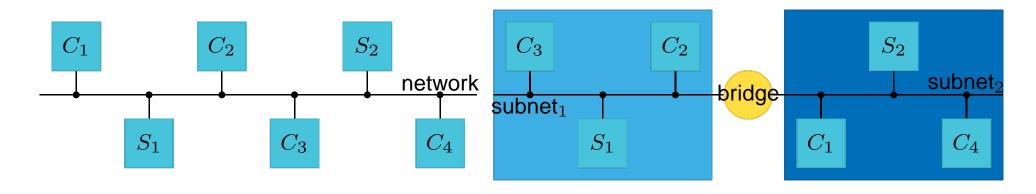


# Bridge (3)



also called switched hub

#### **Worksheet: Bridges Reduce Network Load**



	$C_1$	$C_2$	$C_3$	$C_4$
$S_1$	5	22	22	5
$S_2$	17	5	5	19

Total load= 5+22+22+5+17+5+5+19=100%

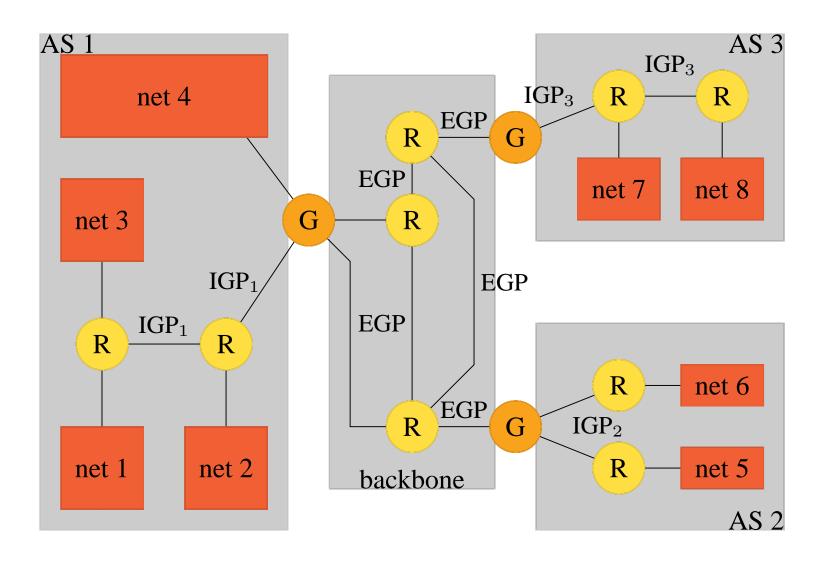
Subnet 1=Total load-Subnet 2=100-(17+19)=64%

Subnet 2=Total load-Subnet 1=100-(22+22)=56%

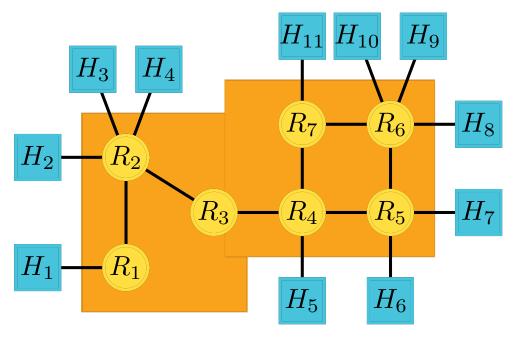
# **Bridges and Network Load**

		То					
		$H_1$	$H_2$	$H_3$	$H_4$	$H_5$	$H_6$
From	$H_1$	0	300	10	0	20	0
	$H_2$	300	0	10	0	100	10
	$H_3$	10	10	0	200	5	95
	$H_4$	10	5	300	0	5	15
	$H_5$	60	815	15	5	0	5
	$H_6$	0	15	500	5	15	0

