

Efficient SimRank Computation on Large Graphs



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SimRank Overview

• SimRank

- An appealing link-based similarity measure (KDD '02)
- Basic philosophy

Two vertices are similar if they are referenced by similar vertices.

• Two Forms

- Original form (KDD '02)

$$s(a, a) = 1$$

$$s(a, b) = \frac{C}{|\mathcal{I}(a)| |\mathcal{I}(b)|} \sum_{j \in \mathcal{I}(b)} \sum_{i \in \mathcal{I}(a)} s(i, j)$$

damping factor

similarity btw. nodes a and b

- Matrix form (EDBT '10)

$$\mathbf{S} = C \cdot (\mathbf{Q} \cdot \mathbf{S} \cdot \mathbf{Q}^T) + (1 - C) \cdot \mathbf{I}_n$$

in-neighbor set of node b

Motivation

• Prior Work (VLDB J. '10)

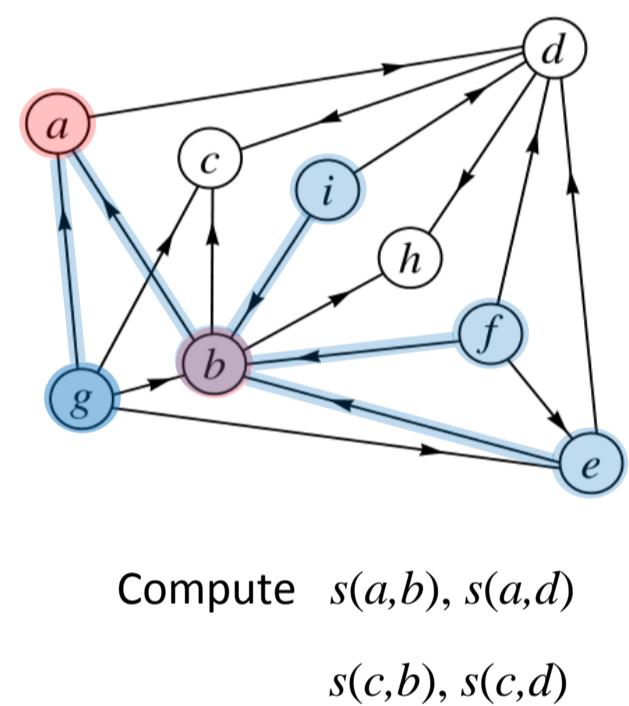
- High time complexity: $O(Kdn^2)$
 - Duplicate computation among partial sums memoization
- Slow (geometric) convergence rate
 - Require $K = \lceil \log_C \epsilon \rceil$ iterations to guarantee accuracy ϵ

• Our Contributions

- Propose an adaptive clustering strategy to reduce the time from $O(Kdn^2)$ to $O(Kd'n^2)$, where $d' \leq d$.
- Introduce a new notion of SimRank to accelerate convergence from geometric to exponential rate.

