## Explicit Connection Actions in Multiparty Session Types

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#### Outline

- Background: multiparty session types (MPST)
  - Scribble: ongoing work on implementing and applying MPST to practice
  - × Standard MPST do not support sessions with *dynamic or optional involvement of participants*

#### MPST with explicit connection actions

- MP sessions as a dynamically evolving configuration of binary connections
  - Modelling-based well-formedness for MPST protocols
  - Session subtyping and role progress
  - Multiparty correlation of binary connections
- Motivating examples
  - Web services choreography (Travel Agency)
  - Microservices industry use case (Supplier Info)
  - Standardised application-layer protocol (FTP)

#### Standard presentation: three-layer framework

Global type

 $G = A \rightarrow B : \langle U_1 \rangle . B \rightarrow C : \langle U_2 \rangle . C \rightarrow A : \langle U_3 \rangle$ 

Global description of multiparty *message passing* protocol/choreography Participants abstracted as *roles* 

Local types

 $T_A = \langle B, U_1 \rangle . \langle C, U_3 \rangle$ 

Localised view of the protocol for each role

Endpoint processes

 $P_A = a[A](x) \cdot x! \langle B, u_1 \rangle \cdot x? (C, y)$ 

Perform I/O via special primitives on channels

# [G] Projection $T_{Alice}$ $T_{Bob}$ $T_{Carol}$ $T_{Carol}$ $T_{Carol}$ $P_{Carol}$ $T_{Carol}$ $T_{Carol}$

#### Communication safety is ensured for a parallel composition of well-typed endpoints

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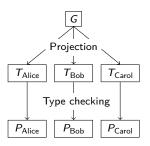
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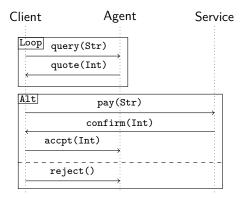
Perform I/O via special primitives on channels

 Communication safety is ensured for a parallel composition of well-typed endpoints



#### Web Services use case: Travel Agency

W3C Choreography working group requirements use case https://www.w3.org/TR/ws-chor-reqs/#UC-001



- [ECOOP06] Session Types for Object-Oriented Languages. Dezani-Ciancaglini, Mostrous, Yoshida and Drossopoulou. "Buyer-Seller-Shipper"
- [CONCUR08] Global progress in dynamically interleaved multiparty sessions. Bettini, Coppo, D'Antoni, Luca, Dezani-Ciancaglini and Yoshida. "Three-Buyer"
  - [FTPL16] Behavioral Types in Programming Languages. Ancona et al. "Customer-Agency"

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global protocol Travel(role C, role A, role S) {
    choice at C {
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        quote(Int) from A to C;
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- imes Not a valid global type
  - "Unfinished role" error
  - Ruled out by syntactic well-formedness:
    - Each involved participant must be present in all choice cases

▶ A session cannot have dynamic or optional involvement of participants

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Session initation by a global atomic synchronisation

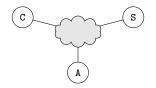
Sender-asynchronous (non-blocking output), reliable, role-to-role ordering



▶ MPST safety: run-time session execution is safe from

- Reception errors
- Deadlocks
- Orphan messages

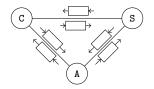
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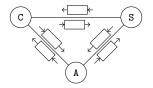
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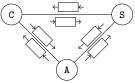
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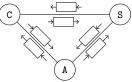


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 Practical protocol specifications include explicit connection request/accept (and disconnect) actions

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explicit global protocol Travel(role C, role A, role S) {
   connect C to A;
   do Main(C. A. S):
}
aux global protocol Main(role C, role A, role S) {
   choice at C {
       query(Str) from C to A;
       quote(Int) from A to C;
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   } or {
       connect C to S;
                                              С
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```

S

- Dynamically established binary connections
  - Role involvement guarded by initial connection accept
  - imes Cannot apply standard (conservative) syntactic MPST well-formedness
- Previous works have studied MPST safety in terms of CFSM-based well-formedness conditions (*multiparty compatibility*)

[ICALP13] Multiparty Compatibility in Communicating Automata. Deniélou and Yoshida. [CONCUR15] Meeting Deadlines Together. Bocchi, Lange and Yoshida.