#### **Using Routing Algorithms & Internetworking**



### **Routing Algorithm: Objectives**

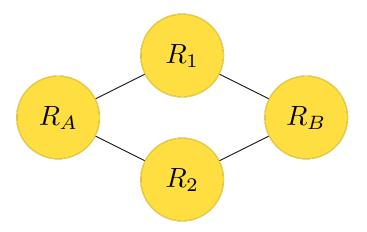


correctness it finds a valid route

robustness	should some part of the network fail, it still finds a route		
simplicity	must be able to analyse		
stability	no 'stampedes' down a link		
fairness	each host gets a share if network busy		
optimality	route minimises the network resources used		

### **Costing a route**

Often there will be a choice of how to get from A to B



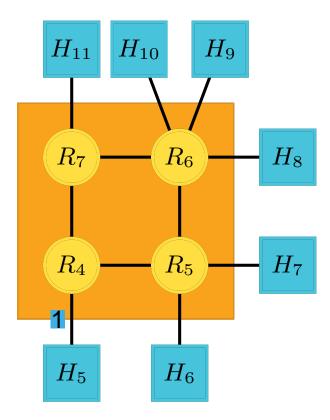
bandwidth avoid links with less free bandwidth

delay	take links which take shorter time		
price	use links with less charge per byte		
priority	prefer some data over a particular link		

### **Two Basic Approaches**

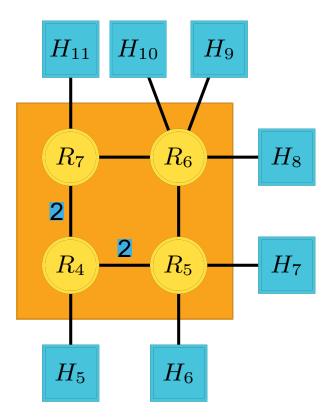
- **adaptive routing** has each router adapting its routing according to network load.
  - **isolated adaptive routing** makes decisions purely on information about the local node
  - distributed adaptive routing pass information between the nodes
- nonadaptive routing
  - routers always taking the same action to route packets
  - aka static routing

# Flood Routing (1)

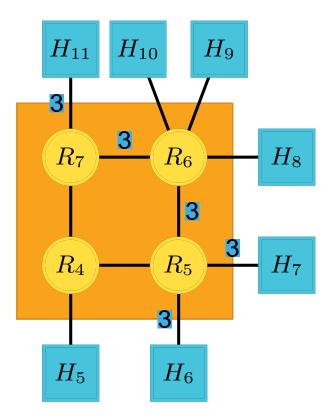


■ Send the data down all links except that it arrived on

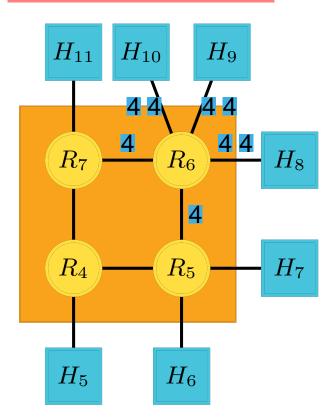
# Flood Routing (2)



