

RWset: Attacking Path Explosion in Constraint-Based Test Generation

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Constraint-Based Test Generation

- Goal: generate inputs that explore (ideally) all paths of a program
- Run program on **symbolic** input, whose initial value is *anything*
- At conditionals that use symbolic inputs, fork execution and follow both paths:
 - On true branch, add constraint that condition is true
 - On false, that it is not
- When a path terminates, generate a test case by solving the constraints on that path



Constraint-Based Test Generation

- EGT / EXE / KLEE
- DART [Godefroid/Klarlund/Sen]
- CUTE [Sen et al.]
- SAGE, Pex [Godefroid et al.]
- Vigilante [Castro et al]
- BitScope [Song et al.]

- RWset applicable to any of these



EXE Results

- ❑ Effective bug-finding tool
 - File system code
 - ext2, ext3, JFS
 - Networking applications
 - bpf, udhcpd
 - Library code
 - PCRE, Pintos
 - Device drivers
 - Minix



Scalability challenge

- Exponential space!
 - Relatively small number of interesting paths: e.g., those that achieve maximum branch coverage
- Mixed symbolic/concrete execution (EXE/DART)
- Search heuristics
 - Best First Search (EXE)
 - Generational Search (SAGE)
- Symbolic execution + random testing (Hybrid CUTE)
- Caching function summaries (SMART)
- Demand-Driven Compositional Symbolic Execution (Pex)



RWset (read-write set) analysis

- Determine whether continuing to execute the current program path will explore new states
- Only a value observed by the program can determine the execution of new program states

```
{data, arg1, arg2} = unconstrained
```

```
flag = 0;
```

```
if (arg1 > 100)
```

```
    flag = 1;
```

```
if (arg2 > 100)
```

```
    flag = 1;
```

```
process(data, flag);
```

flag = 0

{data, arg1, arg2} = unconstrained

flag = 0;

if (arg1 > 100)

 flag = 1;

if (arg2 > 100)

 flag = 1;

process(data, flag);

{data, arg1, arg2} = unconstrained

flag = 0;

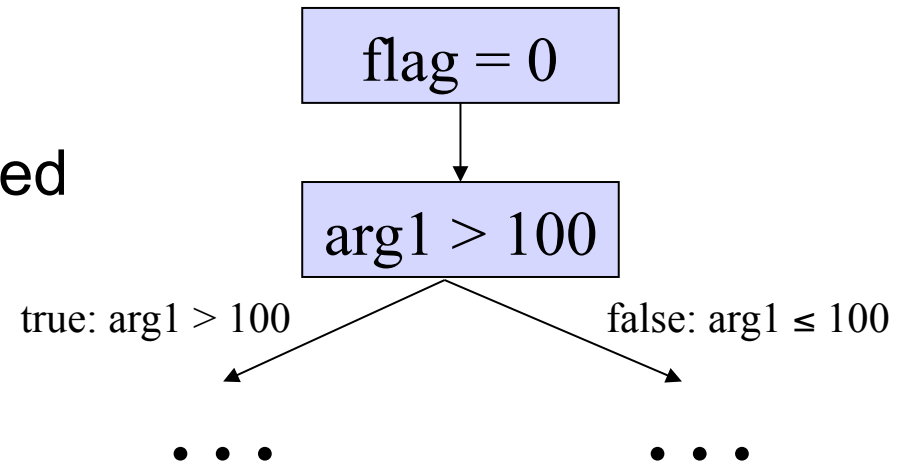
if (arg1 > 100)

flag = 1;

if (arg2 > 100)

flag = 1;

process(data, flag);



{data, arg1, arg2} = unconstrained

flag = 0;

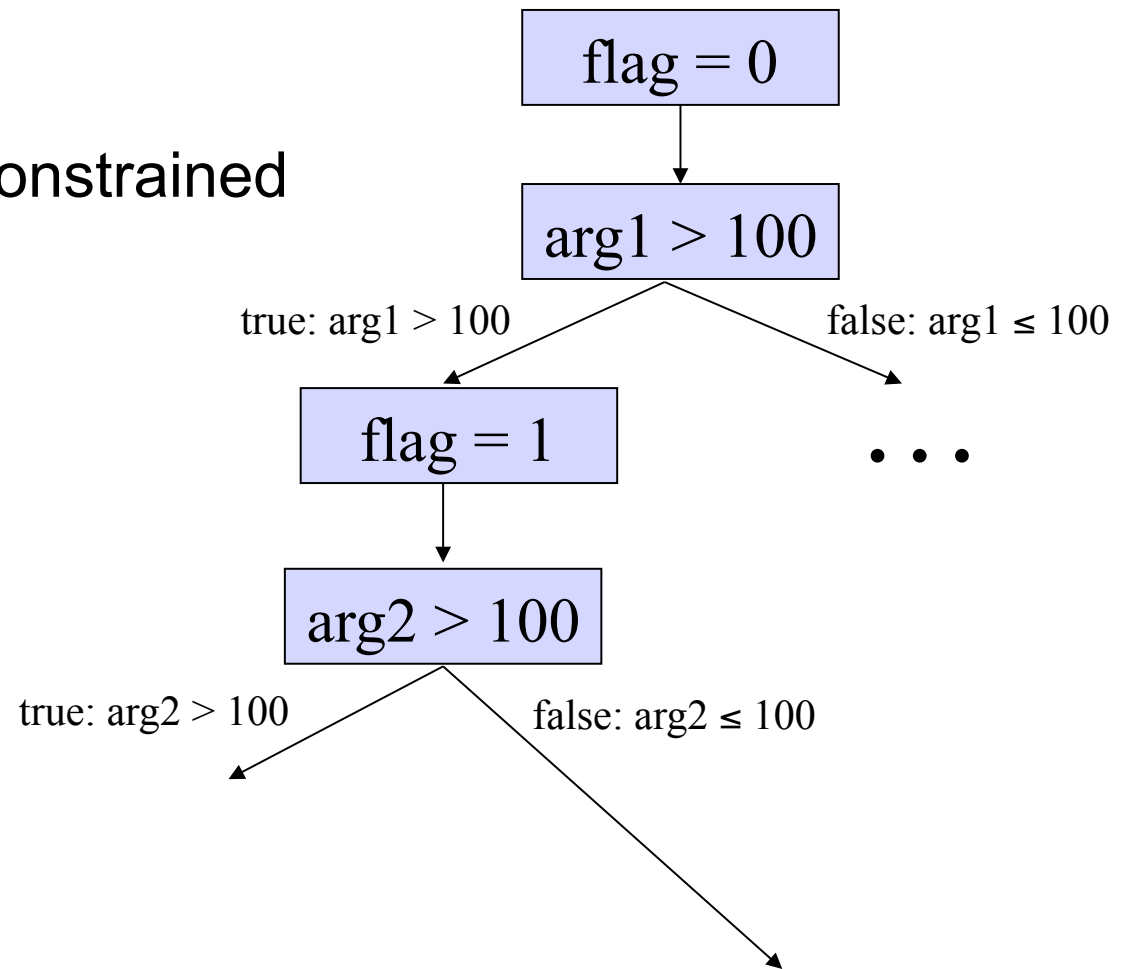
if (arg1 > 100)

flag = 1;

if (arg2 > 100)

flag = 1;

process(data, flag);