- Our approach: MPST protocol validation by a combination of syntactic constraints and explicit error checking
  - Adapt basic MPST syntactic constraints to our extended setting...
  - ...that ensure soundness of checking a 1-bounded model of the protocol

#### Syntactic constraints

- MPST error checking
- Global type grammar
- Role enabling
- Consistent external choices

```
choice at A {
   1() from A to B;
   1() connect A to C;
   ...
} or {
   2() from A to B;
   2() connect A to C;
   ...
}
   Globally-paired interactions
```

Deterministic choices

#### Syntactic constraints

- MPST error checking
- Global type grammar

#### Role enabling

Consistent external choices

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choice at A {
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    //1() from A to C;
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} or {
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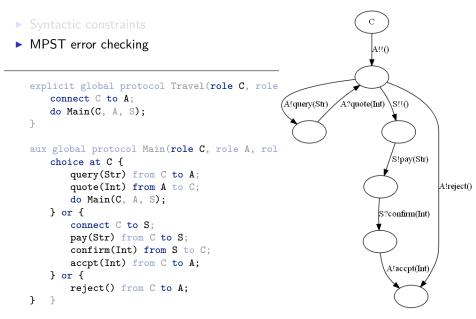
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                                                            Projection
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                                                       T_{C}
                                                                         T_{\rm S}
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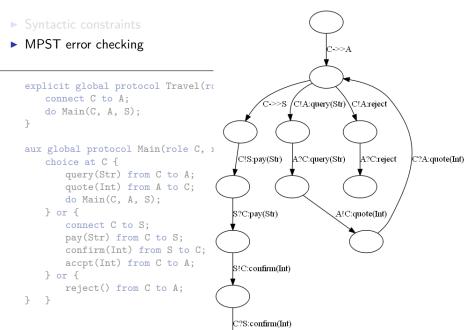


# Validating MPST with explicit connections

- Syntactic constraints
- MPST error checking

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                                                              S
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}
                                                               C??0
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   choice at C {
       query(Str) from C to A;
       quote(Int) from A to C;
                                                               C?pay(Str)
       do Main(C, A, S);
   f or f
       connect C to S:
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                                                               C!confirm(Int)
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## Validating MPST with explicit connections



# Validating MPST with explicit connections

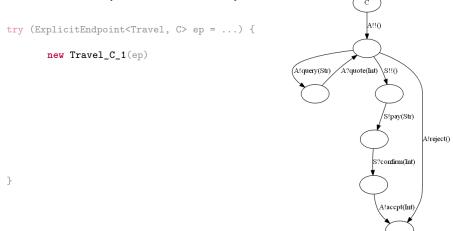
- Syntactic constraints
- MPST error checking
- MPST safety
  - Reception errors, Orphan messages
  - Unfinished roles, Connection/Disconnect/Unconnected errors

```
explicit global protocol Foo(role A, role B) {
    connect A to B; ... disconnect A and B; do Foo(A, B); }
```

- MPST progress
  - Eventual Reception, Role progress, Eventual Connection
- Soundness of 1-bounded MPST validation
   Let S<sub>0</sub> be the initial session of a wf(G) that is 1-safe and satisfies 1-progress.
   Then S<sub>0</sub> is safe and satisfies progress.

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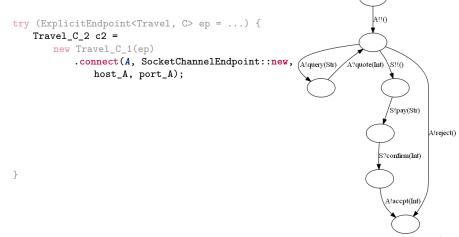
- Validated MPST used to generate Java endpoint APIs
  - Generated APIs promote hybrid approach to session safety [FASE16]
    - Endpoint FSM structures captured as statically-typed call-chaining APIs
    - Usage contract of API is to use every "state channel" instance exactly once
    - Enforced by run-time channel linearity checks