# Flood Routing (3)

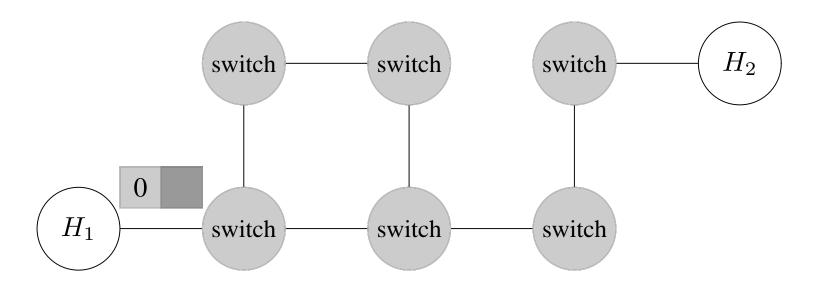


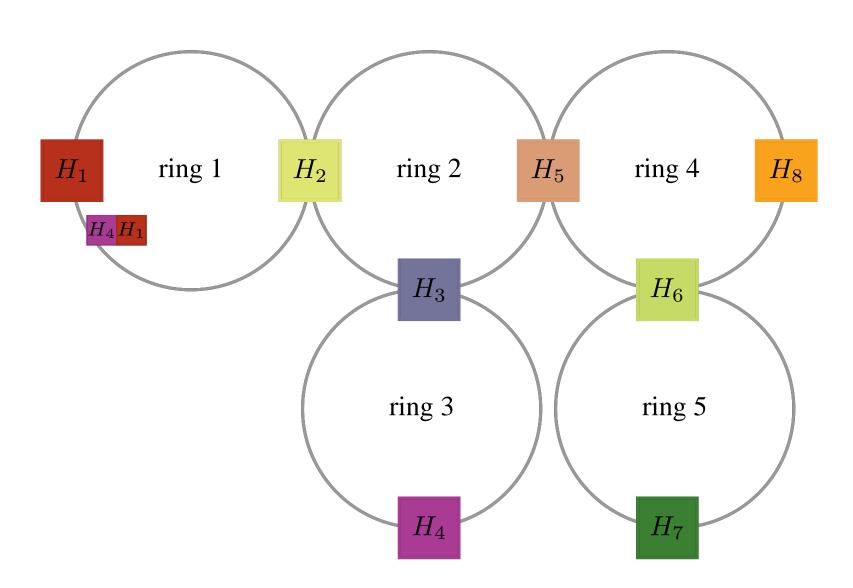
## Flood Routing (4)