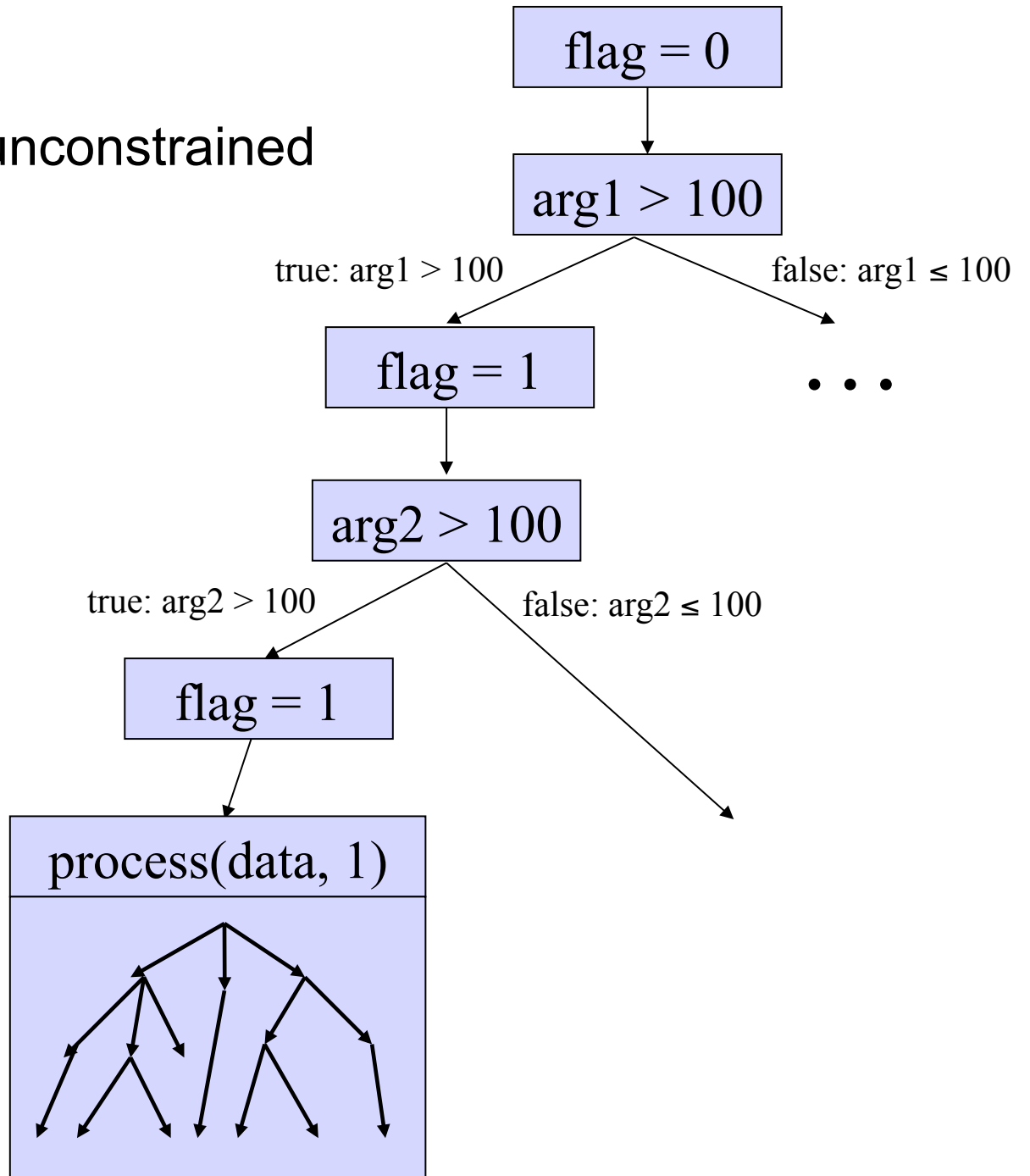
{data, arg1, arg2} = unconstrained

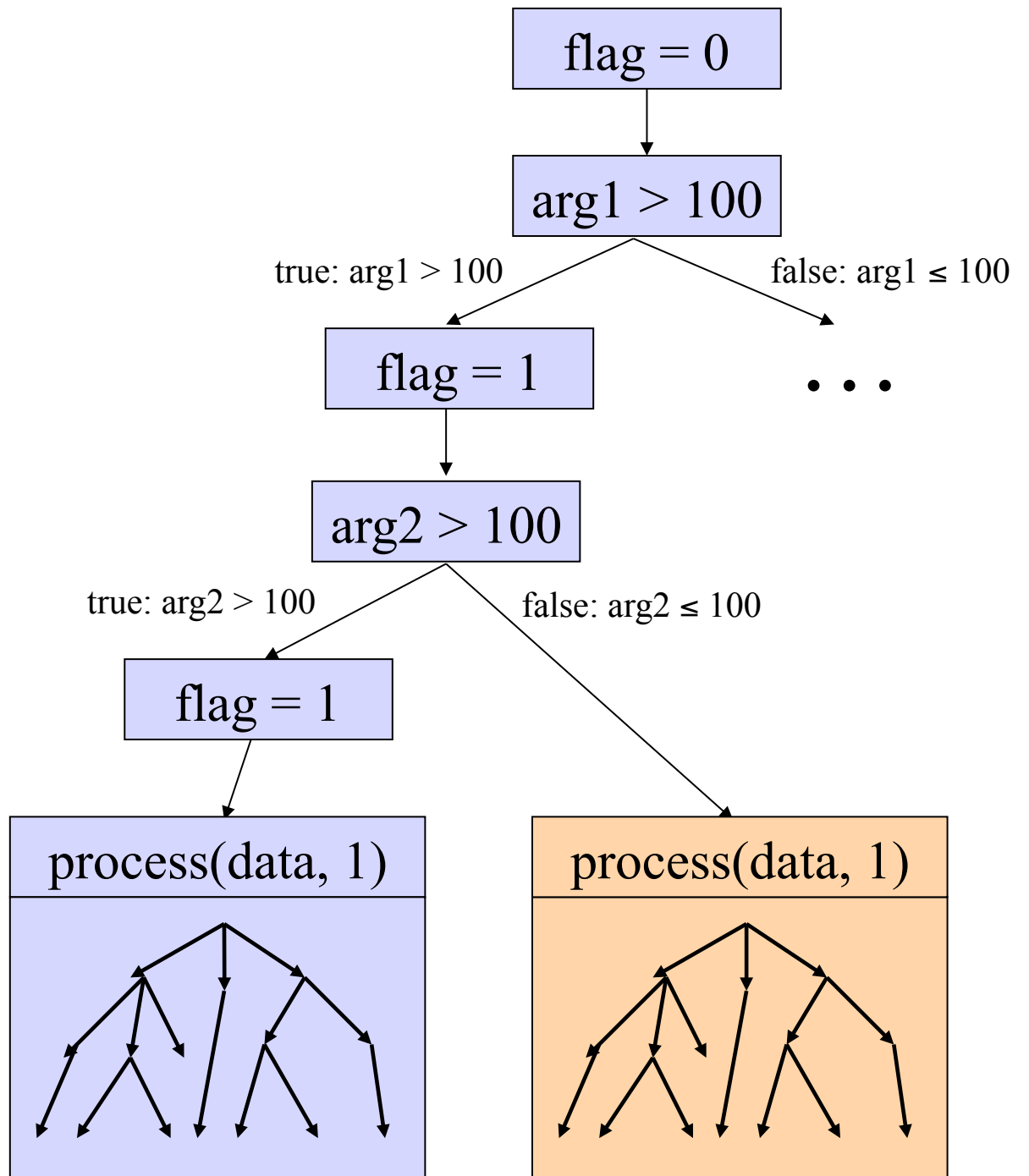
flag = 0;

if (arg1 > 100)
 flag = 1;

if (arg2 > 100)
 flag = 1;

process(data, flag);





If arg1, arg2 not read



Write-set analysis

- Look at values written in the past
- State abstraction in model checking
- Memory state = *write set* up to current progr. pt.
 - Concrete writes: *concr loc = concr val*
 - Symbolic writes: *constraint(sym loc)*
- Program point P, two paths w/ same write-set
 - prune the second one



Write-sets

- Precision: reason at the byte-level
- Minimize write-set size
 1. Overwrites
 2. Dead locations
 3. Alpha renaming
- Complicated in the symbolic domain

```
a[i] = 17;  
a[j] = 20;
```

- Cannot discard all constraints on dead locations



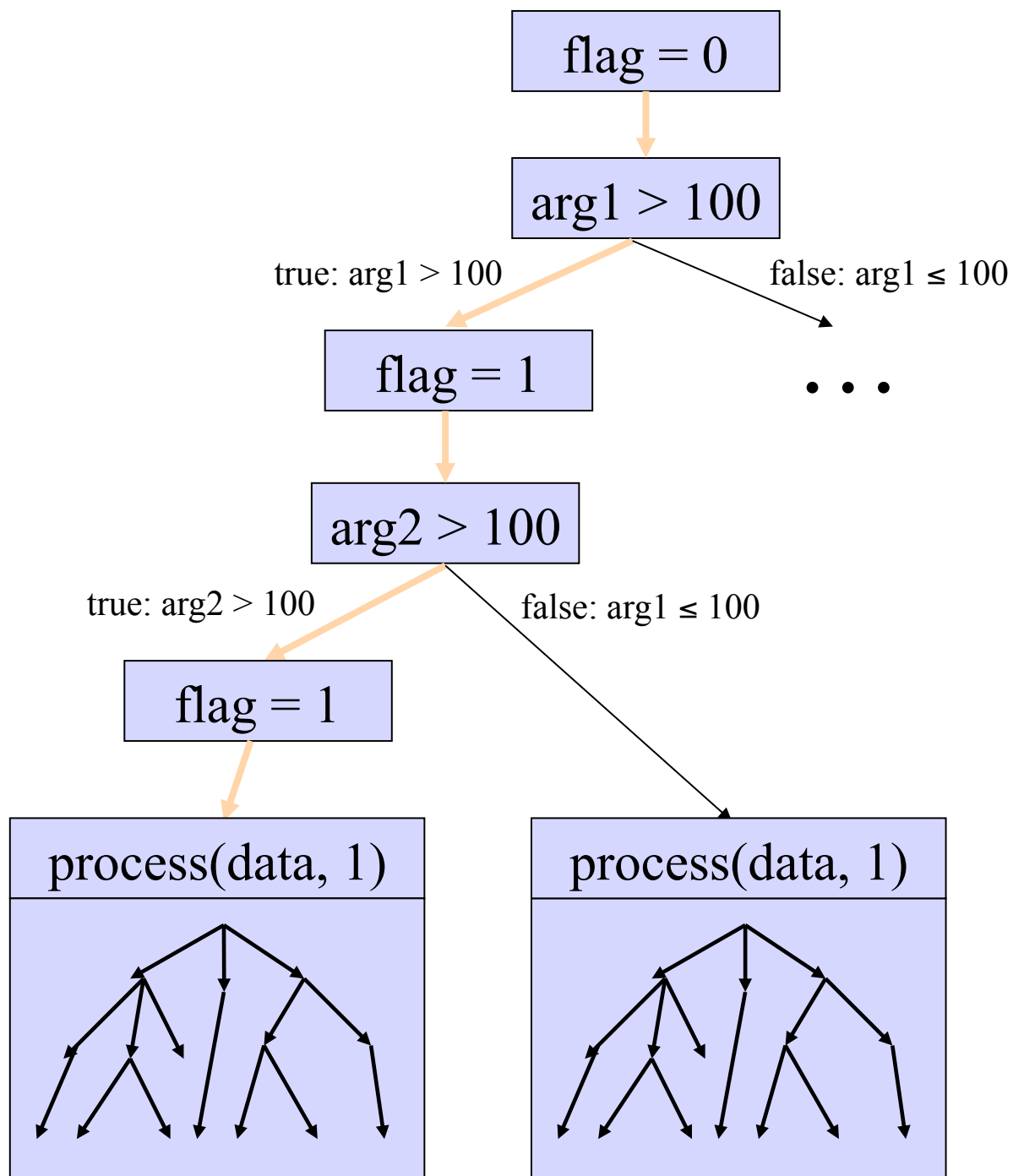
Read-set analysis

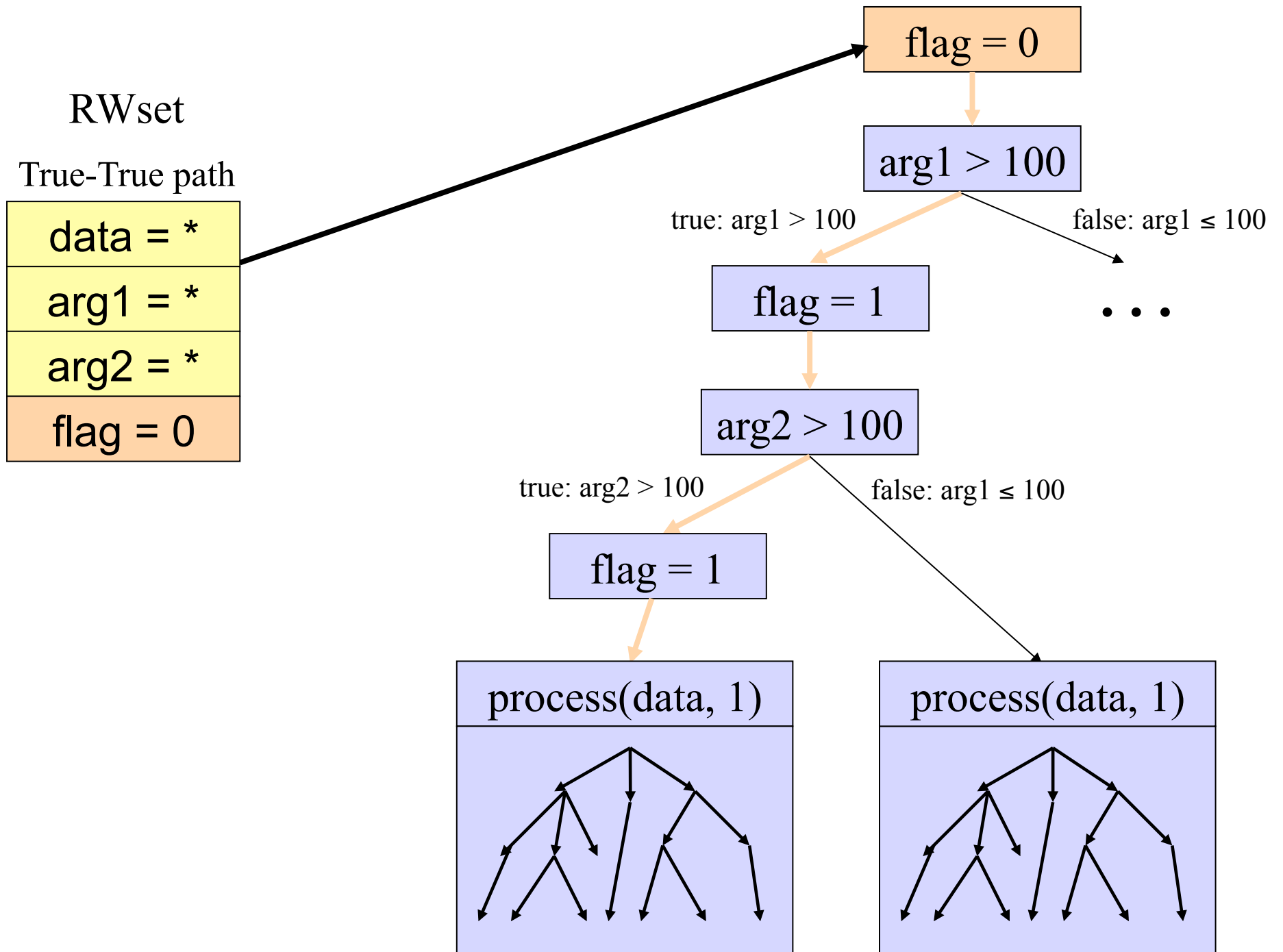
- *Key idea*: can ignore from write-set writes to locations that are never *read* again
 - Definition of *read* driven by goal to achieve high branch coverage
 - Location read if can hit a new branch by changing its value