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A!accpt(Inf

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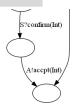
- connect(S role, Callable<? extends BinaryChannelEndpoint> cons, String host, int port) : Travel\_C\_4 Travel\_C\_2
- send(A role, reject op) : EndSocket Travel\_C\_2
- send(A role, query op, String arg0) : Travel\_C\_3 Travel\_C\_2



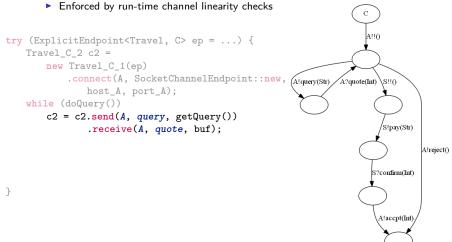
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receive(A role, quote op, Buf<? super Integer> arg1) : Travel\_C\_2 - Travel\_C\_3



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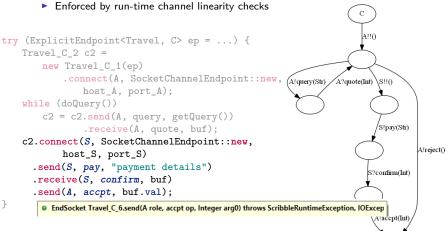


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    - Usage contract of API is to use every "state channel" instance exactly once
  - Enforced by run-time channel linearity checks A!!0 try (ExplicitEndpoint<Travel, C> ep = ...) {  $Travel_C_2 c2 =$ new Travel C 1(ep) .connect(A, SocketChannelEndpoint::new, (A!query(Str)) A?quote(Int) \s!!0 host A. port A): while (doQuery()) c2 = c2.send(A, query, getQuery()) .receive(A, quote, buf); S!pay(Str) c2.connect(S, SocketChannelEndpoint::new, host\_S, port\_S) .send(S, pay, "payment details") S?confirm(Int) .receive(S, confirm, buf) .send(A, accpt, buf.val);

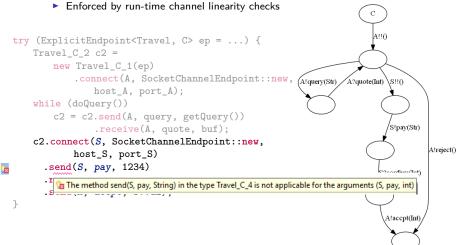
A!reject()

A!accpt(Int)

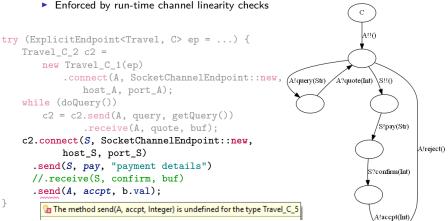
- Validated MPST used to generate Java endpoint APIs
  - Generated APIs promote hybrid approach to session safety [FASE16]
    - Endpoint FSM structures captured as statically-typed call-chaining APIs
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- Validated MPST used to generate Java endpoint APIs
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### "Naive" session subtyping

- > The first naive attempt at Travel Agency is invalid
  - (Non-explicit protocols checked by assuming all roles pre-connected)

```
global protocol Travel(role C, role A, role S) {
    choice at C {
        query(Str) from C to A;
        quote(Int) from A to C;
        do Travel(C, A, S);
    } or {
        pay(Str) from C to S;
        confirm(Int) from S to C;
        accpt(Int) from C to A;
    } or {
        reject from C to A;
    }
}
```

- $\times\,$  s is "unfinished" in the <code>reject</code> case
- × Role progress may also be violated for s in the query case (without assuming some notion of fairness)

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        do Travel(C, A, S);
    } or {
        pay(Str) from C to S;
        confirm(Int) from S to C;
        accpt(Int) from C to A;
    } or {
        reject from C to A;
    }
}
```

imes s is "unfinished" in the reject case

 × Role progress may also be violated for s in the query case (without assuming some notion of fairness)

```
    Related to session subtyping
```

```
global protocol Foo(role A, role B, role C) {
   choice at A { 1() from A to B; 1() from A to C; }
        or { 2() from A to B; 2() from A to C; }
   do Foo(A, B, C);
}
```

[ACTA05] Subtyping for session types in the pi calculus. Gay and Hole. [MSCS16] Fair subtyping for multiparty sessions. Padovani.

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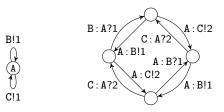
```
global protocol Foo(role A, role B, role C) {
   choice at A { 1() from A to B; }
        or { 2() from A to C; }
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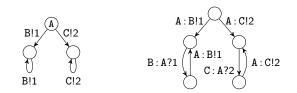
We implement two basic views:

