

1 Syntax

$$\begin{array}{llll}
 e ::= c \in \mathbb{Z} & e_e ::= \cdot +_1 e & s \in stat ::= skip & s_e ::= x :=_1 \cdot \\
 | x \in Var & | \cdot +_2 \cdot & | s_1; s_2 & | \cdot ;_1 s_2 \\
 | e_1 + e_2 & | @_1(e_2) & | x := e & | if_1 s_1 s_2 \\
 | \lambda x.s & | @_2 & | if (e > 0) s_1 s_2 & | while_1 (e > 0) s \\
 | e_1(e_2) & | @_3 & | while (e > 0) s & | while_2 (e > 0) s \\
 | alloc & | .f & | return e & | return_1. \\
 | e.f & | f.in_1 & | e_1.f := e_2 & | .f :=_1 e_2 \\
 | f.in.e & & | delete e.f & | .f :=_2 \cdot \\
 & & & | delete_1 .f
 \end{array}$$

2 Abstract Semantics

2.1 Expressions

$$\begin{array}{c}
 \text{RED-CONST}(c) \\
 \hline
 \overline{(- \mid emp, \eta_e, \eta_c), c \Downarrow (- \mid emp, \eta_e, \alpha(c))} \\
 \\
 \text{RED-VAR-LOCAL}(x) \\
 \hline
 \overline{(- \mid \eta_e \mapsto E_e^\# \star \eta_c \mapsto E_c^\#, \eta_e, \eta_c), x \Downarrow (- \mid \eta_e \mapsto E_e^\# \star \eta_c \mapsto E_c^\#, \eta_e, E_c^\#[x])} \quad x \in \text{dom}(E_c^\#) \\
 \\
 \text{RED-VAR-LOCAL}(x) \\
 \hline
 \overline{(- \mid \eta \mapsto E^\#, \eta, \eta), x \Downarrow (- \mid \eta \mapsto E^\#, \eta, E^\#[x])} \quad x \in \text{dom}(E^\#) \\
 \\
 \text{RED-VAR-GLOBAL}(x) \\
 \hline
 \overline{(- \mid \eta_e \mapsto E_e^\# \star \eta_c \mapsto E_c^\#, \eta_e, \eta_c), x \Downarrow (- \mid \eta_e \mapsto E_e^\# \star \eta_c \mapsto E_c^\#, \eta_e, E_e^\#[x])} \quad x \in \text{dom}(E_e^\#) \wedge x \notin \text{dom}(E_c^\#) \\
 \\
 \text{RED-VAR-UNDEF}(x) \\
 \hline
 \overline{(- \mid \eta_e \mapsto E_e^\# \star \eta_c \mapsto E_c^\#, \eta_e, \eta_c), x \Downarrow (- \mid \eta_e \mapsto E_e^\# \star \eta_c \mapsto E_c^\#, err^\#)} \quad x \notin \text{dom}(E_e^\#) \wedge x \notin \text{dom}(E_c^\#) \\
 \\
 \text{RED-VAR-UNDEF}(x) \\
 \hline
 \overline{(- \mid \eta \mapsto E^\#, \eta, \eta), x \Downarrow (- \mid \eta \mapsto E^\#, err^\#)} \quad x \notin \text{dom}(E^\#)
 \end{array}$$

$$\begin{array}{c}
\text{RED-ADD}(e_1, e_2) \\
\frac{(- \mid \phi, \eta_e, \eta_c), e_1 \Downarrow (M \mid \phi', r^\sharp) \quad (M \mid \phi', M(\eta_c), r^\sharp), \cdot +_1 e_2 \Downarrow \Phi}{(- \mid \phi, \eta_e, \eta_c), e_1 + e_2 \Downarrow \Phi} \\
\\
\text{RED-ADD-1}(e_2) \\
\frac{(- \mid \phi, \eta_e, \eta_c), e_2 \Downarrow (M \mid \phi', r^\sharp) \quad (M \mid \phi', M(v_1^\sharp), r^\sharp), \cdot +_2 \cdot \Downarrow \Phi}{(- \mid \phi, \eta_e, (v_1^\sharp)), \cdot +_1 e_2 \Downarrow \Phi} \\
\\
\text{RED-ADD-2} \\
\frac{\text{RED-LAMBDA}(x, s)}{(- \mid emp, v_1^\sharp, (v_2^\sharp)), \cdot +_2 \cdot \Downarrow (- \mid emp, \eta_e, v_1^\sharp +^\sharp v_2^\sharp)} \\
\\
\text{RED-LAMBDA}(x, s) \\
\frac{\text{RED-APP}(e_1, e_2)}{(- \mid \phi, \eta_e, \eta_c), e_1 \Downarrow (M \mid \phi', r^\sharp) \quad (M \mid \phi', M(\eta_c), r^\sharp), @_1(e_2) \Downarrow \Phi} \\
\\
\text{RED-APP-1}(e_2) \\
\frac{(- \mid \phi, \eta_e, \eta_c), e_2 \Downarrow (M \mid \phi', r^\sharp) \quad (M \mid \phi', M(v_1^\sharp), x, s, r^\sharp), @_2 \Downarrow \Phi}{(- \mid \phi, \eta_c, (v_1^\sharp, @_2(e_2))) @_1(e_2) \Downarrow \Phi} \\
\\
\text{RED-APP-2}(s) \\
\frac{\text{RED-APP-3-RET}}{(- \mid \eta \mapsto E^\sharp[x \leftarrow v^\sharp] * \eta \mapsto E^\sharp * \phi, \eta_e, \eta_c), s \Downarrow \Phi \quad \Phi, @_3 \Downarrow \Phi'} \\
\\
\text{RED-APP-3-NO-RET} \\
\frac{\text{RED-APP-3-NO-RET}}{(- \mid emp, ret(\eta_e, v^\sharp)), @_3 \Downarrow (- \mid emp, \eta_e, v^\sharp)} \\
\\
\text{RED-APP-3-NO-RET} \\
\frac{\text{RED-APP-3-NO-RET}}{(- \mid emp, \eta_e, \eta_c), @_3 \Downarrow (- \mid emp, err)}
\end{array}$$