Duplicate Computations among Partial Sums

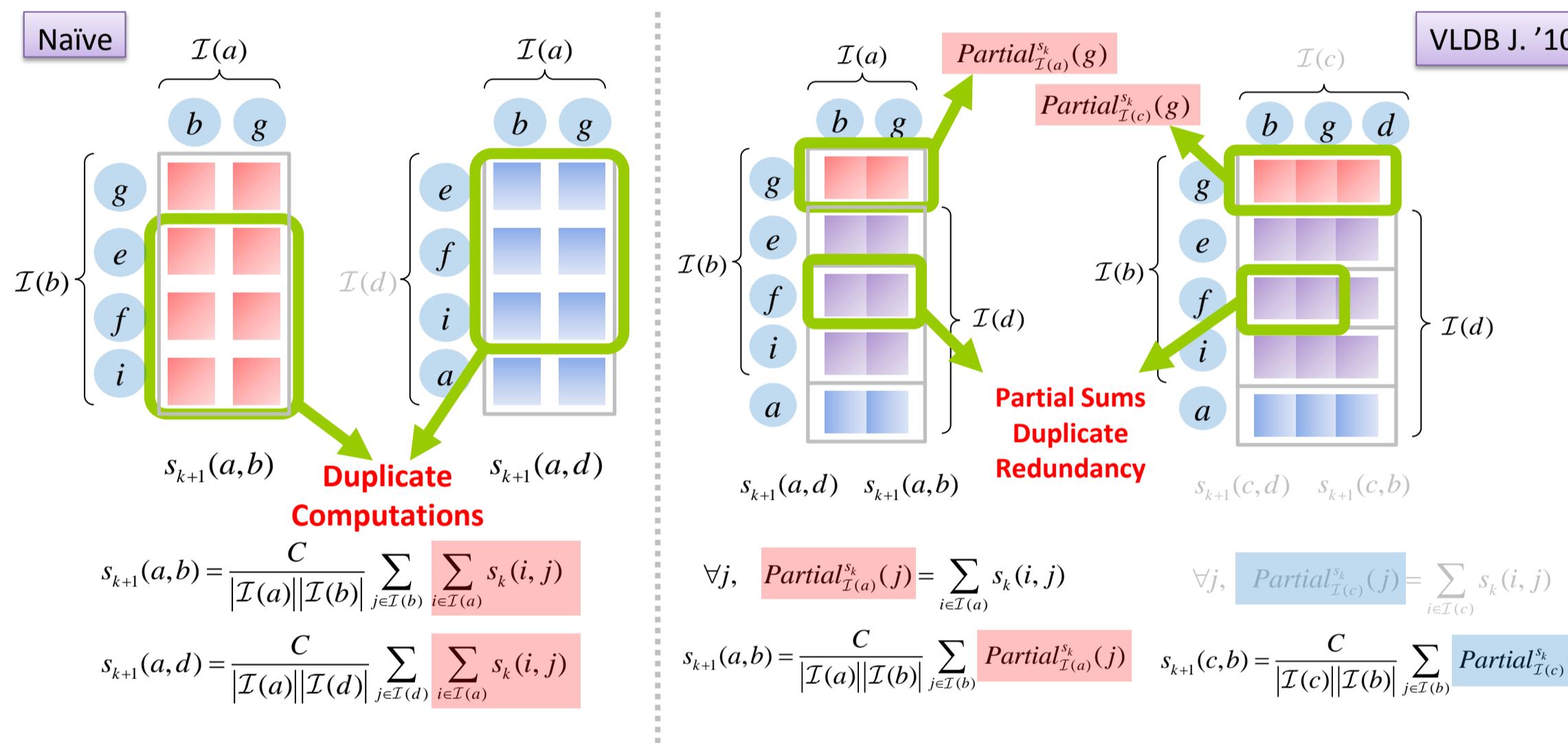
• Example:



Compute $s(a, b), s(a, d)$
 $s(c, b), s(c, d)$

$$s_{k+1}(a, b) = \frac{C}{|\mathcal{I}(a)| |\mathcal{I}(b)|} \sum_{j \in \mathcal{I}(b)} \sum_{i \in \mathcal{I}(a)} s_k(i, j)$$

$$s_{k+1}(a, d) = \frac{C}{|\mathcal{I}(a)| |\mathcal{I}(d)|} \sum_{i \in \mathcal{I}(d)} \sum_{j \in \mathcal{I}(a)} s_k(i, j)$$



Slow Convergence

• Existing Approach (VLDB J. '10)

$$\|\mathbf{S}_k - \mathbf{S}\|_{\max} \leq C^{k+1}$$

Geometric Rate

For $C = 0.8$, $\epsilon = 0.0001$,
 $K = \lceil \log_{0.8} 0.0001 \rceil = 41$ iterations.

• Our Approach

$$\|\hat{\mathbf{S}}_k - \hat{\mathbf{S}}\|_{\max} \leq \frac{C^{k+1}}{(k+1)!}$$

Exponential Rate

For $C = 0.8$, $\epsilon = 0.0001$,
there are only 7 iterations.

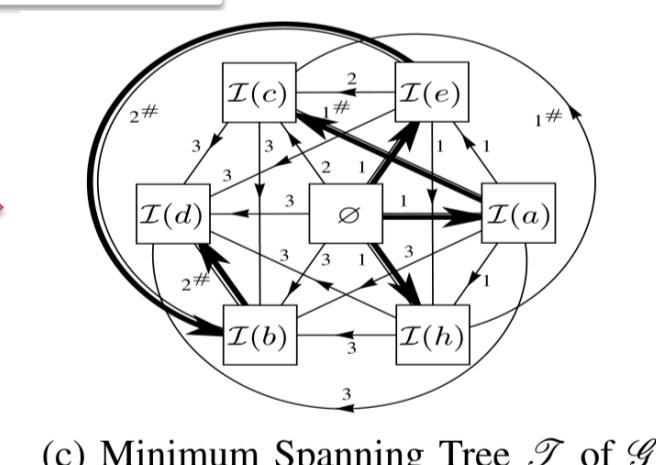
Eliminating Partial Sums Redundancy

vertex	$\mathcal{I}(\star)$
a	{ b, g }
e	{ f, g }
h	{ b, d }
c	{ b, d, g }
b	{ f, g, e, i }
d	{ f, a, e, i }

