

# Rational Agents

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Introduction to Artificial Intelligence  
2nd Part

# What you have seen

You have seen procedures for computational problem-solving:

- searching
- learning
- planning

# What we will be looking at

An **agent**, a mathematical entity acting in a simple world

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  - Imperfect and incomplete information
  - Quantifying uncertainty, attaching probabilities
  - Going for uncertain outcomes, calculating expected utility

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- Able to take decisions under uncertainty
  - Imperfect and incomplete information
  - Quantifying uncertainty, attaching probabilities
  - Going for uncertain outcomes, calculating expected utility
- Able to update his (or her) beliefs when confronted with new information (learning)

# What is rationality?



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Nobel Prize Winner  
Economics

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Agents (not only humans) *can* be rational!



# The lectures one by one

- Logical Agents I
- Logical Agents II
- An Uncertain World
- Making Sense of Uncertainty
- Making (Good) Decisions
- Making Good Decisions in time
- Learning from Experience I
- Learning from Experience II

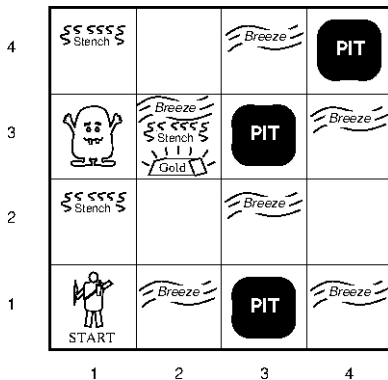
# Logical Agents I

# The main reference



Stuart Russell and Peter Norvig  
Artificial Intelligence: a modern approach  
Chapters 7-9

# The Wumpus World



# Agents

**Sensors** Breeze, Glitter, Smell

**Actuators** Turn L/R, Go, Grab, Release, Shoot, Climb

**Rewards** 1000 escaping with gold, -1000 dying, -10 using arrow, -1 walking

- Environment**
- Squares adjacent to Wumpus are smelly
  - Squares adjacent to pit are breezy
  - Glitter iff gold is in the same square
  - Shooting kills Wumpus if you are facing it
  - Shooting uses up the only arrow
  - Grabbing picks up gold if in same square
  - Releasing drops the gold in same square

# Knowledge base

- A set of sentences representing what the agent thinks about the world.
  - "I am in [2,1]"
  - "I am out of arrows"
  - "I smell Wumpus"
  - "I'd better not go forward"
- We interpret it as what the agent *knows*, but it can very well work for what the agent *believes*.

# Updating the knowledge base

- What we TELL the knowledge base
- What we ASK the knowledge base

```
function KB-AGENT( percept ) returns an action  
  static: KB, a knowledge base  
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  action ← ASK(KB, MAKE-ACTION-QUERY(t))  
  TELL(KB, MAKE-ACTION-SENTENCE(action, t))  
  t ← t + 1  
  return action
```

# Rational explorations

- The starting state...

OK <input type="checkbox"/> A			



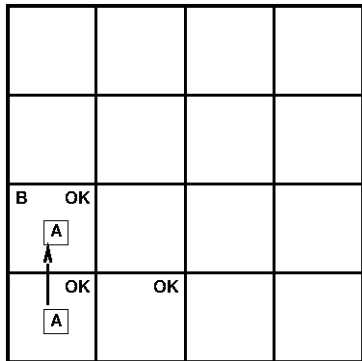
# Rational explorations

- and what we know.

OK			
OK <span style="border: 1px solid black; padding: 2px;">A</span>	OK		

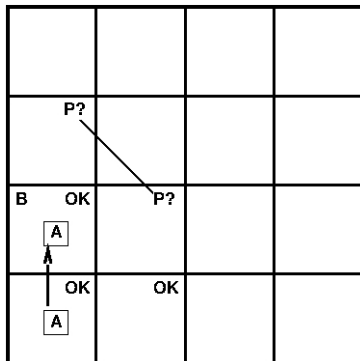
# Rational explorations

- B stands for Breeze



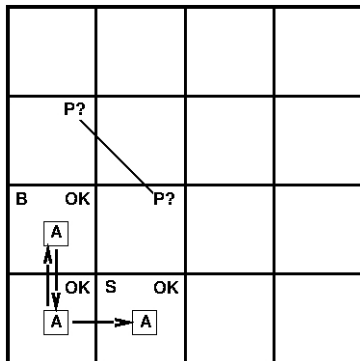
# Rational explorations

- Where is the pit?
- We are ruling out one square!