$$\begin{array}{c}
\text{RED-NEW-OBJ} \\
\hline
\overline{(- \mid emp, \eta_e, \eta_c) , alloc \Downarrow (-, \bullet \rightarrow l \mid l \mapsto \{_ : \boxtimes\}, \eta_e, l)}
\\
\\
\frac{\text{RED-FIELD}(e, f)}{(- \mid \phi, \eta_e, \eta_c) , e \Downarrow \Phi \quad \Phi, .f \Downarrow \Phi'} \\
\hline
\overline{(- \mid \phi, \eta_e, \eta_c) , e.f \Downarrow \Phi'}
\\
\\
\frac{\text{RED-FIELD-1}(f)}{(- \mid l \mapsto \{f : u^\# \}, \eta_e, \eta_c) .f \Downarrow (- \mid l \mapsto \{f : u^\# \}, \eta_e, u^\# \mid_{Val^\#})}
\\
\\
\frac{\text{RED-IN}(f, e)}{(- \mid \phi, \eta_e, \eta_c) , e \Downarrow \Phi \quad \Phi, f \text{ in } e \Downarrow \Phi'} \frac{\text{RED-IN-1-TRUE}(f)}{(- \mid l \mapsto \{f : u^\# \}, \eta_e, l) , f \text{ in } e \cdot \Downarrow (- \mid l \mapsto \{f : u^\# \}, \eta_e, +)} \quad u^\# \Big|_{Val^\#} \neq \perp
\\
\\
\frac{\text{RED-IN-1-FALSE}(f)}{(- \mid l \mapsto \{f : u^\# \}, \eta_e, l) , f \text{ in } e \cdot \Downarrow (- \mid l \mapsto \{f : u^\# \}, \eta_e, 0)} \quad \boxtimes \sqsubseteq u^\#
\end{array}$$

2.2 Statements

RED-SKIP

$$\overline{(- \mid emp, \eta_e, \eta_c) , skip \Downarrow (- \mid emp, \eta_e, \eta_c)}$$

$$\frac{\text{RED-SEQ}(s_1, s_2)}{(- \mid \phi, \eta_e, \eta_c) , s_1 \Downarrow \Phi \quad \Phi, \cdot;_1 s_2 \Downarrow \Phi'} \quad \frac{\text{RED-SEQ-1}(s_2)}{(- \mid \phi, \eta_e, \eta_c) , s_2 \Downarrow \Phi}$$

$$\begin{array}{c}
\text{RED-ASN}(x, e) \\
\frac{(- \mid \phi, \eta_e, \eta_c), e \Downarrow (M \mid \phi', x^\#) \quad (M \mid \phi', M(\eta_c), x^\#), x :=_1 \cdot \Downarrow \Phi}{(- \mid \phi, \eta_e, \eta_c), x := e \Downarrow \Phi}
\end{array}$$

$$\begin{array}{c}
\text{RED-ASN-1}(x) \\
\frac{\eta'_e \text{ fresh}}{(- \mid \eta_e \mapsto E_e^\# \star \eta_c \mapsto E_c^\#, \eta_c, (\eta_e, v^\#)), x :=_1 \cdot \Downarrow (-, \bullet \rightarrow \eta'_e \mid \eta_e \mapsto E_e^\# \star \eta_c \mapsto E_c^\# \star \eta'_e \mapsto E_e^\# [x \leftarrow v^\#], \eta'_e, \eta_c) \quad x \notin \text{dom}(E_c^\#)}
\end{array}$$

$$\begin{array}{c}
\text{RED-ASN-1}(x) \\
\frac{\eta_e \text{ fresh}}{(- \mid \eta \mapsto E^\#, \eta, (\eta, v^\#)), x :=_1 \cdot \Downarrow (-, \bullet \rightarrow \eta_e \mid \eta \mapsto E^\# \star \eta_e \mapsto E^\# [x \leftarrow v^\#], \eta_e, \eta) \quad x \notin \text{dom}(E^\#)}
\end{array}$$