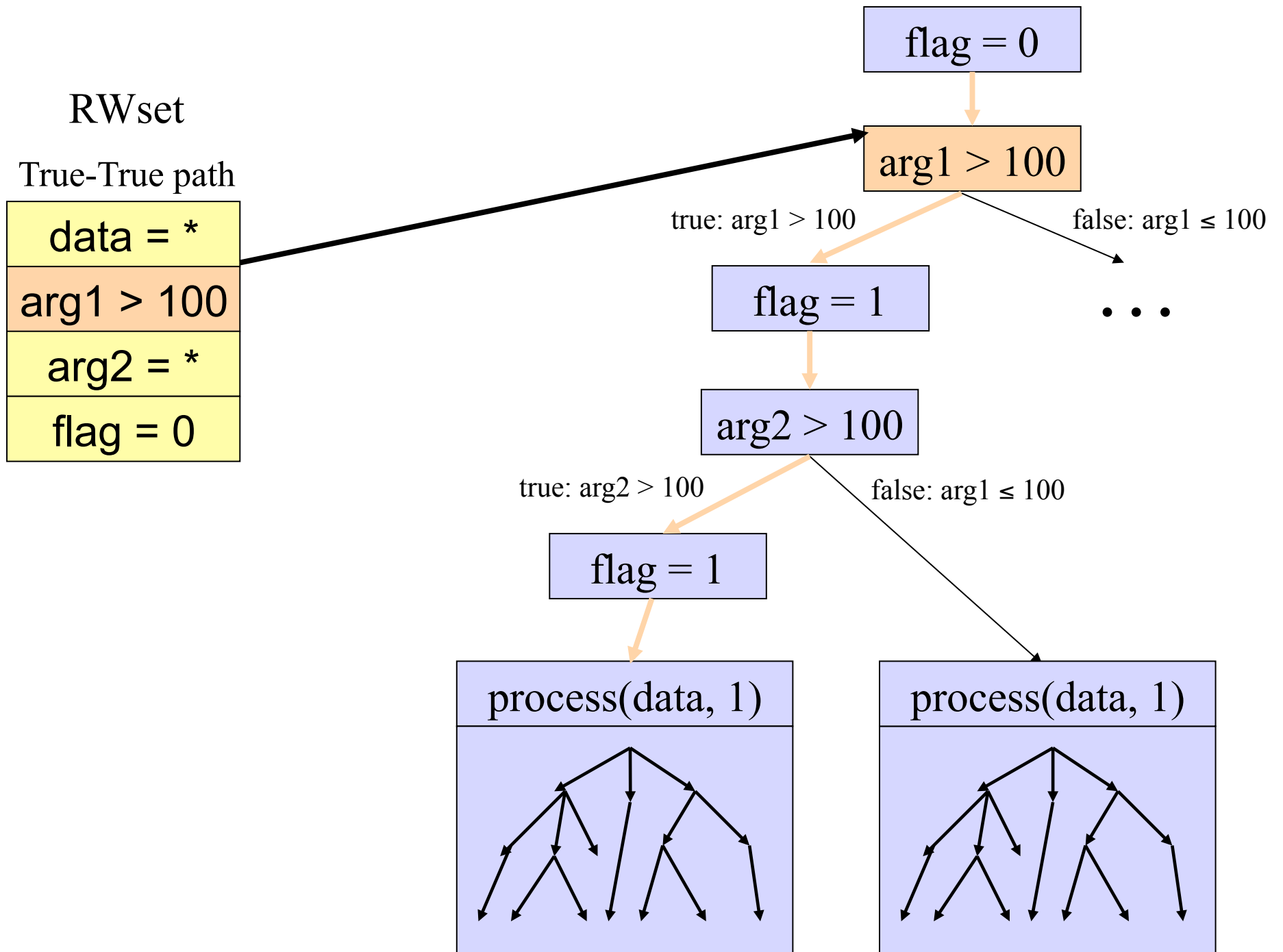
RWset

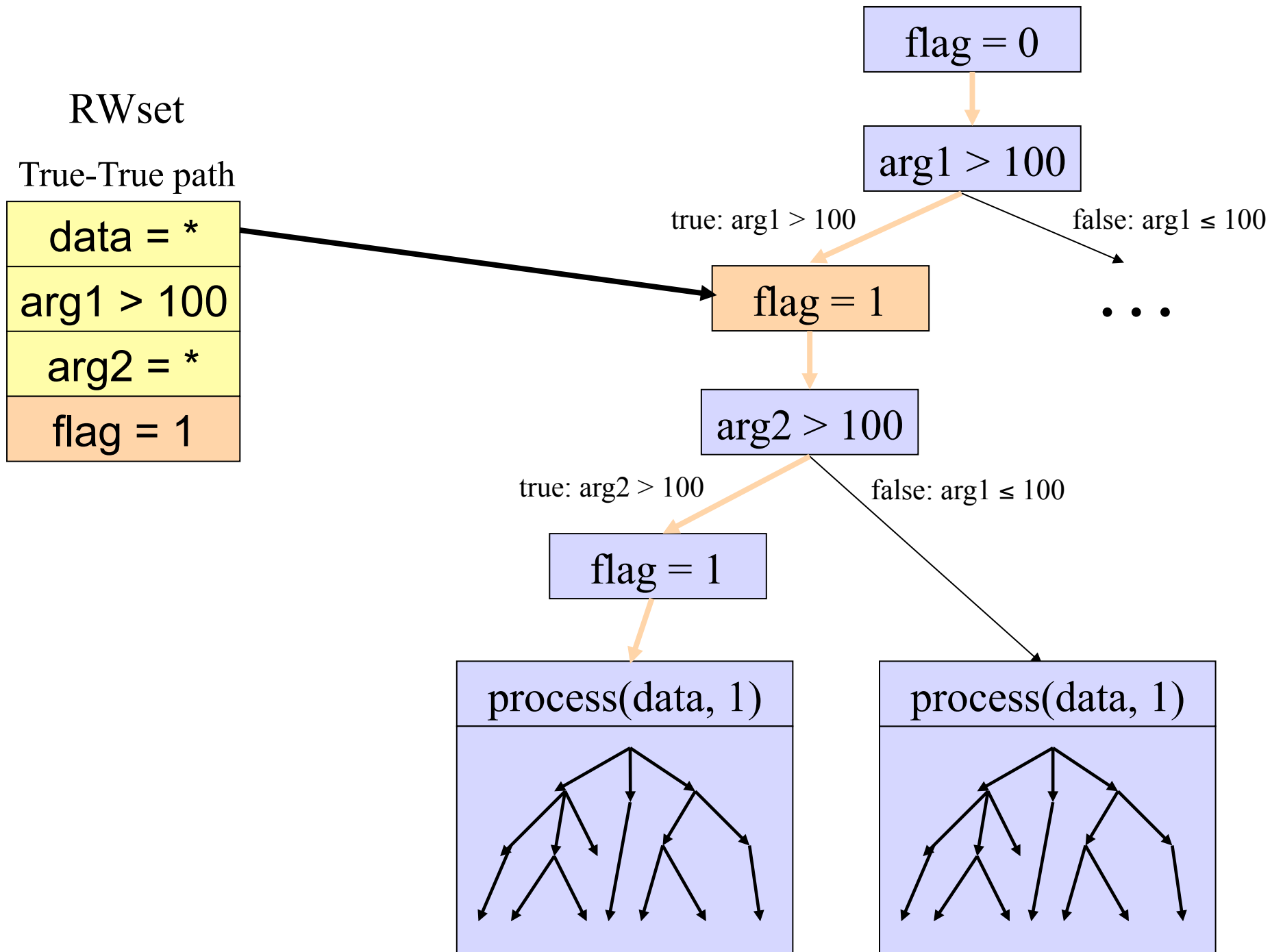
True-True path

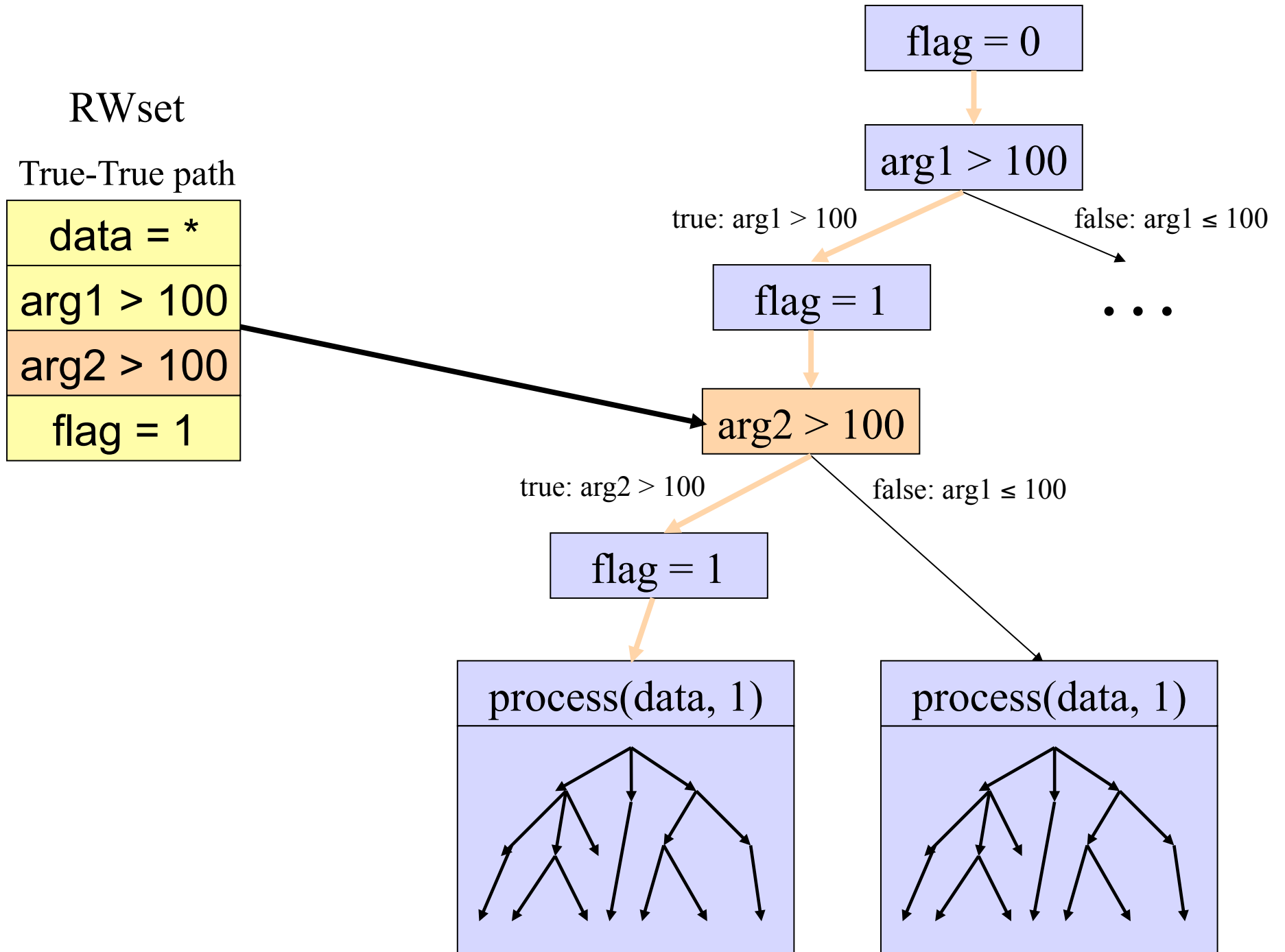
data = *
arg1 = *
arg2 = *

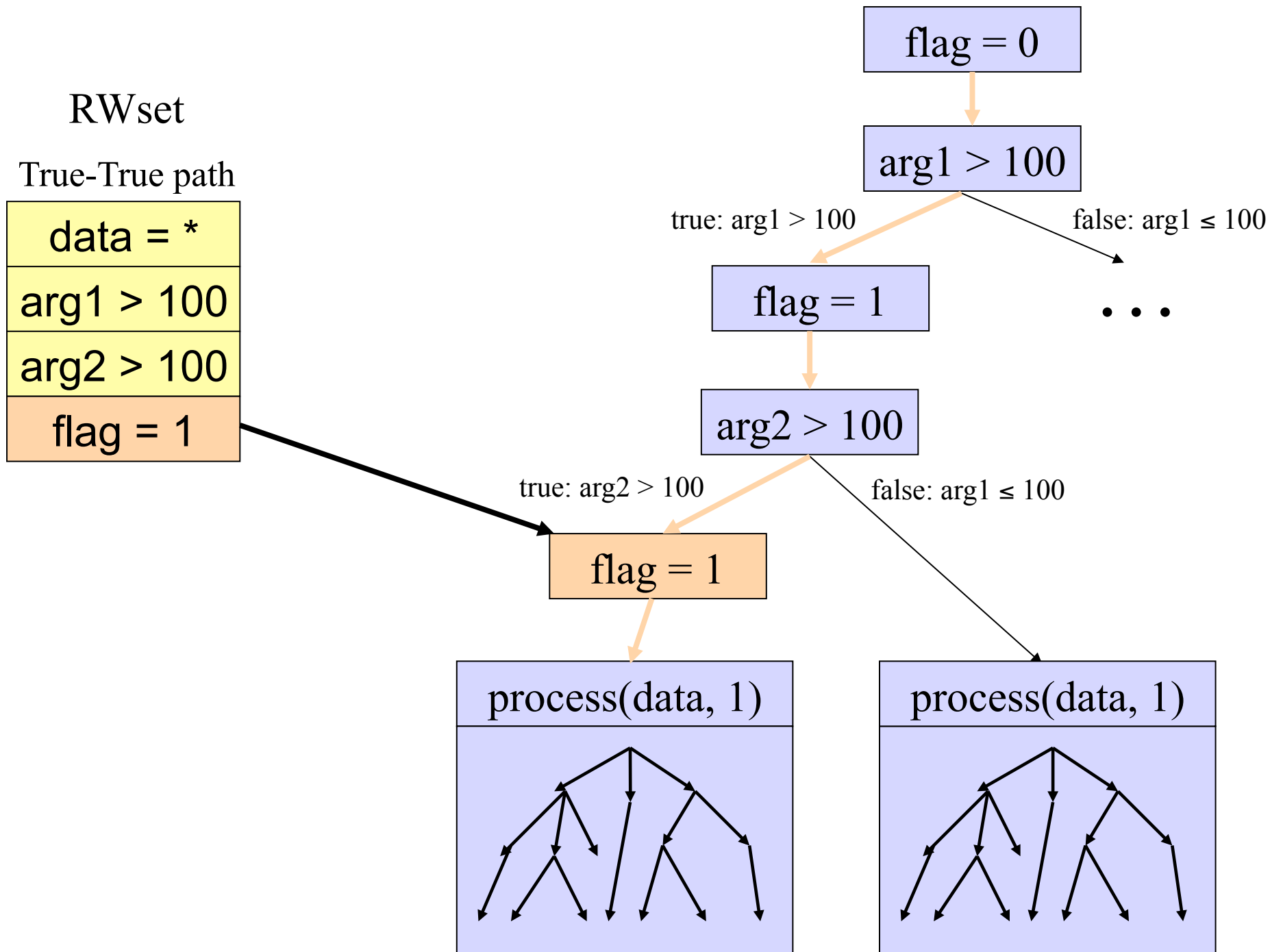












RWset

True-True path

data = *
arg1 > 100
arg2 > 100
flag = 1

flag = 0

arg1 > 100

true: arg1 > 100

false: arg1 ≤ 100

