Imperial College London



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RDMA Tutorial

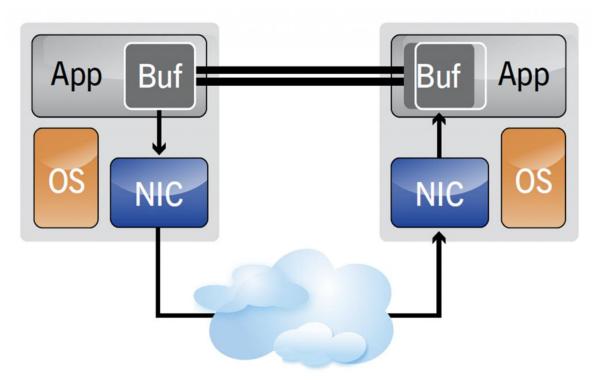
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Remote Direct Memory Access

RDMA is a hardware *mechanism* through which the network card (*NIC*) can *directly access* all or parts of the *main memory* of a *remote node without involving the processor*.



RDMA properties

Remote – data is transferred between nodes in a network
 Direct – no CPU or OS kernel is involved in the data transfer
 Memory – data transferred between two apps and their virtual address spaces
 Access – support to send, receive, read, write, and do atomic operations

Main highlights of RDMA

- Zero-copy data
- Bypasses the CPU
- Bypasses the OS kernel
- Message based transactions

Benefits of using RDMA

- High throughput (bandwidth)
- Low end-to-end latencies

Low CPU utilization

One-sided RDMA operations do not involve the remote CPU at all.

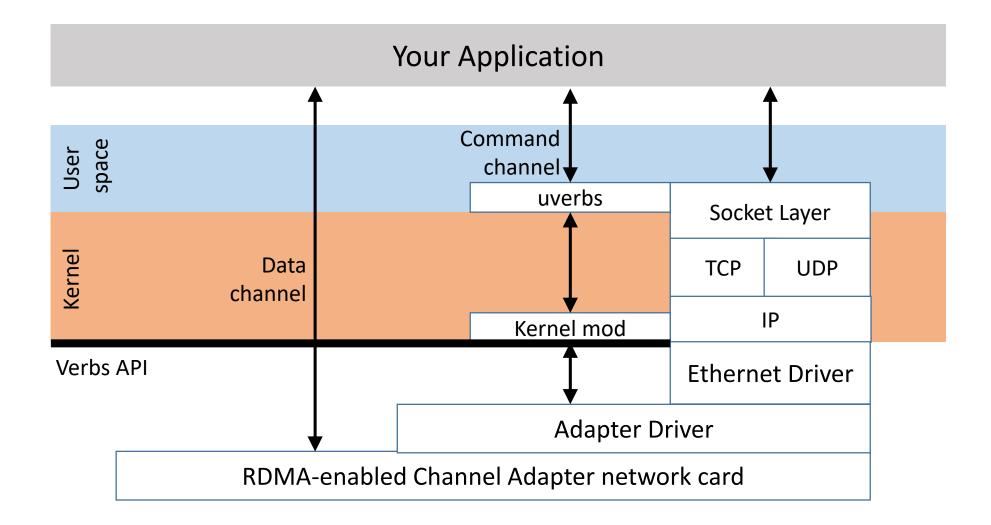
✓ Low memory bus contention

No data is copied between the user space and kernel, and the other way around.

Asynchronous operations

Great for overlapping communication and computation.

Traditional TCP/IP sockets vs RDMA



Setting up the RDMA data channels

Buffers need to be *registered* with the *network card* before used

During the registration process:

- *Pin memory* so that it cannot be swapped by the Operating System.
- Store the address translation information in the NIC.
- Set permissions for the memory region.
- Return a remote and local key, which are used by the adapters when executing the RDMA operations.

Work Queues

RDMA communication is based on a set of three queues

Send

work queues, always created as a Queue Pair (QP)

- Receive
- Completion

The send and receive queues are there to schedule the work to be done.

A completion queue is used to notify when the work has been completed.