- Fair output choices (as modelled so far)
- "Most unfair" while still session type safe
  - Endpoints commit to a single case in any output choice (Extreme "naive" output choice subtyping)
  - Modelled by a transformation on endpoint FSMs

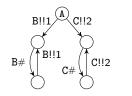
```
global protocol Foo(role A, role B, role C) {
   choice at A { 1() from A to B; }
        or { 2() from A to C; }
   do Foo(A, B, C);
}
```



```
global protocol Foo(role A, role B, role C) {
   choice at A { 1() from A to B; }
        or { 2() from A to C; }
   do Foo(A, B, C);
}
```



```
explicit global protocol Foo(role A, role B, role C) {
    choice at A { 1() connect A to B; disconnect A and B; }
        or { 2() connect A to C; disconnect A and C; }
    do Foo(A, B, C);
}
```





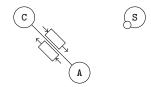
OK if fairness assumed

```
global protocol Foo(role A, role B, role C) {
   choice at A {
      1() from A to B;
      do Foo(A, B, C);
   } or {
      2() from A to B;
      2() from B to C;
   }
}
```

- Modelling based on a single session with one endpoint process per role
  - Connection mechanism (in particular, addressing) left abstract
- In practice: correlation by session identifier tags, port coordination, ... connect A to B; .. connect A to C; .. connect B to C; ...
- Travel Agency (accpt case) with dynamic port forwarding

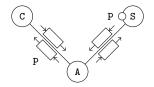
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```
accpt() from C to A;
connect A to S;
port(Int) from S to A;
port(Int) from A to C;
connect C to S;
pay(Str) from C to S;
confirm(Int) from S to C;
```



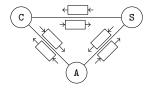
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confirm(Int) from S to C;
...
```



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```
accpt() from C to A;
connect A to S;
port(p:Int) from S to A; @"open=p:C"
port(p) from A to C;
connect C to S; @"port=p"
pay(Str) from C to S;
confirm(Int) from S to C;
...
```

[RV13] Practical interruptible conversations. Hu, Neykova, Yoshida, Demangeon and Honda.

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```
accpt() from C to A;
connect A to S;
port(p:Int) from S to A; @"open=p:C"
//port(p) from A to C;
connect C to S; @"port=p" X
pay(Str) from C to S;
confirm(Int) from S to C;
...
```

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accpt() from C to A;
connect A to S;
port(p:Int) from S to A; @"open=p:C"
port(p) from A to C;
connect C to S; @"port=p"
pay(Str) from C to S;
confirm(Int) from S to C;
...
s_C.receive(A, port).connect(S, ..., host_S, pay, "payment details")...
s_A.receive(S, port).send(C, port)...
s_S.send(A, port).accept(C, pay, b)...
```

- Modelling based on a single session with one endpoint process per role
  - Connection mechanism (in particular, addressing) left abstract
- In practice: correlation by session identifier tags, port coordination, ... connect A to B; .. connect A to C; ... connect B to C; ...
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accpt() from C to A;
connect A to S;
port(p:Int) from S to A; @"open=p:C"
port(p) from A to C;
connect C to S; @"port=p"
pay(Str) from C to S;
confirm(Int) from S to C;
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s_C.receive(A, port).connect(S, ..., host_S, pay, "payment details")...
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accpt() from C to A;
connect A to S;
port(p:Int) from S to A; @"open=p:C"
port(p) from A to C;
connect C to S; @"port=p"
pay(Str) from C to S;
confirm(Int) from S to C;
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s_C.receive(A, port).connect(S, ..., host_S, pay, "payment details")...
s_A.receive(S, port).send(C, port)...
s_S.send(A, port).accept(C, pay, b)...
```