flag = 1

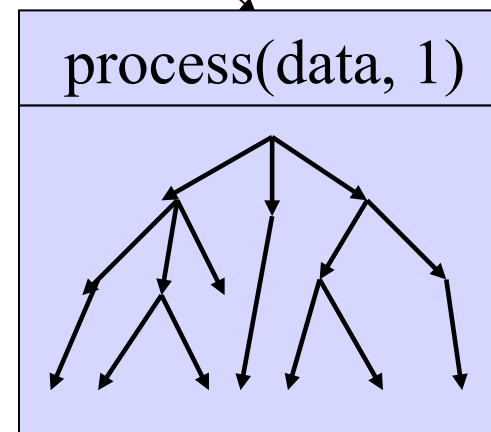
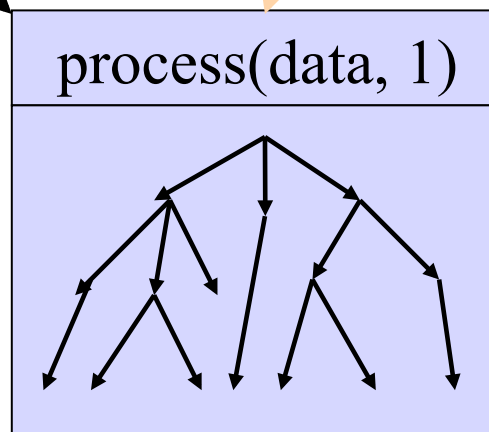
...

arg2 > 100

true: arg2 > 100

false: arg1 ≤ 100

flag = 1



RWset

True-True path

data = *
arg1 > 100
arg2 > 100
flag = 1

flag = 0

arg1 > 100

true: arg1 > 100

false: arg1 ≤ 100

flag = 1

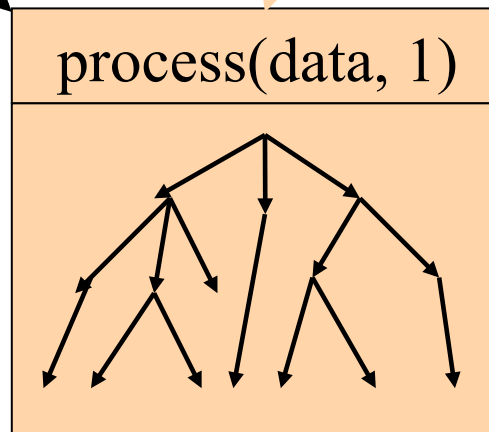
...