Applications issue a job using a *work request* or a *work queue element*

A work request is a small *struct* with a *pointer to a buffer*:

- In a send queue it's a pointer to a message to be sent.
- In a receive queue it's shows where an incoming message should be placed.

Once a work request has been completed, the adapter creates a *completion queue element* and enqueues it in the *completion queue*.

RDMA's network stack overview

Application	 Posts work requests to a queue Each work request is a message, a unit of work
	 Verbs interface – allows the application to request services
RDMA adapter driver	 Maintains the work queues Manages address translation Provides completion and even mechanisms
RDMA-supporting NIC and network protocols	 Transport layer: reliable/unreliable, datagram, etc. Packetizes messages Implements the RDMA protocol Implements end-to-end reliability Assures reliable delivery

src: InfiniBand Trade Association: Introduction to IB for end users

Network protocols supporting RDMA

InfiniBand (IB)

- QDR 4x 32 Gbps
- FDR 4x 54 Gbps
- EDR 4x 100 Gbps

RoCE – RDMA over Converged Ethernet

- 10 Gbps
- 40 Gbps

iWARP – internet Wide Area RDMA Protocol

RDMA is just a *mechanism*

Does *not* specify the *semantics* of a data transfer

RDMA networks support two types of memory access models:

- One sided RDMA read and write + atomic operations
- Two sided RDMA send and receive

RDMA Send and Receive

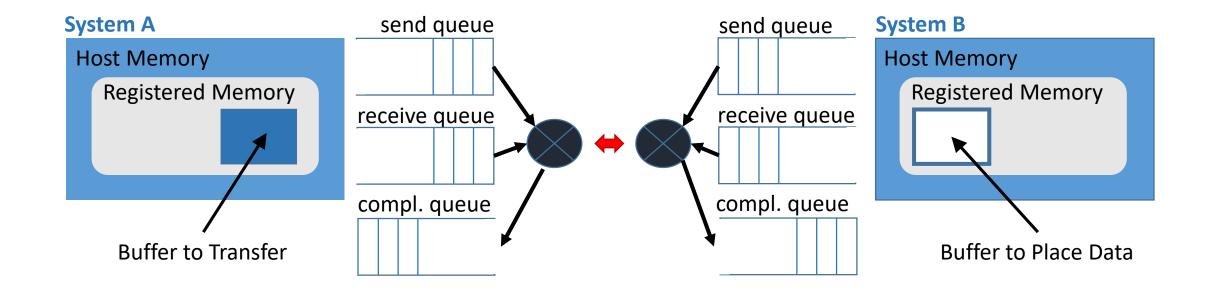
Traditional message passing where *both* the *source* and the *destination* processes are *active*ly involved in the communication.

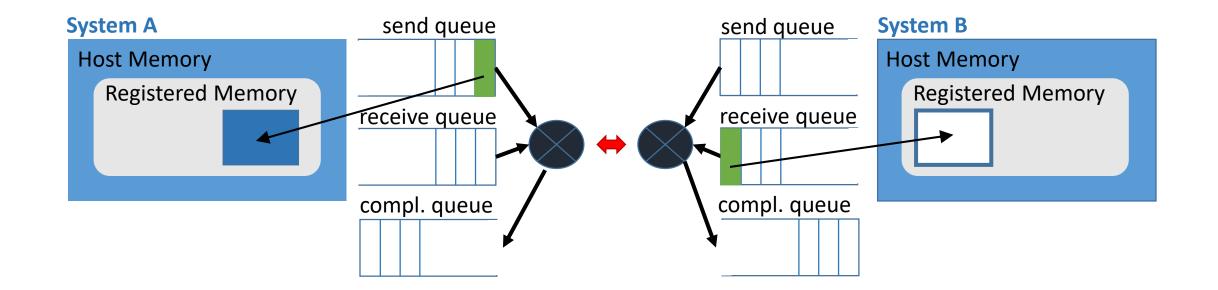
Both need to have **created** their queues:

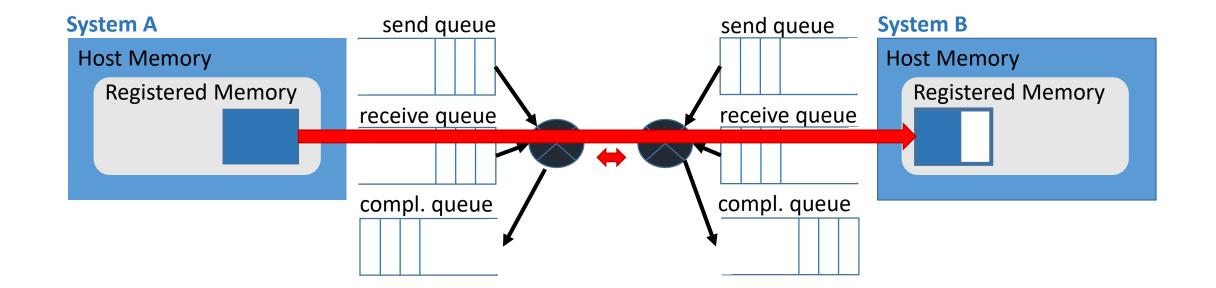
- A queue pair of a send and a receive queue.
- A *completion queue* for the queue pair.

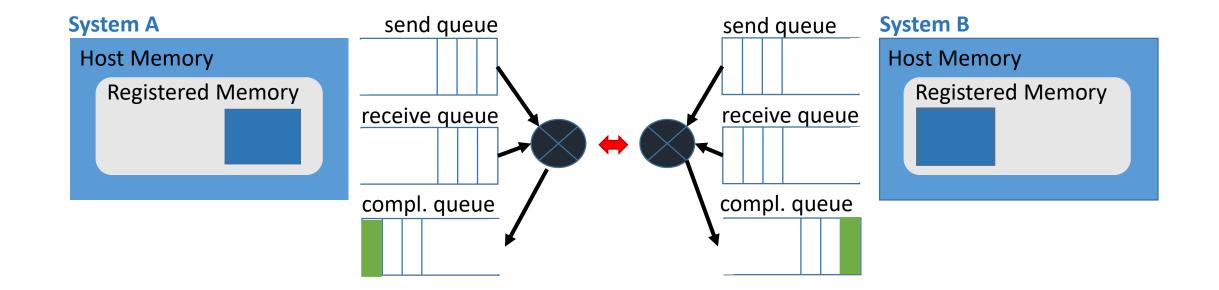
Sender's work request has a pointer to a buffer that it wants to send. The WQE is enqueued in the send queue.

Receiver's work request has a pointer to an empty buffer for receiving the message. The WQE is enqueued in the receive queue.









Only the *sender* side is *active*; the *receiver* is *passive*.

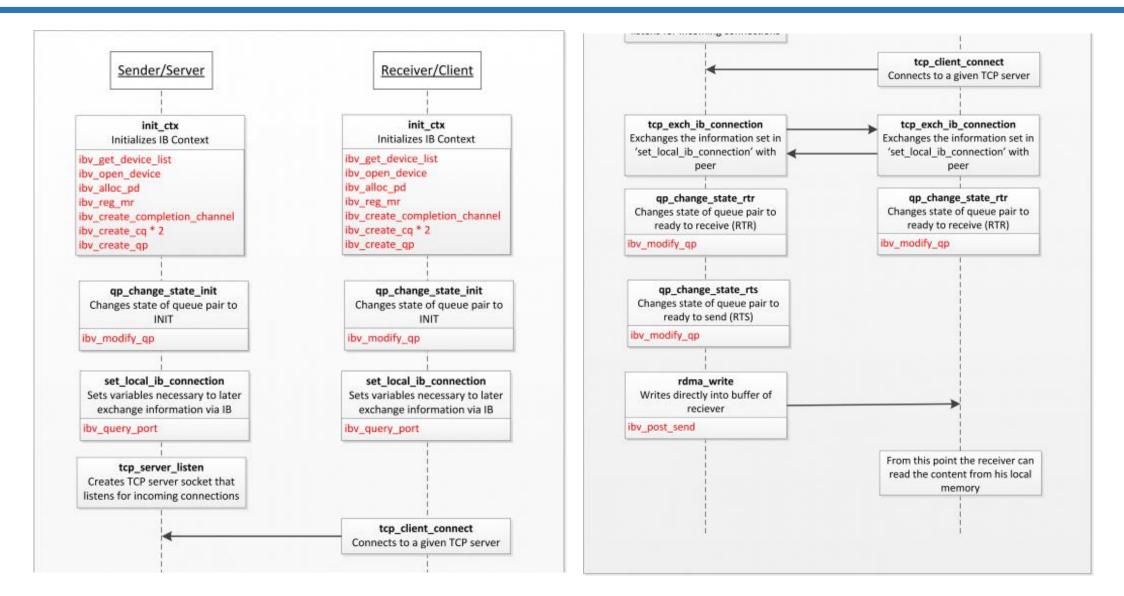
The passive side issues no operation, uses no CPU cycles, gets no indication that a "read" or a "write" happened.