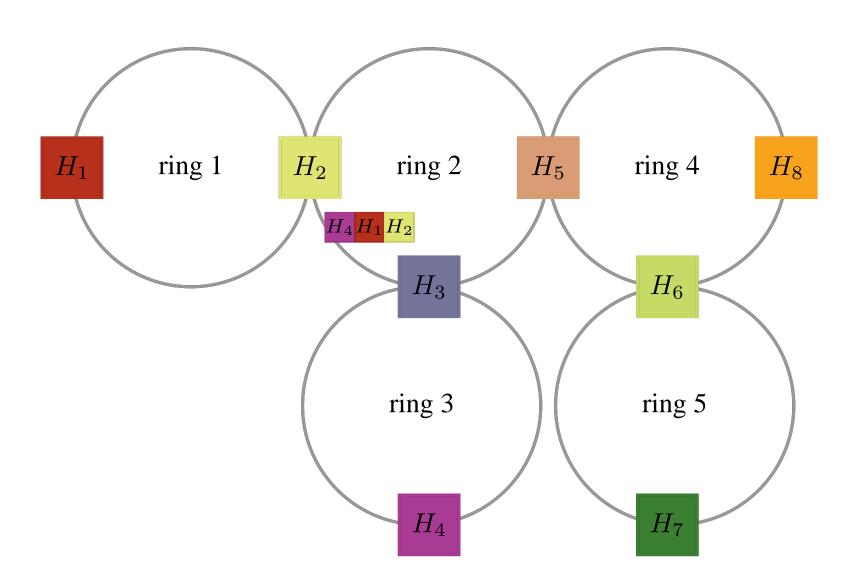


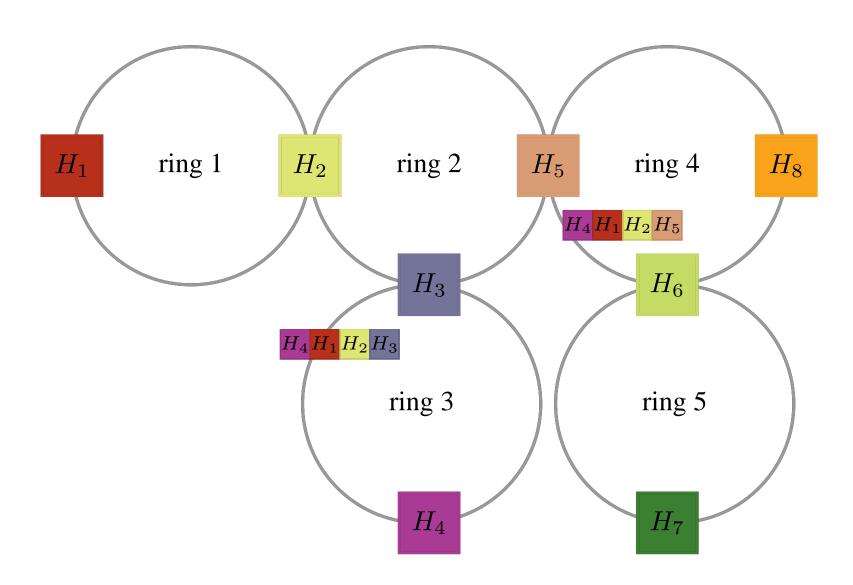
- Avoid loops by
  - recording hop count
  - recording where packet has been