arg2 > 100

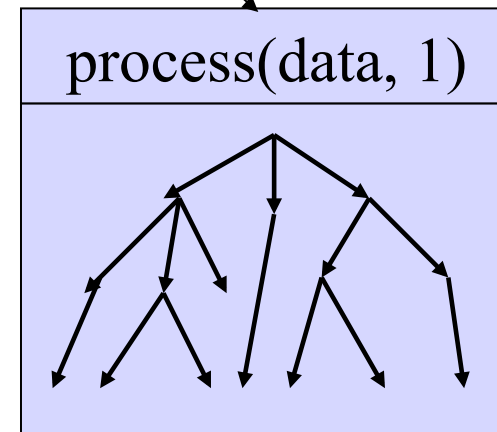
true: arg2 > 100

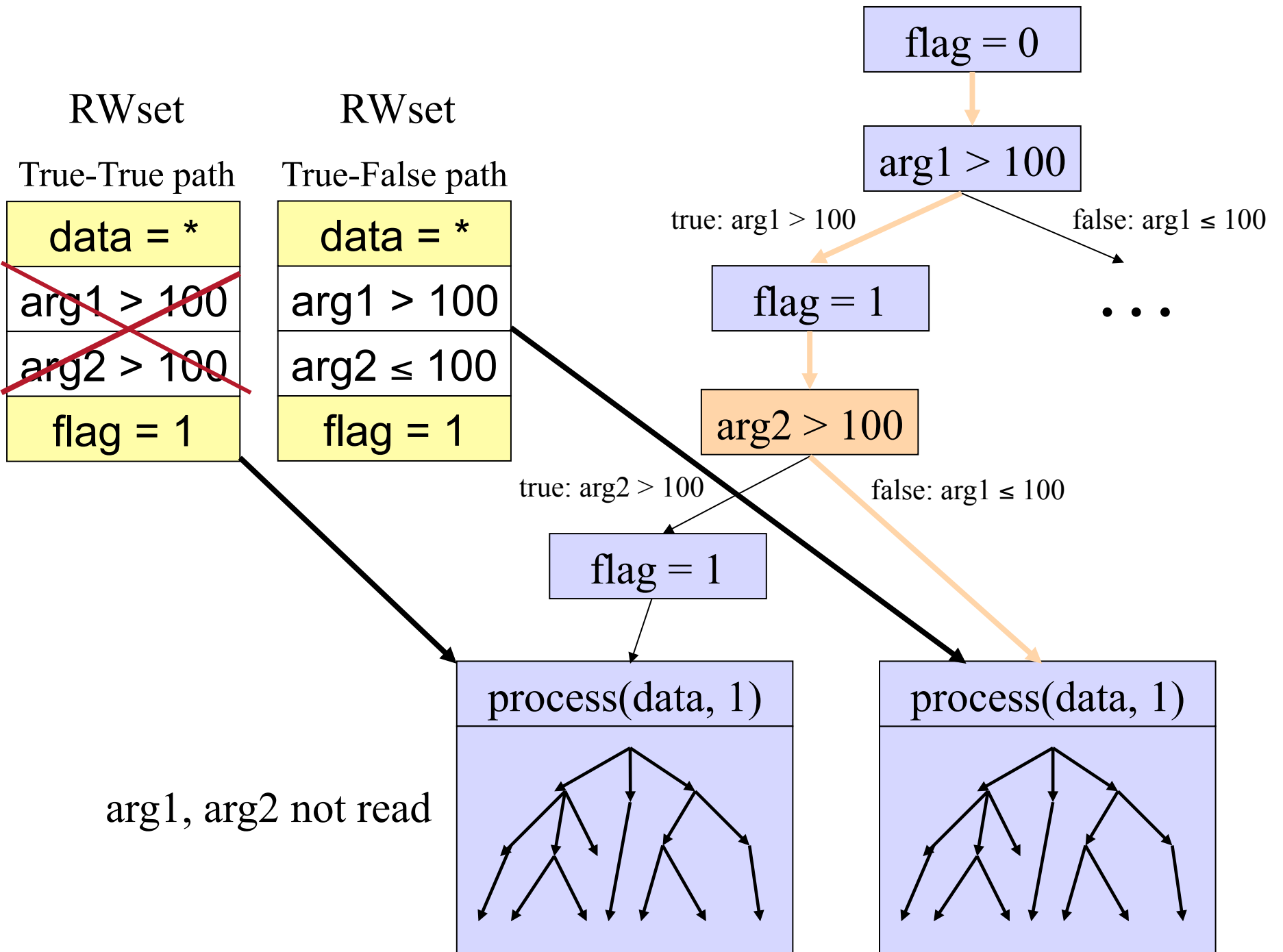
false: arg1 ≤ 100

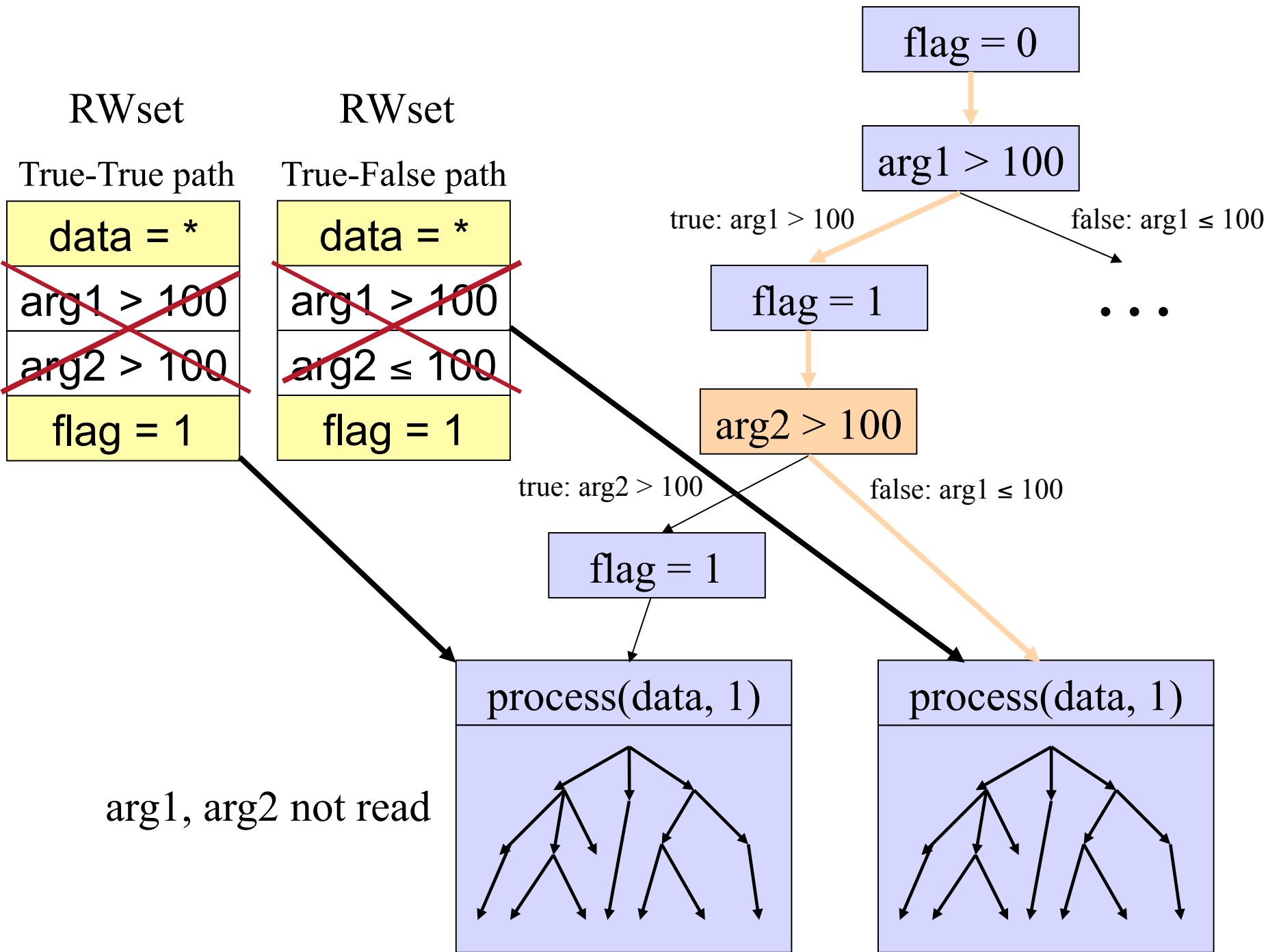
flag = 1

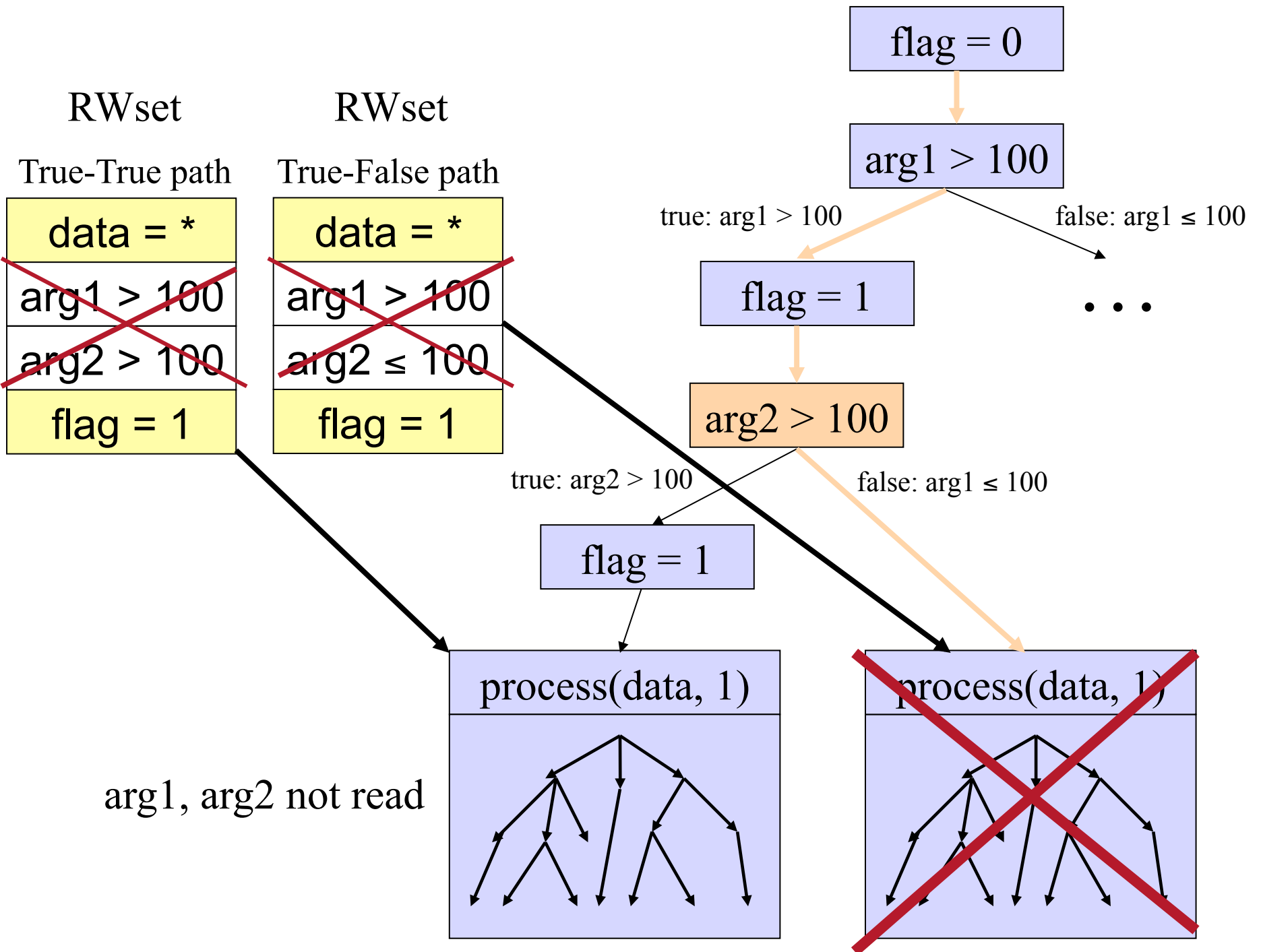


arg1, arg2 not read





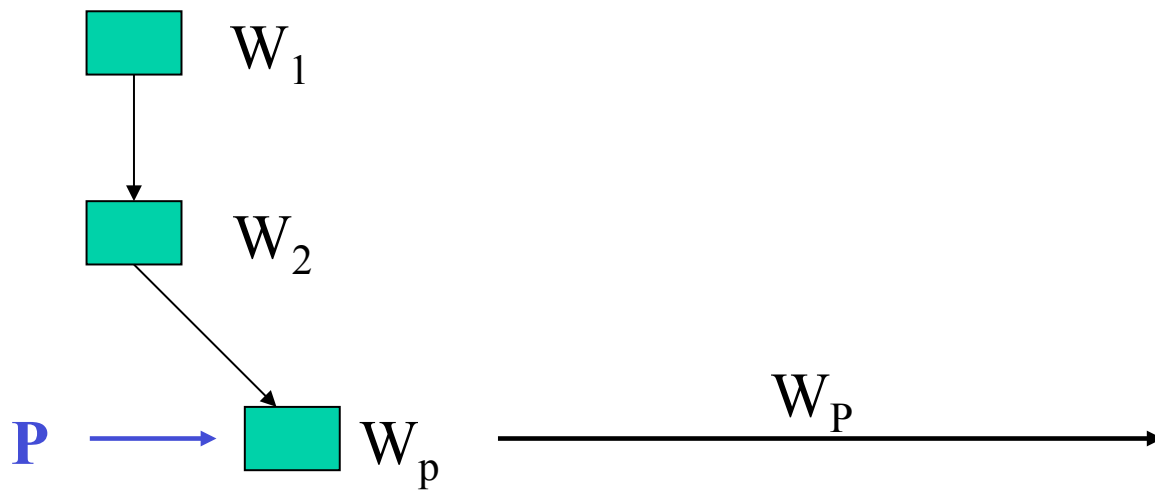




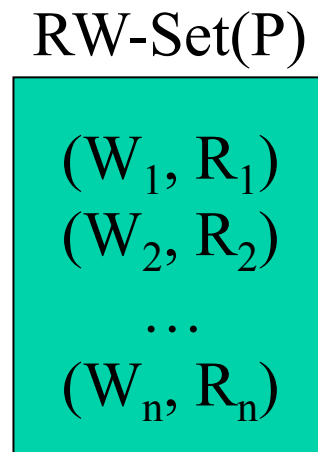


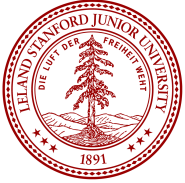
Implementation

Execution path



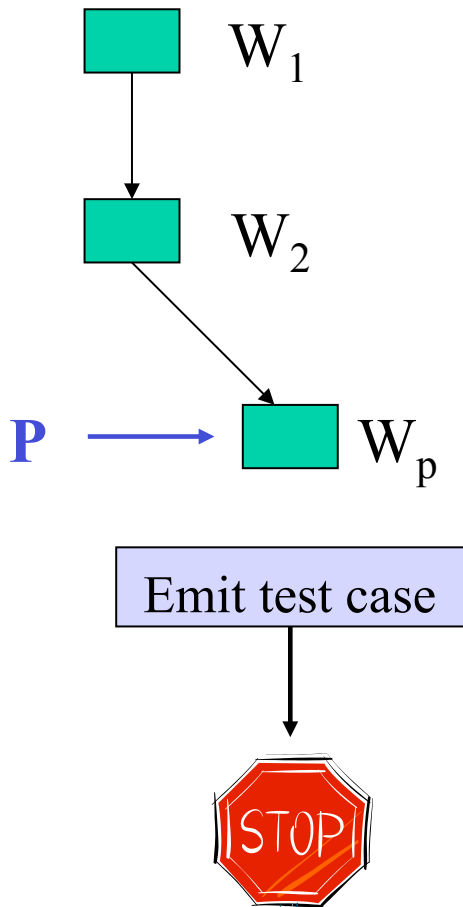
Global cache





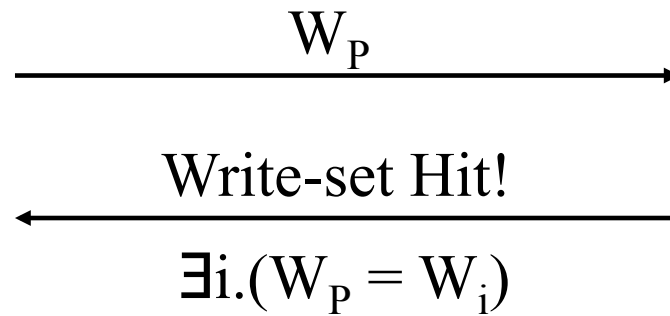
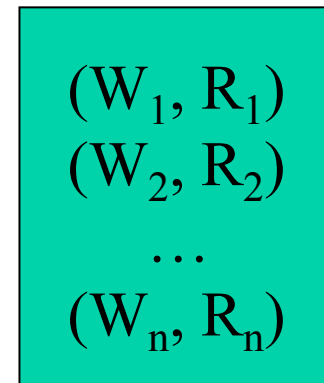
Implementation