# FTP (active/passive modes)

```
220 from S to C;
USER from C to S;
choice at S {
 331 from S to C;
 PASS from C to S;
 choice at S {
   230 from S to C:
   choice at C {
     PASV from C to S; // Passive mode
     choice at S {
       227(p:Int) from S to C; @"open=p:C"
       . . .
     } or {
     }
   f or f
     PORT(q:Int) from C to S; @"open=q:S" // Active mode
     choice at S {
       200 from S to C;
       . . .
     7
   } or {
1 1 1
```

#### Microservices use case

```
explicit global protocol InfoAuth
     (role LoginSvc, role Client, role AuthSvc,
      role Filtersvc, role SupplierSvc,
      role ContractSvc) {
  connect Client to LoginSvc;
 login(UserName, password) from Client to LoginSvc;
  choice at LoginSvc {
   loginfailure() from LoginSvc to Client;
  } or f
   loginsuccess() from LoginSvc to Client;
   disconnect Client and LoginSvc:
   connect Client to AuthSvc:
   do Main(Client, AuthSvc, Filtersvc,
       SupplierSvc, ContractSvc);
3 3
aux global protocol Main
     (role Client, role AuthSvc.
      role Filtersvc, role SupplierSvc,
      role ContractSvc) {
  choice at Client {
   getsuppliers(UUID) from Client to AuthSvc:
   do SuppInfo(Client, AuthSvc,
       Filtersvc, SupplierSvc);
  } or f
   getcontracts() from Client to AuthSvc;
   do ContractInfo(Client, AuthSvc,
       Filtersvc, ContractSvc);
 do Main(Client, AuthSvc, Filtersvc,
     SupplierSvc, ContractSvc);
```

aux global protocol SuppInfo (role Client, role AuthSvc, role Filtersvc, role SupplierSvc) { choice at AuthSvc { deny() from AuthSvc to Client; } or f connect AuthSvc to SupplierSvc; getsuppliers() from AuthSvc to SupplierSvc; suppliers() from SupplierSvc to AuthSvc; disconnect AuthSvc and SupplierSvc; do FilterInfo <Filtersuppliers(UserContext, Filters, SupplierDetails)> (AuthSvc. Filtersvc); suppliers() from AuthSvc to Client: 3

```
aux global protocol ContractInfo
     (role Client, role AuthSvc.
      role Filtersvc, role ContractSvc) {
 choice at AuthSvc {
```

```
} }
```

7

3

aux global protocol FilterInfo <sig Querv> (role AuthSvc, role Filtersvc) { Query connect AuthSyc to Filtersyc: filtered() from Filtersvc to AuthSvc; disconnect AuthSvc and Filtersvc;

3

#### Microservices use case

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      role Filtersvc, role SupplierSvc,
      role ContractSvc) {
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 login(UserName, password) from Client to LoginSvc;
  choice at LoginSvc {
   loginfailure() from LoginSvc to Client;
   loginsuccess() from LoginSvc to Client;
   disconnect Client and LoginSvc:
   connect Client to AuthSvc:
   do Main(Client, AuthSvc, Filtersvc,
       SupplierSvc, ContractSvc):
```

```
aux global protocol Main
     (role Client, role AuthSvc.
      role Filtersvc, role SupplierSvc,
      role ContractSvc) {
 choice at Client {
   getsuppliers(UUID) from Client to AuthSvc:
   do SuppInfo(Client, AuthSvc,
 } or {
   getcontracts() from Client to AuthSvc;
   do ContractInfo(Client, AuthSvc.
      Filtersvc. ContractSvc):
 do Main(Client, AuthSvc, Filtersvc,
```

(role Client, role AuthSvc, choice at AuthSvc { deny() from AuthSvc to Client; connect AuthSvc to SupplierSvc; getsuppliers() from AuthSvc to SupplierSvc; suppliers() from SupplierSvc to AuthSvc; disconnect AuthSvc and SupplierSvc; do FilterInfo Filters, SupplierDetails)> (AuthSvc, Filtersvc); suppliers() from AuthSvc to Client;

```
aux global protocol ContractInfo
     (role Client, role AuthSvc.
      role Filtersvc, role ContractSvc) {
 choice at AuthSvc {
```