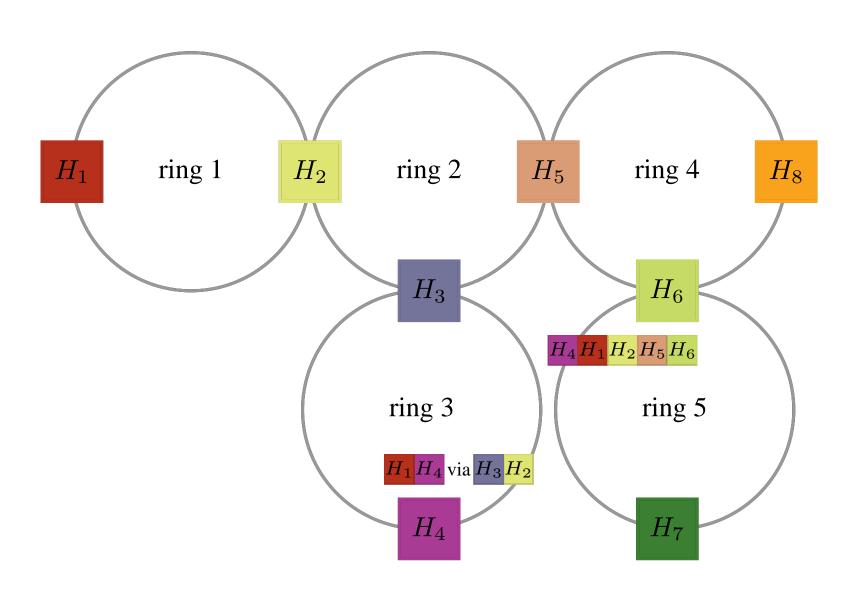
# **Worksheet: Flood Routing**



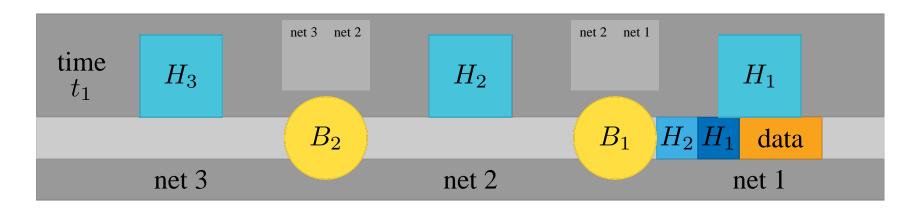






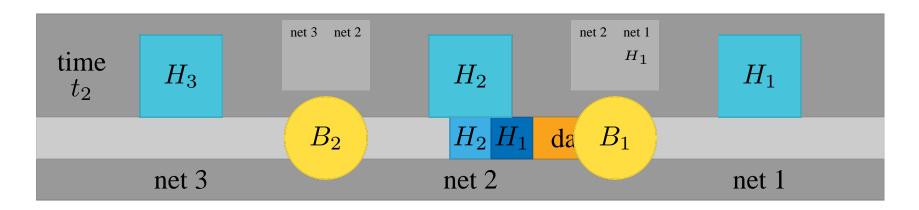


### **Backwards Learning and Transparent Bridges (1)**



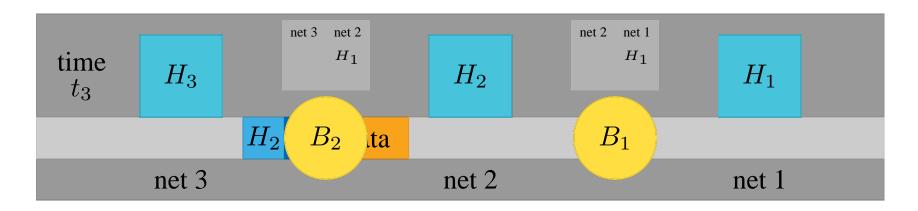
- if don't know where host is, flood network
- bridges learn where you are

#### **Backwards Learning and Transparent Bridges (2)**



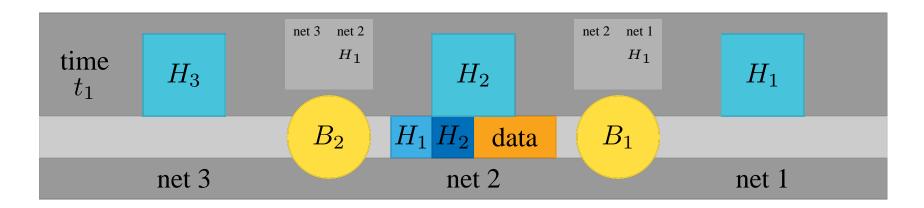
- if don't know where host is, flood network
- bridges learn where you are

#### **Backwards Learning and Transparent Bridges (3)**



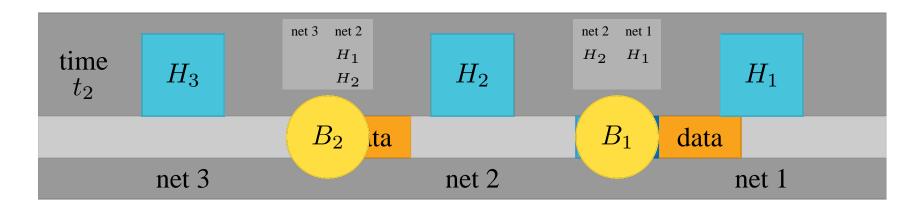
- if don't know where host is, flood network
- bridges learn where you are

#### **Backwards Learning and Transparent Bridges (4)**



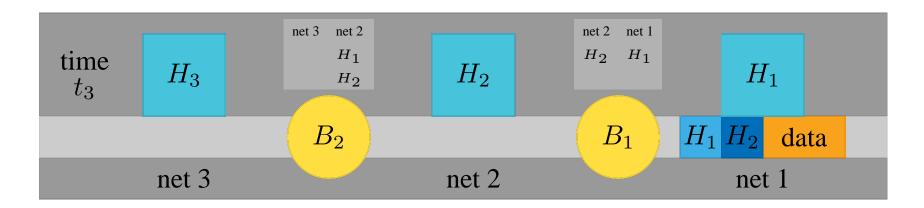
reply teaches bridges where distination was