Execution path



Global cache

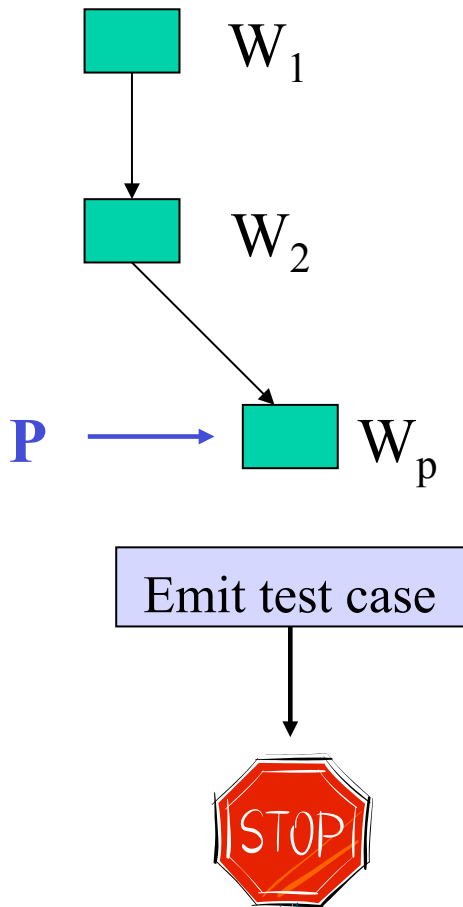
RW-Set(P)





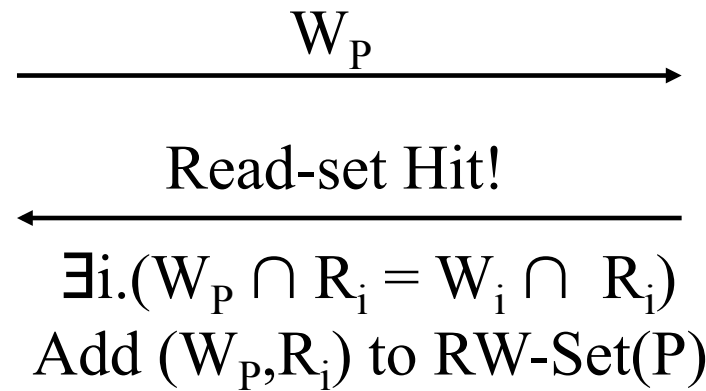
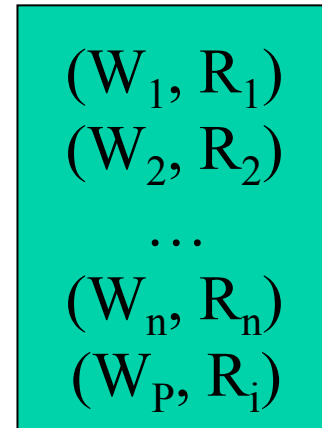
Implementation

Execution path



Global cache

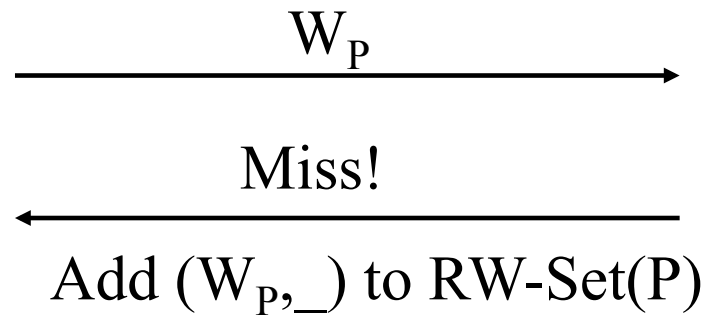
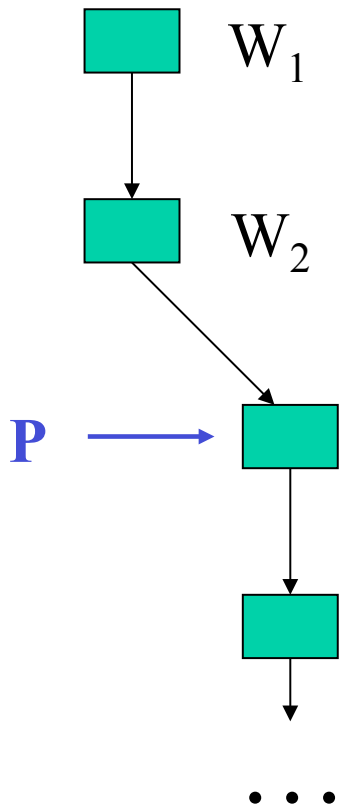
RW-Set(P)





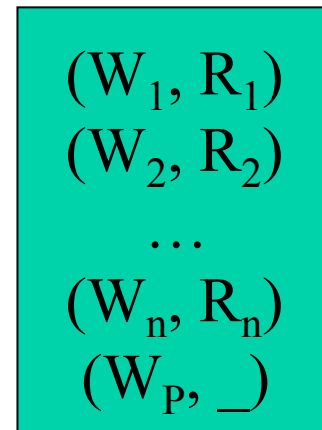
Implementation

Execution path



Global cache

RW-Set(P)





Evaluation

- Medium-sized open source benchmarks
 - bpf, udhcpd, expat, tcpdump, pcre
- Minix 3 device drivers
 - lance, pci, sb16



Medium-sized apps

- *bpf*: Berkeley Packet Filter
- *expat*: XML parsing library
- *pcre*: Perl compatible reg exp library
- *tcpdump*: tool for printing packet headers
- *udhcpd*: a DHCPD server

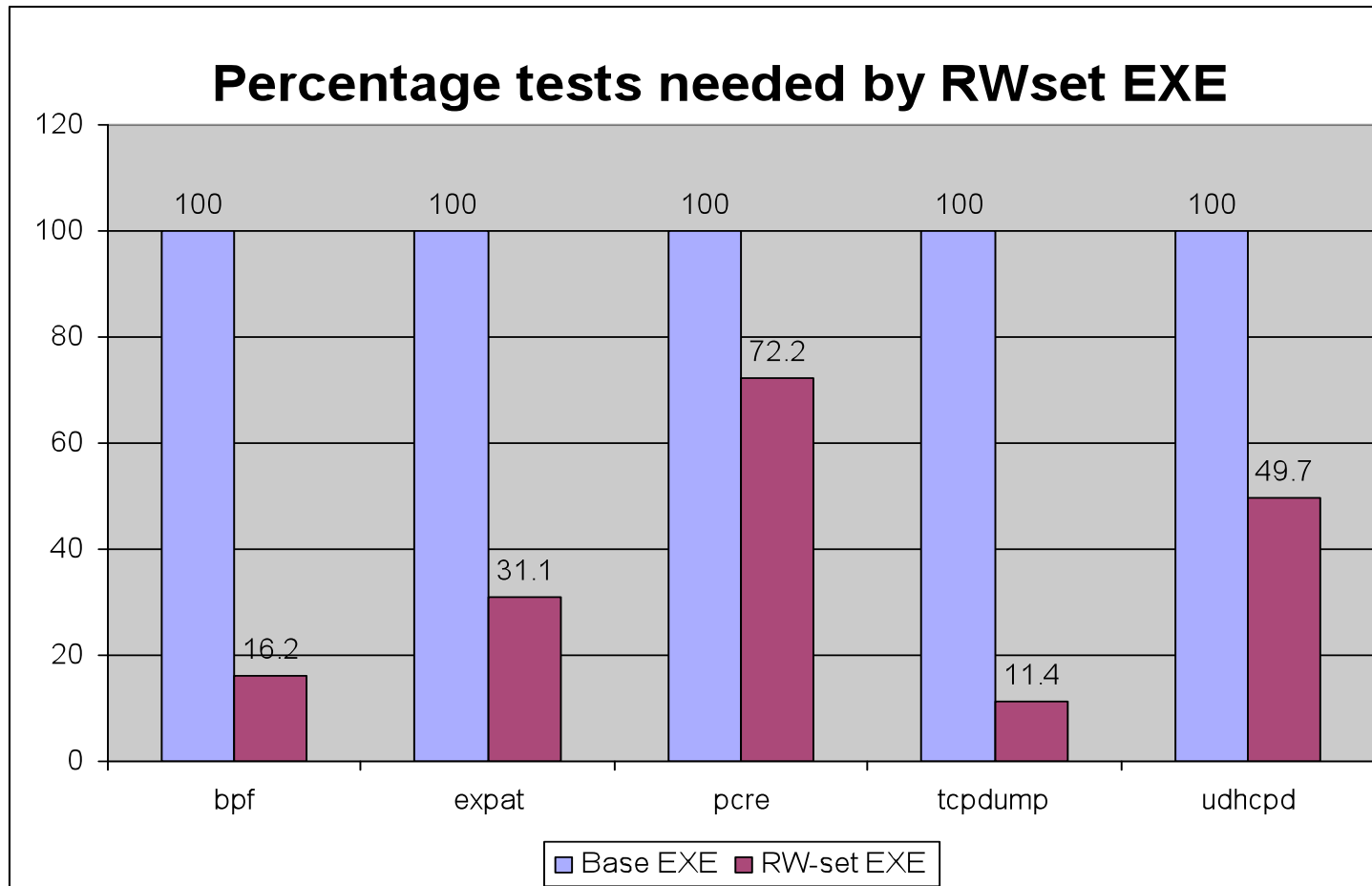


Medium-sized apps

- Ran base EXE on each app for 30 minutes
 - Except 30,000 test cases for PCRE
- Recorded number of branches hit
- Reran in RWset mode until we reached the same branch coverage