To issue an RDMA *read* or a *write*, the work request *must include:*

- 1. the *remote* side's *virtual memory address* and
- 2. the *remote* side's *memory registration key*.

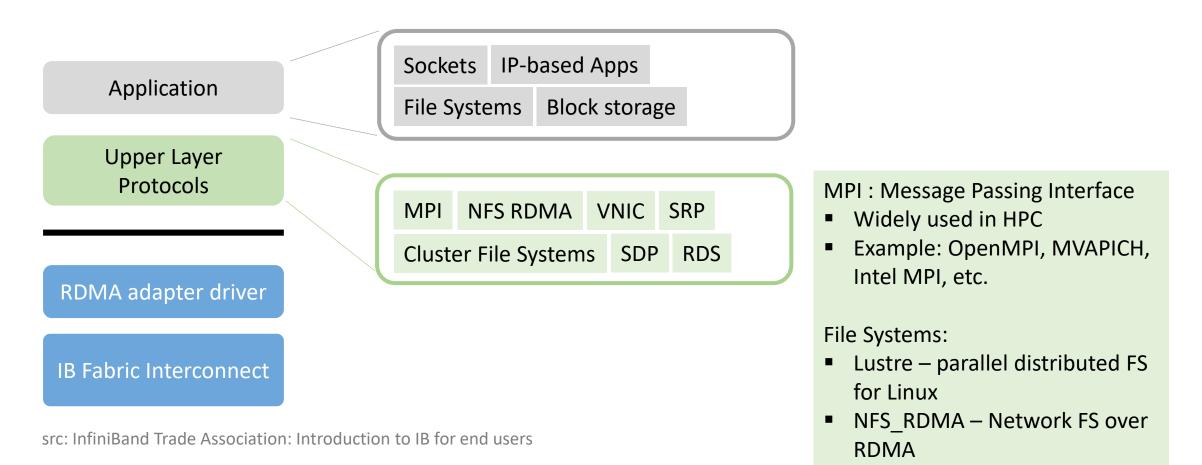
The active side must obtain the passive side's address and key beforehand. Typically, the traditional RDMA send/receive mechanisms are used.

Using the verbs API



Challenges of using RDMA

Added extra complexity for the developer to use the Verbs API



RDMA References

IB trade introduction <u>https://cw.infinibandta.org/document/dl/7268</u>

- First steps for programming with IB verbs <u>https://thegeekinthecorner.wordpress.com/2010/08/13/building-an-rdma-capable-application-with-ib-verbs-part-1-basics/</u>
- Figures from <u>https://zcopy.wordpress.com/category/getting-started/</u>
- More details at <u>http://www.mellanox.com/related-</u> <u>docs/prod_software/RDMA_Aware_Programming_user_manual.pdf</u>

Overview of our new EDR cluster

- EDR InfiniBand
- 36-port Mellanox switch
- 18 nodes cluster (EDR NICs)
- I server with 4 Xeon E5-5660 v4 processors:
 - 64 cores (128 with HT enabled)
 - 512 GB RAM
 - 2 EDR NICs, 1 x 10G NIC, 1 x 1G NIC

8 servers with 2 Xeon E5-2630 v4 processors:

- 20 cores (40 with HT enabled)
- 32 GB RAM
- 2 EDR NICs