(a) In-neighbors in \mathcal{G}

	$\mathcal{I}(a)$	$\mathcal{I}(c)$	$\mathcal{I}(h)$	$\mathcal{I}(c)$	$\mathcal{I}(b)$	$\mathcal{I}(d)$
\emptyset	1	1	1	2	3	3
$\mathcal{I}(a)$		1	1	1 [#]	3	3
$\mathcal{I}(e)$			1	2	2 [#]	3
$\mathcal{I}(h)$				1 [#]	3	3
$\mathcal{I}(c)$					3	3
$\mathcal{I}(b)$						2 [#]

(b) Transition Costs (Edge Weights) in \mathcal{G}



(c) Minimum Spanning Tree \mathcal{T} of \mathcal{G}

	$\mathcal{P}(\star)$
$\mathcal{I}(a)$	{ b, g }
$\mathcal{I}(e)$	{ f, g }
$\mathcal{I}(h)$	{ b, d }
$\mathcal{I}(c)$	{ $\mathcal{I}(a), \{d\}$ }
$\mathcal{I}(b)$	{ $\mathcal{I}(e), \{e, i\}$ }
$\mathcal{I}(d)$	{ $\mathcal{I}(b) \setminus \{g\}, \{a\}$ }

• (Inner) partial sums sharing

$$\text{Partial}_{\mathcal{I}(a)}^{s_k}(\star) = s_k(b, \star) + s_k(g, \star)$$

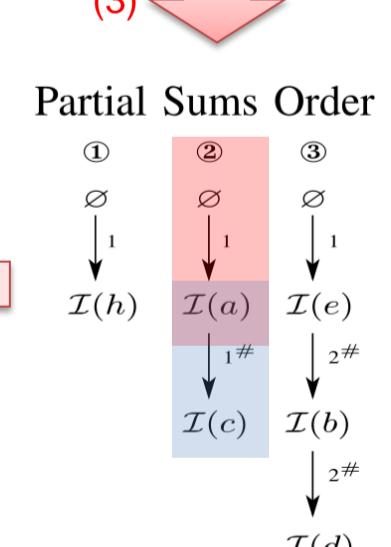
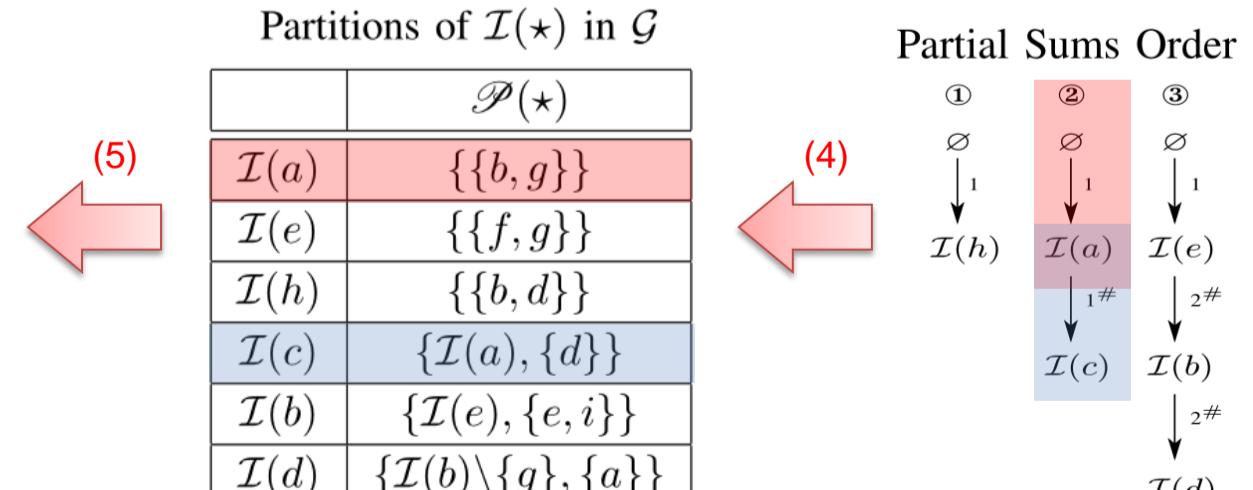
$$\text{Partial}_{\mathcal{I}(a)}^{s_k}(\star) = \text{Partial}_{\mathcal{I}(a)}^{s_k}(\star) + s_k(d, \star)$$