### **Backwards Learning and Transparent Bridges (5)**



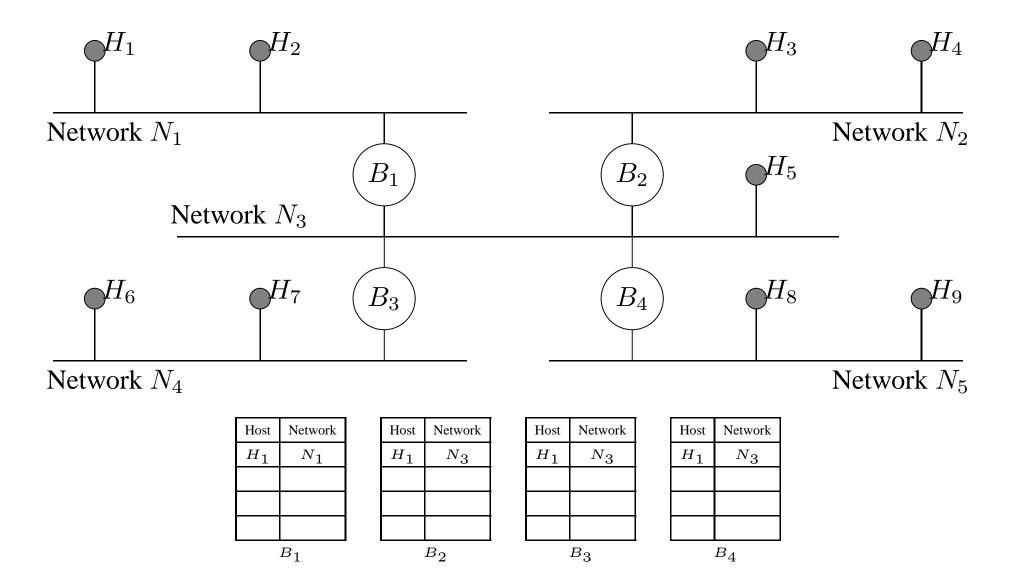
reply teaches bridges where distination was

### **Backwards Learning and Transparent Bridges (6)**

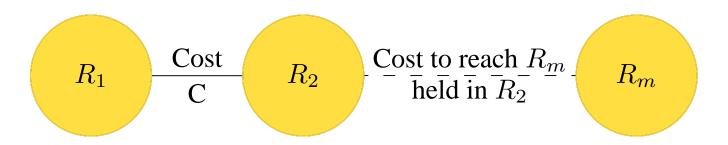


reply teaches bridges where distination was

### **Worksheet: Backwards Learning**

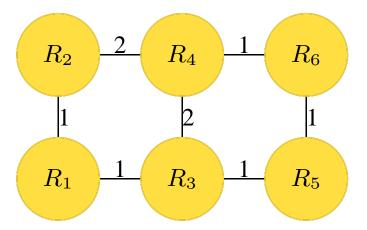


#### **Distance Vector**



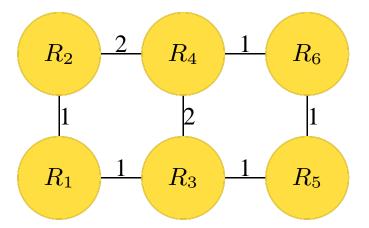
- Cost C=cost to neighbour+value in neighbour's distance vector
- Recompute local distance vector as distribute to neighbours

# **Distance Vector (1)**



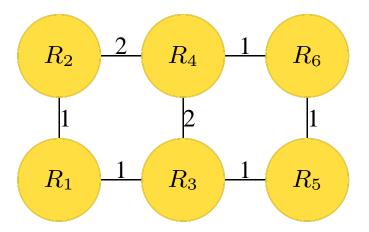
				T	o						То						
		$R_1$	$R_2$	$R_3$	$R_4$	$R_5$	$R_6$				$R_1$	$R_2$	$R_3$	$R_4$	$R_5$	$R_6$	
	$R_1$	0	0	0	0	0	0			$R_1$	0	1	1	1	1	1	
	$R_2$	0	0	0	0	0	0	\_		$R_2$	1	0	1	1	1	1	
From	$R_3$	0	0	0	0	0	0	<b>L</b> /	From	$R_3$	1	1	0	1	1	1	
	$R_4$	0	0	0	0	0	0			$R_4$	1	1	1	0	1	1	
	$R_5$	0	0	0	0	0	0			$R_5$	1	1	1	1	0	1	
	$R_6$	0	0	0	0	0	0			$R_6$	1	1	1	1	1	0	