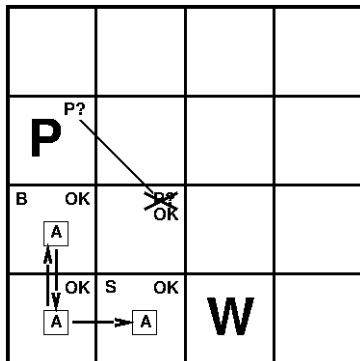
# Rational explorations

- S stands for smell
- What do we know?



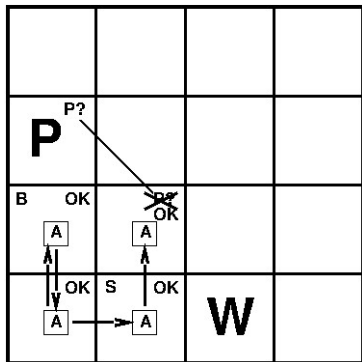
# Rational explorations

- Logic is the key



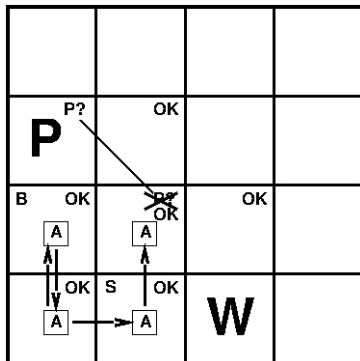
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- The further we go the more we know



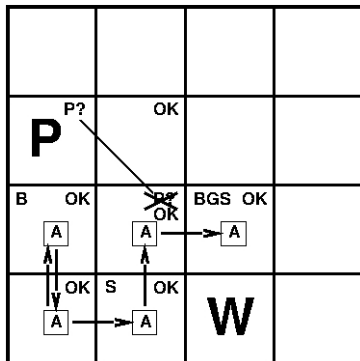
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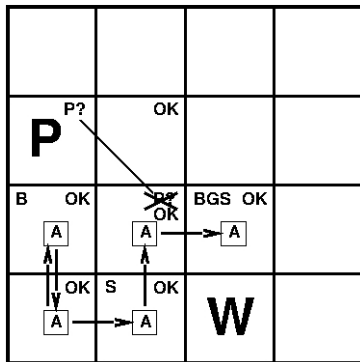
- Gold!





# Rational explorations

- We know the way out
- Game over



## Reasoning in the Wumpus World

Let  $P_{i,j}$  be true if there is a pit in  $[i,j]$ .

Let  $B_{i,j}$  be true if there is a breeze in  $[i,j]$ .

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“Pits cause breezes in adjacent squares”

$$B_{1,1} \Leftrightarrow (P_{1,2} \vee P_{2,1})$$

$$B_{2,1} \Leftrightarrow (P_{1,1} \vee P_{2,2} \vee P_{3,1})$$

“A square is breezy **if and only if** there is an adjacent pit”

# Expressivity: at what cost?

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Propositional Logic lacks expressive power

# First order logic

- Massive increase of expressivity
- But there are costs, e.g., decidability
- We will see how to exploit the gains while limiting the costs



# FOL KB

- We can encode the KB at each particular time point using FOL

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 $\forall t \text{ } AtGold(t) \wedge \neg Holding(Gold, t) \Rightarrow Action(Grab, t)$

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 $\forall t \text{ } AtGold(t) \wedge \neg Holding(Gold, t) \Rightarrow Action(Grab, t)$
  - Axioms from knowledge to knowledge, e.g.,  
 $\forall t \text{ } AtGold(t) \wedge Action(Grab, t) \Rightarrow Holding(Gold, t + 1)$

# Describing the world

Perception  $\forall s, b, t \text{ Percept}([s, b, \textit{Glitter}], t) \Rightarrow \textit{AtGold}(t)$



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Decision-making  $\forall t \textit{AtGold}(t) \Rightarrow \textit{Action}(\textit{Grab}, t)$

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Decision-making  $\forall t \textit{AtGold}(t) \Rightarrow \textit{Action}(\textit{Grab}, t)$

Internal reflection  $\forall t \textit{AtGold}(t) \wedge$   
 $\neg \textit{Holding}(\textit{Gold}, t) \Rightarrow \textit{Action}(\textit{Grab}, t)$ , do we have  
 gold already? (notice we cannot observe if we are  
 holding gold, we need to track it)