• Outer partial sums sharing

$$\text{OuterPartial}_{\mathcal{I}(a)}^{s_k, s_k} = \sum_{y \in \mathcal{I}(b, g)} \text{Partial}_{\mathcal{I}(a)}^{s_k}(y)$$

$$\text{OuterPartial}_{\mathcal{I}(a)}^{s_k, s_k} = \text{OuterPartial}_{\mathcal{I}(a)}^{s_k} + \text{Partial}_{\mathcal{I}(a)}^{s_k}(d)$$

$$s_{k+1}(a, \star) = \frac{C}{|\mathcal{I}(a)| |\mathcal{I}(\star)|} \text{OuterPartial}_{\mathcal{I}(a)}^{s_k, s_k}$$



From Geometric to Exponential Rate

$$\mathbf{S} = C \cdot (\mathbf{Q} \cdot \mathbf{S} \cdot \mathbf{Q}^T) + (1 - C) \cdot \mathbf{I}_n \quad \Rightarrow \quad \mathbf{S} = (1 - C) \cdot \sum_{i=0}^{\infty} C^i \cdot \mathbf{Q}^i \cdot (\mathbf{Q}^T)^i$$

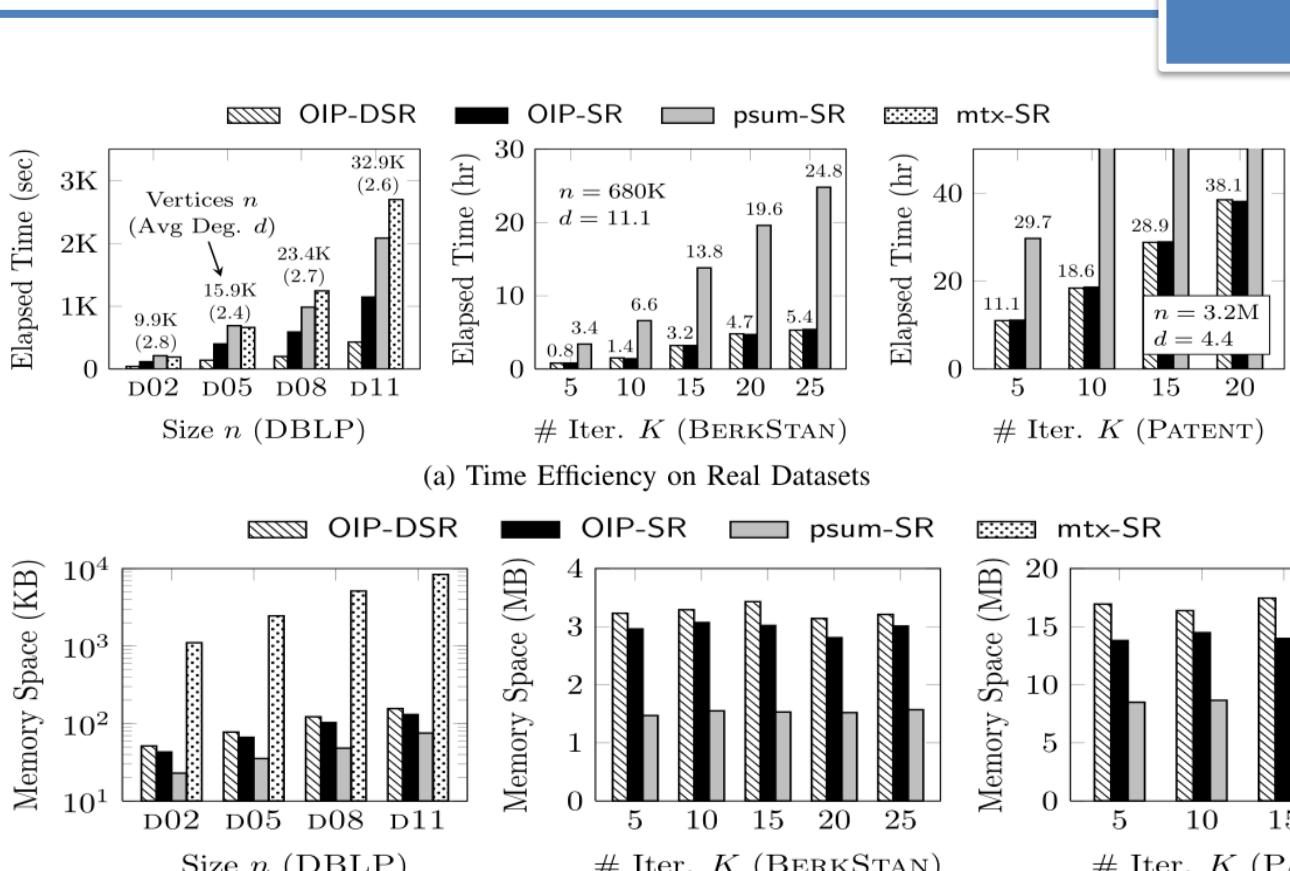
Geometric Sum

$$\frac{d\hat{\mathbf{S}}(t)}{dt} = \mathbf{Q} \cdot \hat{\mathbf{S}}(t) \cdot \mathbf{Q}^T, \quad \hat{\mathbf{S}}(0) = e^{-C} \cdot \mathbf{I}_n. \quad \Rightarrow \quad \hat{\mathbf{S}} = e^{-C} \cdot \sum_{i=0}^{\infty} \frac{C^i}{i!} \cdot \mathbf{Q}^i \cdot (\mathbf{Q}^T)^i$$

Exponential Sum

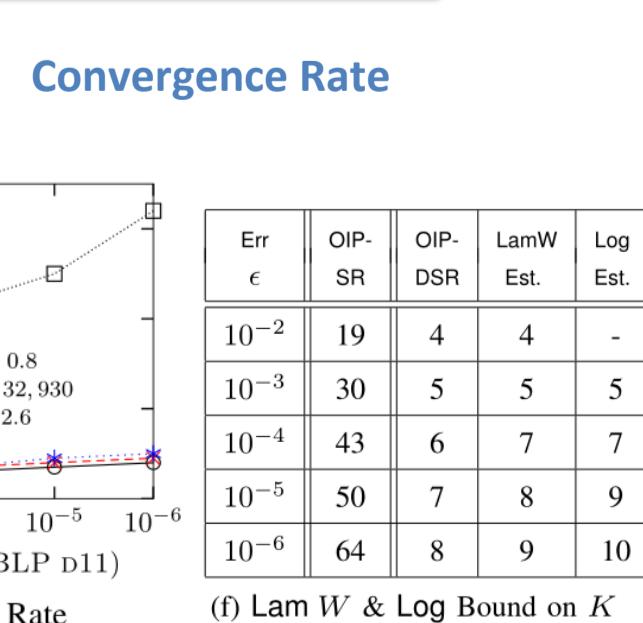
	(Geometric) SimRank	Exponential SimRank