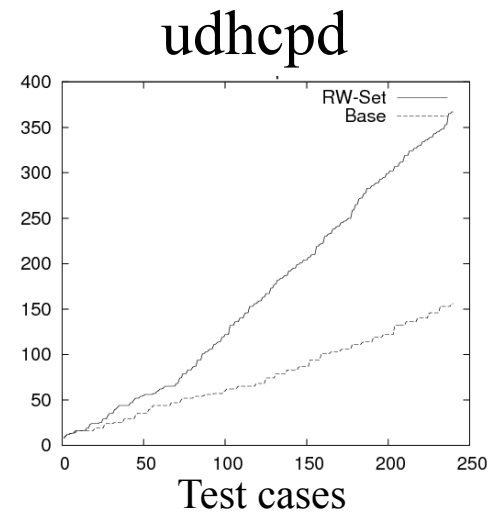
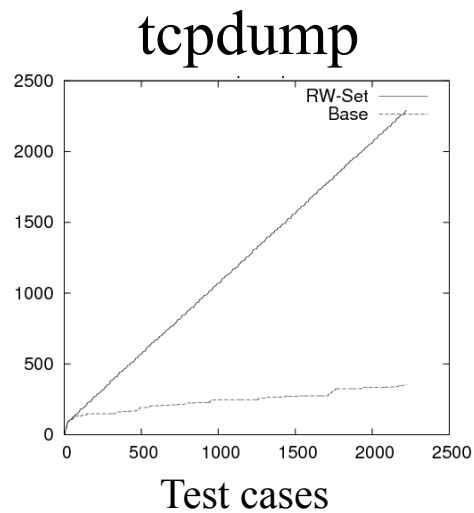
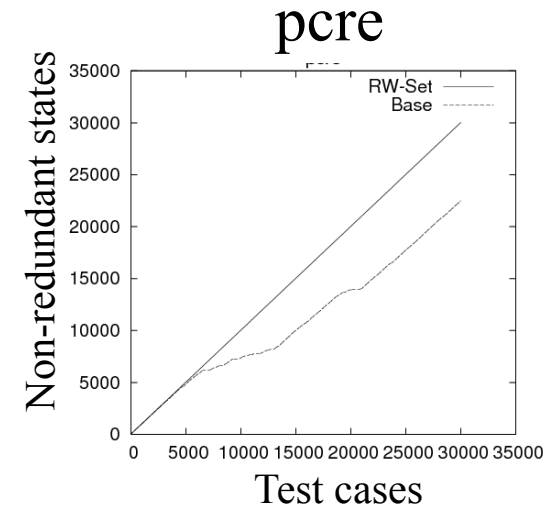
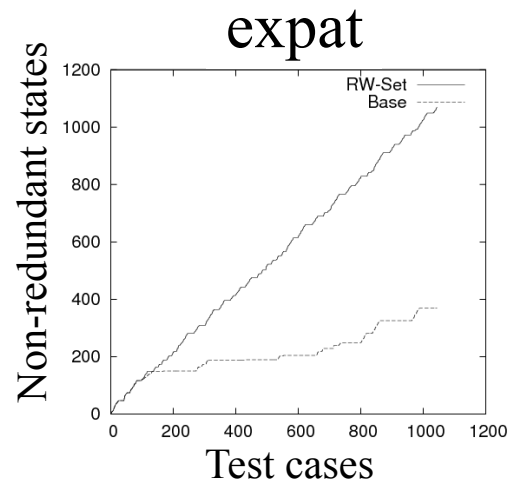
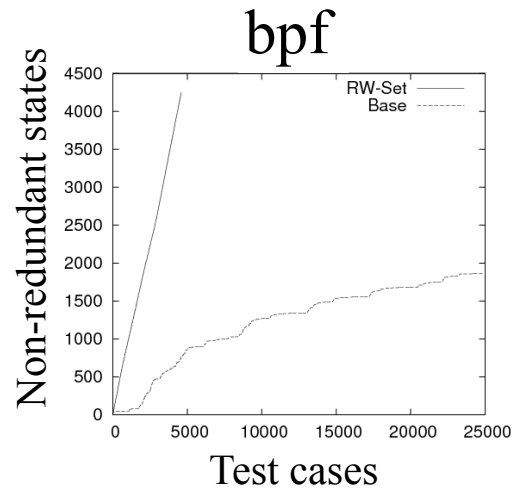


Medium-sized apps



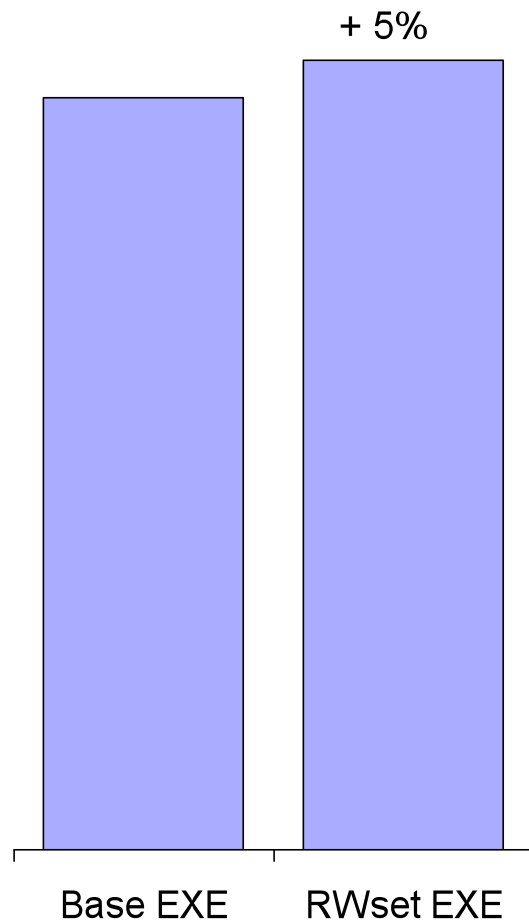


Non-redundant states





Performance



- Dry run for the medium-sized apps: compute write-sets and read-sets but do no pruning



Device drivers

- Hard to check w/o the physical device or outside of the kernel
- Focused on three MINIX 3 drivers:
 - lance: AMD lance ethernet card driver
 - pci: PCI bus driver
 - sb16: Sound Blaster 16 driver



Minix 3 device drivers

- Structured around a main dispatch loop

```
while (1) {  
    /* receive message from other processes,  
       the kernel, or the hardware */  
    message = read_message();  
    process(message);  
}
```



Minix 3 device drivers

- Structured around a main dispatch loop

```
while (1) {  
    /* receive message from other processes,  
       the kernel, or the hardware */  
    message = read_symbolic_data();  
    process(message);  
}
```




Minix 3 device drivers

- Structured around a main dispatch loop

```
for (k=0; k < n; k++) {  
    /* receive message from other processes,  
       the kernel, or the hardware */  
    message = read_symbolic_data();  
    process(message);  
}
```

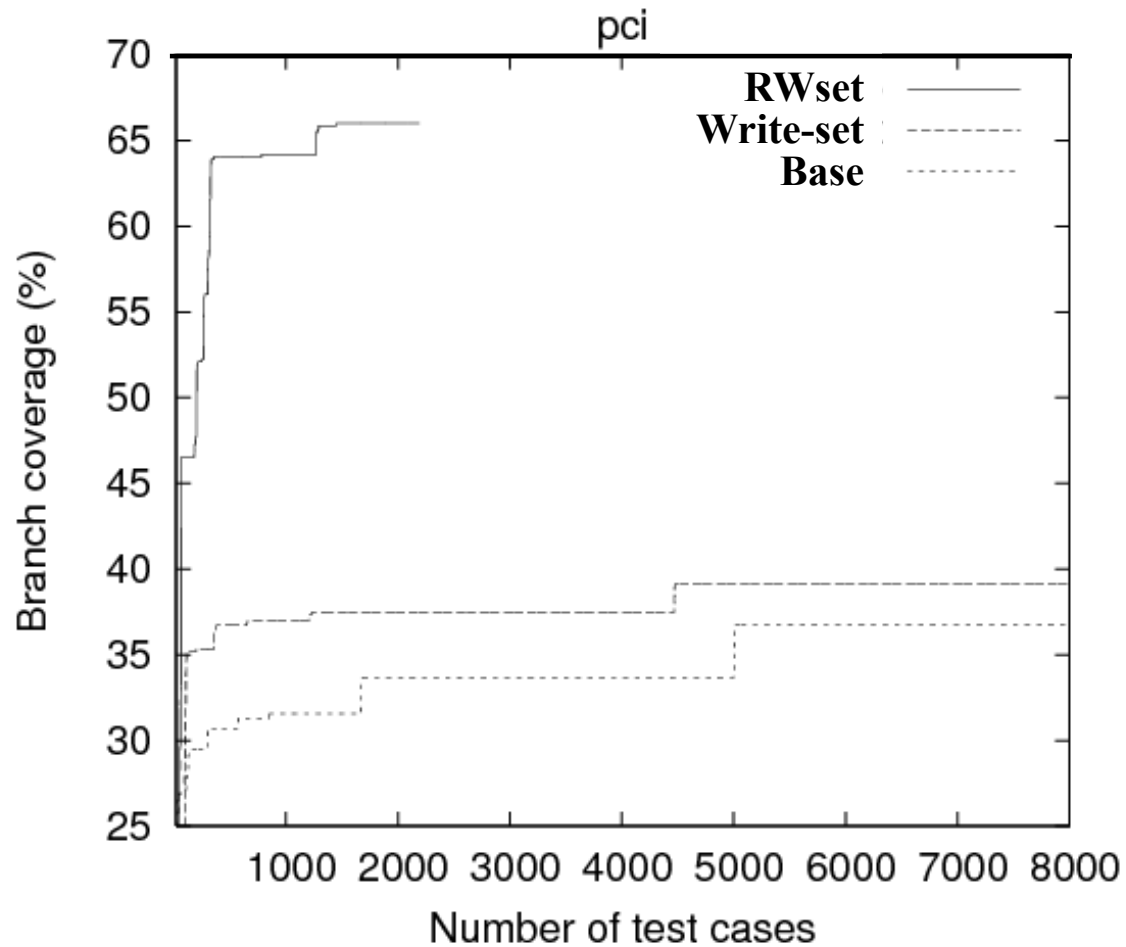


Minix drivers - evaluation

- Three versions of EXE:
 - Base
 - Write-set
 - RWset
- Fixed the # of iterations in each version
 - mainly to have the base version terminate
- Ran each version for an hour
 - recorded branch coverage + non-redundant states