```
aux global protocol FilterInfo
     <sig Querv>
     (role AuthSvc. role Filtersvc) {
  Query connect AuthSyc to Filtersyc:
 filtered() from Filtersvc to AuthSvc;
 disconnect AuthSvc and Filtersvc;
```

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```
explicit global protocol InfoAuth
     (role LoginSvc, role Client, role AuthSvc,
  connect Client to LoginSvc;
 login(UserName, password) from Client to LoginSvc;
  choice at LoginSvc {
   loginfailure() from LoginSvc to Client;
   loginsuccess() from LoginSvc to Client;
   disconnect Client and LoginSvc:
   connect Client to AuthSvc:
   do Main(Client, AuthSvc, Filtersvc,
       SupplierSvc, ContractSvc):
aux global protocol Main
     (role Client, role AuthSvc.
      role ContractSvc) {
  choice at Client {
   getsuppliers(UUID) from Client to AuthSvc:
   do SuppInfo(Client, AuthSvc,
  } or {
   getcontracts() from Client to AuthSvc;
   do ContractInfo(Client, AuthSvc.
       Filtersvc. ContractSvc):
 do Main(Client, AuthSvc, Filtersvc,
```

```
aux global protocol SuppInfo
    (role Client, role AuthSvc,
        role Filtersvc, role SupplierSvc) {
    choice at AuthSvc to SupplierSvc;
    demy() from AuthSvc to SupplierSvc;
    getsuppliers() from AuthSvc to SupplierSvc;
    disconnect AuthSvc and SupplierSvc;
    do FilterInfo
        <filters, SupplierSvc to AuthSvc;
    diters, SupplierSvc;
    suppliers() from AuthSvc to Client;
        Filters, SupplierDetails)>
        (AuthSvc, Filtersvc);
        suppliers() from AuthSvc to Client;
    }
}
aux global protocol ContractInfo
    (role Client, role AuthSvc.
```

```
role Filtersvc, role ContractSvc) {
    choice at AuthSvc {
```

```
} }
```



#### Related work

#### Dynamic participation in sessions/conversations

[ESOP09] Conversation types. Caires and Vieira. [POPL11] Dynamic multirole session types. Deniélou and Yoshida. [CONCUR12] Nested protocols in session types. Demangeon and Honda.

#### > Dynamic message sequence charts and communication automata

 [FSTTCS02] Dynamic message sequence charts. Leucker, Madhusudan and Mukhopadhyay.
 [LATA13] Dynamic communication automata and branching high-level MSCs. Bollig, Cyriac, Hélou et, Jara and Schwentick.

#### CFSM-based well-formedness of choreographies and MPST

[POPL08] Deciding choreography realizability. Basu and Bultan.
 [ICALP13] Multiparty Compatibility in Communicating Automata. Deniélou and Yoshida.
 [CONCUR15] Meeting Deadlines Together. Bocchi, Lange and Yoshida.
 [PLACES16] Multiparty compatibility for concurrent objects. Perera, Lange and Gay.

#### Implementations of session types

Java ([ECOOP08,SCP13,PPDP16,FASE16]), Scala ([ECOOP16,ECOOP17]), Haskell ([PADL04,HASKELL08,PLACES10,POPL16,HASKELL16]), OCaml ([JFP17,ESOP17,COORDINATION17]), SILL ([ESOP13,FoSSaCS15]), Links ([ESOP15]), Python ([RV13]), Rust ([WGP15]), C ([TOOLS12]),

. . .

# Conclusions and future work

- (We can finally do Travel Agency in MPST!)
- Practically-motivated extension for explicit connection actions in MPST
  - Scribble toolchain for MPST validation and Endpoint API generation
    - Integrating MPST with existing model checking techniques and tools [TACAS16] Characteristic Formulae for Session Types. Lange and Yoshida.
    - (The session type system interplay with delegation)
  - Other kinds of communication actions? e.g., SSL/TLS connection wrapping (HTTPS, SMTP, FTPS, ...)
  - Integration of further extensions from MPST theory
    - e.g., time, asynchronous interrupts, nested subsessions, message value assertions, role parameterisations, event handling, . . .
- Thanks!
  - https://github.com/scribble/scribble-java
  - https://www.doc.ic.ac.uk/~rhu/scribble/explicit.html