# Describing the world

## Adjacent squares

$$\begin{aligned} \forall x, y, a, b \text{ Adjacent}([x, y], [a, b]) &\Leftrightarrow \\ (x = a \wedge (y = b - 1 \vee y = b + 1)) \vee \\ (y = b \wedge (x = a - 1 \vee x = a + 1)) \end{aligned}$$

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“A square is breezy **if and only if** there is an adjacent pit”

$$\forall s, \text{Breezy}(s) \Leftrightarrow \exists r (\text{Adjacent}(r, s) \wedge \text{Pit}(r))$$

# Describing the world

- We can go on and describe plans, causal rules, etc.
- But let's do some reasoning now

# Facts and Knowledge Bases



"Richard the Lionheart is a king"

# Facts and Knowledge Bases



"Joffrey Baratheon is a king"



# Telling and Asking

*Tell(KB, King(Joffrey))*

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*Ask*(KB,  $\exists x \text{ Person}(x)$ ) is there a person?

*Askvar*(KB, *Person*( $x$ )) who is a person?

*Askvar* returns a list of **substitutions**:  $\{x/\text{Joffrey}\}, \{x/\text{Jaime}\}$

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$$\sigma = \{x/\text{Tyrion}, y/\text{Joffrey}\}$$

$$S\sigma = \text{Smarter}(\text{Tyrion}, \text{Joffrey})$$

$\text{Askvar}(KB, S)$  returns some/all  $\sigma$  such that  $KB \models S\sigma$

# Unification

$$\forall x \text{ King}(x) \wedge \text{Greedy}(x) \Rightarrow \text{Evil}(x)$$

*King(Joffrey)*

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$$\theta = \{x/\text{Joffrey}, y/\text{Joffrey}\} \text{ works}$$

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$Knows(Joffrey, x)$	$Knows(Joffrey, Sansa)$	
$Knows(Joffrey, x)$	$Knows(y, Sansa)$	
$Knows(Joffrey, x)$	$Knows(y, Mother(Joffrey))$	
$Knows(Joffrey, x)$	$Knows(x, Sansa)$	

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$Knows(Joffrey, x)$	$Knows(y, Mother(Joffrey))$	$\{y/Joffrey, x/Mother(Joffrey)\}$
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$Knows(Joffrey, x)$	$Knows(y, Mother(Joffrey))$	$\{y/Joffrey, x/Mother(Joffrey)\}$
$Knows(Joffrey, x)$	$Knows(x, Sansa)$	fail

# Standardising apart

*Knows(Joffrey, x) & Knows(x, Sansa)* fails

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*Knows(Joffrey, x) & Knows(x, Sansa)* fails

Standardising apart eliminates overlap of variables, e.g.,  
*Knows(z<sub>17</sub>, Sansa)*

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Definite clause:

disjunction of literals, **exactly** one of which positive



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$$\frac{p_1', p_2', \dots, p_n', (p_1 \wedge p_2 \wedge \dots \wedge p_n \Rightarrow q)}{q\theta}$$

where  $p_i'\theta = p_i\theta$  for all  $i$

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Assuming all variables are universally quantified...

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Assuming all variables are universally quantified...

$p_1'$  is *King(Joffrey)*

$p_1$  is *King(x)*

$p_2'$  is *Greedy(y)*

$p_2$  is *Greedy(x)*

$\theta$  is  $\{x/Joffrey, y/Joffrey\}$

$q$  is *Evil(x)*

$q\theta$  is *Evil(Joffrey)*

# Soundness of GMP

Need to show that

$$p_1', \dots, p_n', (p_1 \wedge \dots \wedge p_n \Rightarrow q) \models q\theta$$

provided that  $p_i'\theta = p_i\theta$  for all  $i$

Lemma: If  $\varphi$  is definite clause, then  $\varphi \models \varphi\theta$  by Universal Instantiation.

- 1  $(p_1 \wedge \dots \wedge p_n \Rightarrow q) \models (p_1 \wedge \dots \wedge p_n \Rightarrow q)\theta = (p_1\theta \wedge \dots \wedge p_n\theta \Rightarrow q\theta)$
- 2  $p_1', \dots, p_n' \models p_1' \wedge \dots \wedge p_n' \models p_1'\theta \wedge \dots \wedge p_n'\theta$
- 3 From 1 and 2,  $q\theta$  follows by ordinary Modus Ponens

# Coming next

- Making sound and efficient inferences
- Where to start?
- How to go on?