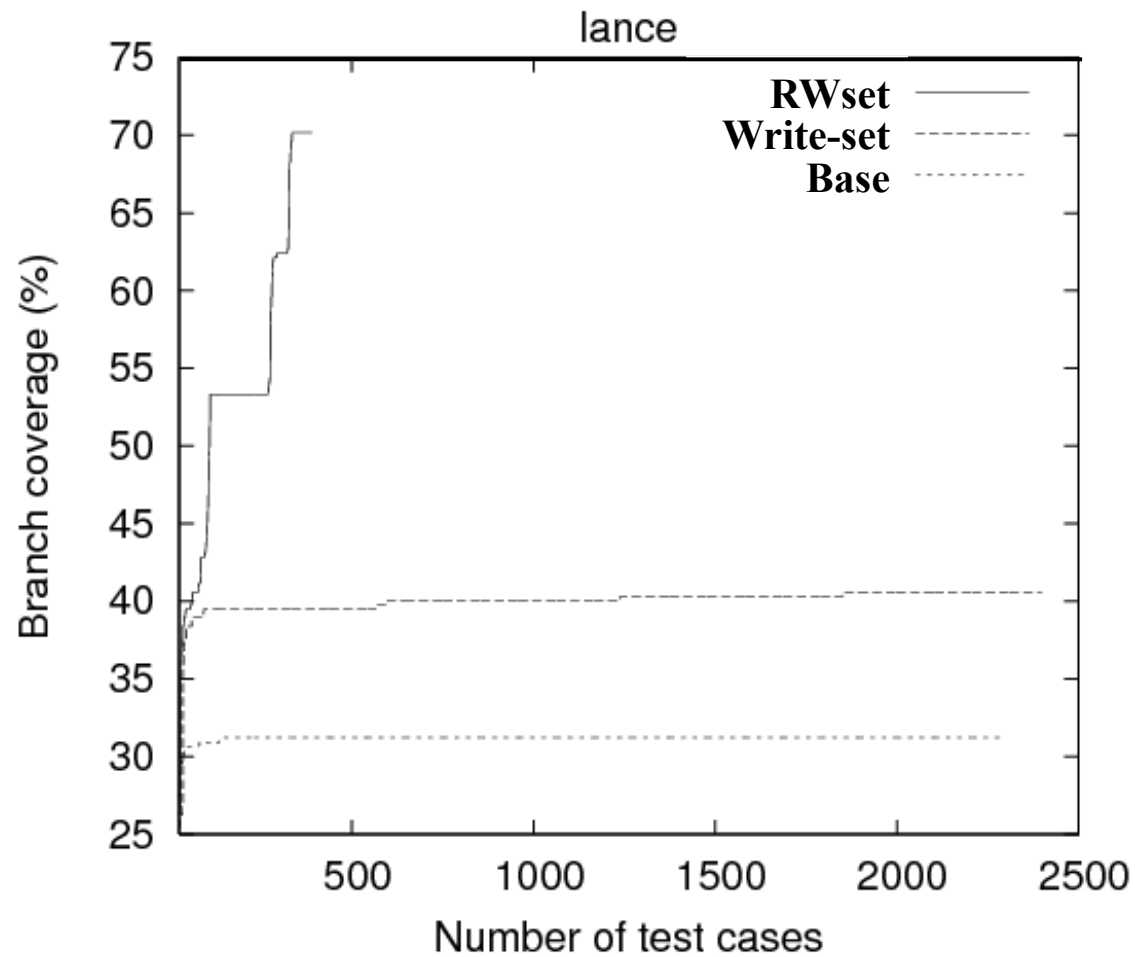


Branch coverage (pci)



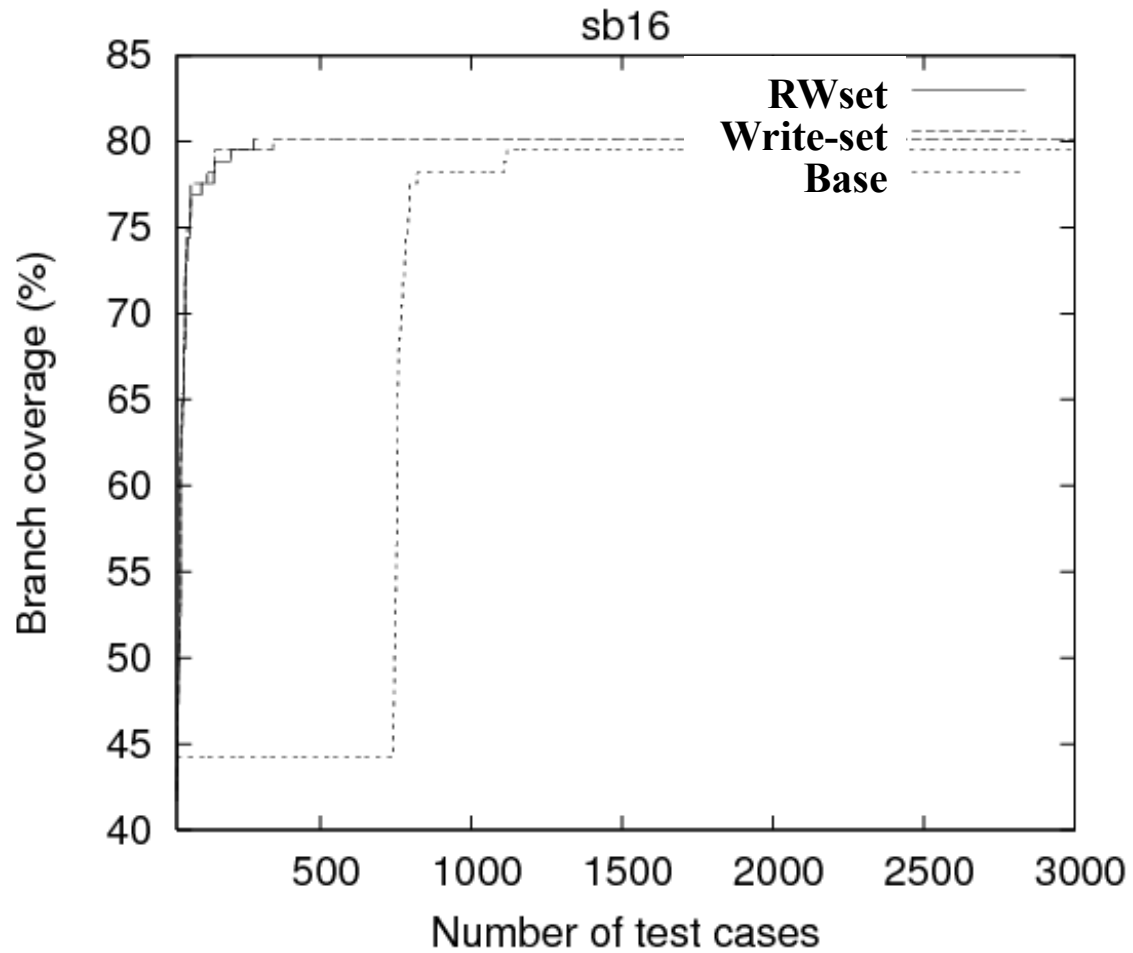


Branch coverage (lance)



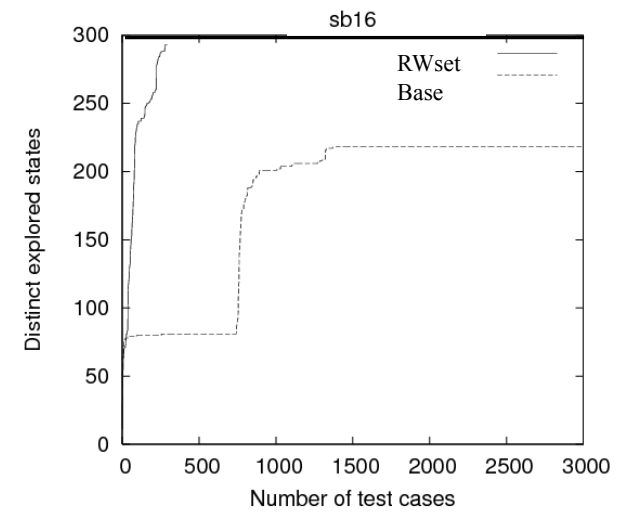
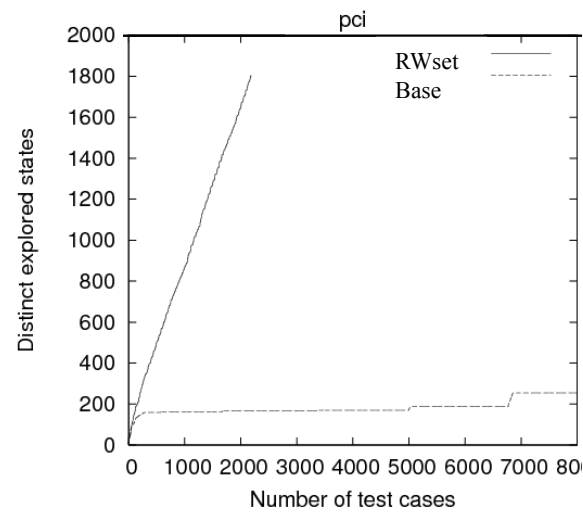
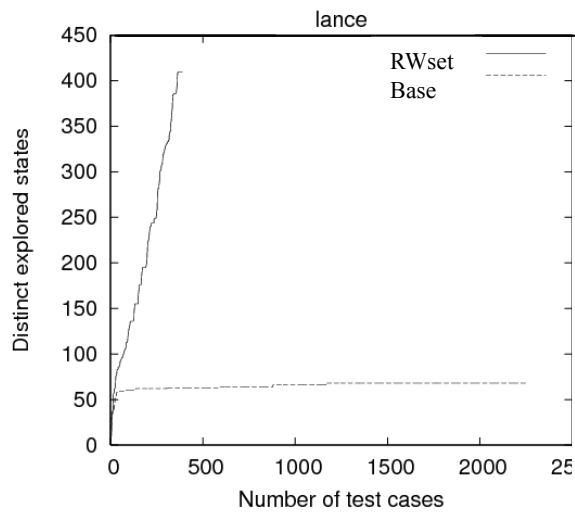


Branch coverage (sb16)



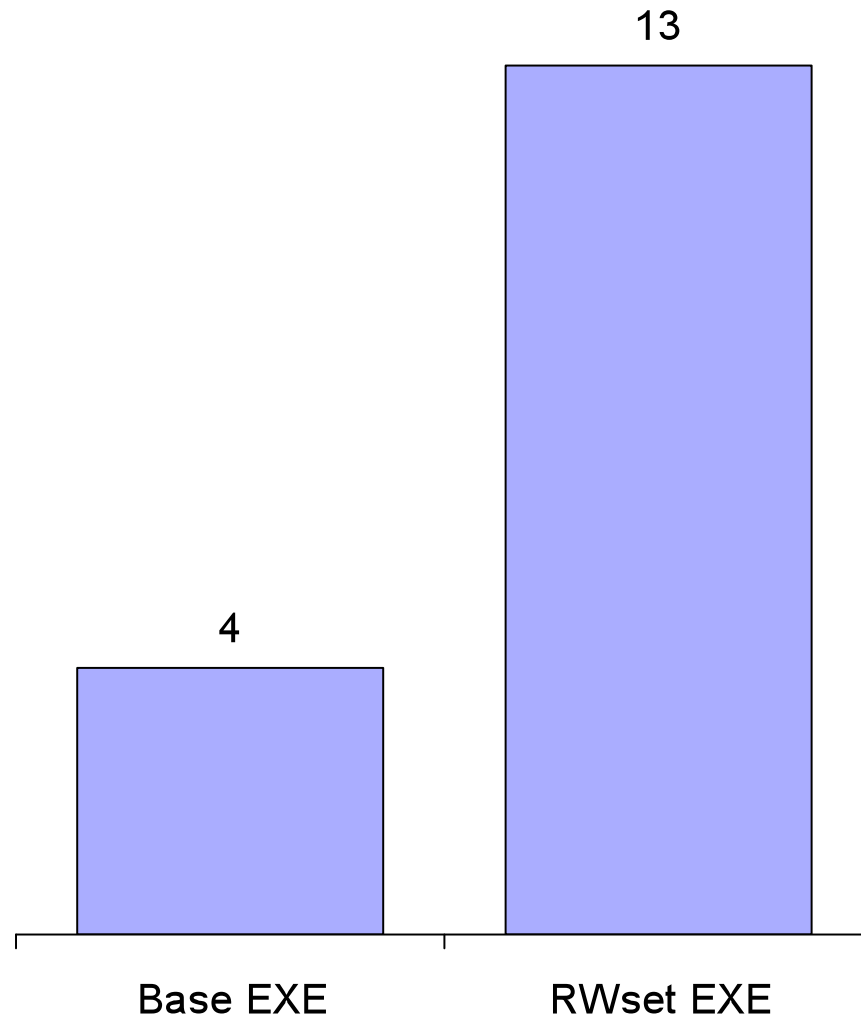


Non-redundant states





Bugs in device drivers





Summary

- RWset analysis
 - efficient way to prune large numbers of redundant paths
 - only values observed by the program trigger the execution of new paths
- Evaluated on Minix drivers and medium size applications
 - big reduction in the number of paths explored



Questions?