Closed Form	$\mathbf{S} = C \cdot (\mathbf{Q} \cdot \mathbf{S} \cdot \mathbf{Q}^T) + (1 - C) \cdot \mathbf{I}_n$	$\frac{d\hat{\mathbf{S}}(t)}{dt} = \mathbf{Q} \cdot \hat{\mathbf{S}}(t) \cdot \mathbf{Q}^T, \quad \hat{\mathbf{S}}(0) = e^{-C} \cdot \mathbf{I}_n.$
Series Form	$\mathbf{S} = (1 - C) \cdot \sum_{i=0}^{\infty} C^i \cdot \mathbf{Q}^i \cdot (\mathbf{Q}^T)^i$	$\hat{\mathbf{S}} = e^{-C} \cdot \sum_{i=0}^{\infty} \frac{C^i}{i!} \cdot \mathbf{Q}^i \cdot (\mathbf{Q}^T)^i$
Iterative Form	$\mathbf{S}_0 = \mathbf{I}_n$ $\mathbf{S}_{k+1} = C \cdot (\mathbf{Q} \cdot \mathbf{S}_k \cdot \mathbf{Q}^T) + (1 - C) \cdot \mathbf{I}_n$	$\left\{ \begin{array}{l} \mathbf{T}_{k+1} = \mathbf{Q} \cdot \mathbf{T}_k \cdot \mathbf{Q}^T \\ \hat{\mathbf{S}}_{k+1} = \hat{\mathbf{S}}_k + e^{-C} \cdot \frac{C^{k+1}}{(k+1)!} \cdot \mathbf{T}_{k+1} \\ \hat{\mathbf{S}}_0 = e^{-C} \cdot \mathbf{I}_n \end{array} \right.$
Error	$\ \mathbf{S}_k - \mathbf{S}\ _{\max} \leq C^{k+1}$	$\ \hat{\mathbf{S}}_k - \hat{\mathbf{S}}\ _{\max} \leq \frac{C^{k+1}}{(k+1)!}$

Time & Space Efficiency



Experimental Evaluations

Convergence Rate



Relative Order Preservation