# **Distance Vector (2)**



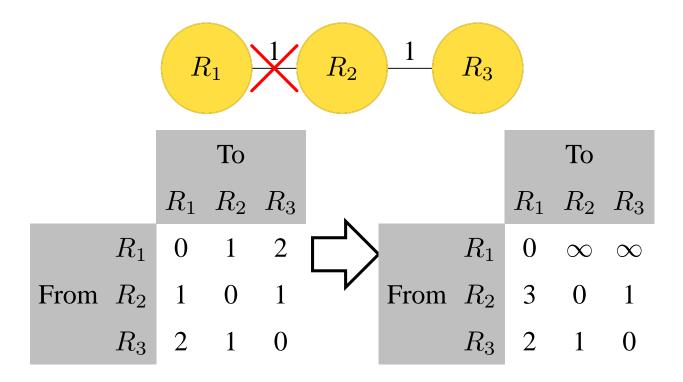
				T	o'						То						
		$R_1$	$R_2$	$R_3$	$R_4$	$R_5$	$R_6$				$R_1$	$R_2$	$R_3$	$R_4$	$R_5$	$R_6$	
	$R_1$	0	1	1	1	1	1		From	$R_1$	0	1	1	2	2	2	
	$R_2$	1	0	1	1	1	1			$R_2$	1	0	2	2	2	2	
From	$R_3$	1	1	0	1	1	1			$R_3$	1	2	0	2	1	2	
	$R_4$	1	1	1	0	1	1			$R_4$	2	2	2	0	2	1	
	$R_5$	1	1	1	1	0	1			$R_5$	2	2	1	2	0	1	
	$R_6$	1	1	1	1	1	0			$R_6$	2	2	2	1	1	0	

# **Distance Vector (3)**

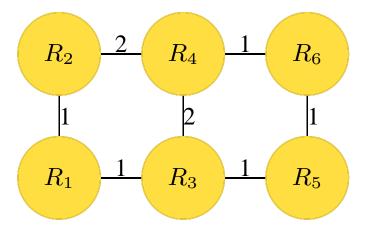


				T	o						То						
		$R_1$	$R_2$	$R_3$	$R_4$	$R_5$	$R_6$				$R_1$	$R_2$	$R_3$	$R_4$	$R_5$	$R_6$	
	$R_1$	0	1	1	2	2	2	F	From	$R_1$	0	1	1	3	2	3	
	$R_2$	1	0	2	2	2	2			$R_2$	1	0	2	2	3	3	
From	$R_3$	1	2	0	2	1	2			$R_3$	1	2	0	2	1	2	
	$R_4$	2	2	2	0	2	1			$R_4$	3	2	2	0	2	1	
	$R_5$	2	2	1	2	0	1			$R_5$	2	3	1	2	0	1	
	$R_6$	2	2	2	1	1	0			$R_6$	3	3	2	1	1	0	

#### **Distance Vector: Poor Response to Failures**



# Link State



				T	o'						То						
		$R_1$	$R_2$	$R_3$	$R_4$	$R_5$	$R_6$				$R_1$	$R_2$	$R_3$	$R_4$	$R_5$	$R_6$	
From	$R_1$	0	1	1	?	?	?	Fro		$R_1$	0	1	1	3	2	3	
	$R_2$	1	0	?	2	?	?			$R_2$	1	0	2	2	3	3	
	$R_3$	1	?	0	2	1	?		From	$R_3$	1	2	0	2	1	2	
	$R_4$	?	2	2	0	?	1			$R_4$	3	2	2	0	2	1	
	$R_5$	?	?	1	?	0	1			$R_5$	2	3	1	2	0	1	
	$R_6$	?	?	?	1	1	0			$R_6$	3	3	2	1	1	0	