$$\begin{array}{c}
\text{RED-ASN-1-LOCAL}(x) \\
\frac{\eta'_c \text{ fresh}}{(- \mid \eta_e \mapsto E_e^\# \star \eta_c \mapsto E_c^\#, \eta_c, (\eta_e, v^\#)), x :=_1 \cdot \Downarrow (-, \bullet \rightarrow \eta'_e \mid \eta_e \mapsto E_e^\# \star \eta_c \mapsto E_c^\# \star \eta'_c \mapsto E_c^\# [x \leftarrow v^\#], \eta_e, \eta'_c) \quad x \in \text{dom}(E_c^\#)}
\end{array}$$

$$\begin{array}{c}
\text{RED-ASN-1-LOCAL}(x) \\
\frac{\eta_c \text{ fresh}}{(- \mid \eta \mapsto E^\#, \eta, (\eta, v^\#)), x :=_1 \cdot \Downarrow (-, \bullet \rightarrow \eta_c \mid \eta \mapsto E^\# \star \eta_c \mapsto E^\# [x \leftarrow v^\#], \eta, \eta_c) \quad x \in \text{dom}(E^\#)}
\end{array}$$

$$\begin{array}{c}
\text{RED-IF}(e, s_1, s_2) \\
\frac{(- \mid \phi, \eta_e, \eta_c), e \Downarrow (M \mid \phi', x^\#) \quad (M \mid \phi', M(\eta_c), x^\#), \text{if}_1 s_1 s_2 \Downarrow \Phi}{(- \mid \phi, \eta_e, \eta_c), \text{if } (e > 0) s_1 s_2 \Downarrow \Phi}
\end{array}$$

$$\begin{array}{c}
\text{RED-IF-1-POS}(s_1, s_2) \\
\frac{(- \mid \phi, \eta_c, \eta_e), s_1 \Downarrow \Phi}{(- \mid \phi, \eta_c, (\eta_e, v^\#)), \text{if}_1 s_1 s_2 \Downarrow \Phi} \quad v^\# \sqcap + \neq \perp
\end{array}
\quad
\begin{array}{c}
\text{RED-IF-1-NEG}(s_1, s_2) \\
\frac{(- \mid \phi, \eta_c, \eta_e), s_2 \Downarrow \Phi}{(- \mid \phi, \eta_c, (\eta_e, v^\#)), \text{if}_1 s_1 s_2 \Downarrow \Phi} \quad v^\# \sqcap -_0 \neq \perp
\end{array}$$

$$\begin{array}{c}
\text{RED-WHILE}(e, s) \\
\frac{(- \mid \phi, \eta_e, \eta_c), e \Downarrow (M \mid \phi', x^\#) \quad (M \mid \phi', M(\eta_c), x^\#), \text{while}_1 (e > 0) s \Downarrow \Phi}{(- \mid \phi, \eta_e, \eta_c), \text{while } (e > 0) s \Downarrow \Phi}
\end{array}$$

$$\begin{array}{c}
\text{RED-WHILE-1-NEG}(e, s) \\
\frac{}{(- \mid \phi, \eta_c, (\eta_e, v^\#)), \text{while}_1 (e > 0) s \Downarrow (- \mid \phi, \eta_e, \eta_c)} \quad v^\# \sqcap -_0 \neq \perp
\end{array}$$

$$\begin{array}{c}
\text{RED-WHILE-1-POS}(e, s) \\
\frac{(- \mid \phi, \eta_e, \eta_c), s \Downarrow \Phi \quad \Phi, \text{while}_2 (e > 0) s \Downarrow \Phi'}{(- \mid \phi, \eta_c, (\eta_e, v^\#)), \text{while}_1 (e > 0) s \Downarrow \Phi'} \quad v^\# \sqcap + \neq \perp
\end{array}
\quad
\begin{array}{c}
\text{RED-WHILE-2}(e, s) \\
\frac{(- \mid \phi, \eta_e, \eta_c), \text{while } (e > 0) s \Downarrow \Phi}{(- \mid \phi, \eta_e, \eta_c), \text{while}_2 (e > 0) s \Downarrow \Phi}
\end{array}$$

$$\begin{array}{c}
\text{RED-RETURN}(e) \\
\frac{}{(- \mid \phi, \eta_e, \eta_c), e \Downarrow \Phi \quad \Phi, return_1 \cdot \Downarrow \Phi'} \\
\text{RED-RETURN-1} \\
\frac{}{(- \mid emp, \eta_e, v^\#), return_1 \cdot \Downarrow (- \mid emp, ret(\eta_e, v^\#))} \\
\text{RED-FIELD-ASN}(e_1, f, e_2) \\
\frac{(- \mid \phi, \eta_e, \eta_c), e_1 \Downarrow (M \mid \phi', r^\#) \quad (M \mid \phi', M(\eta_c), r^\#), .f :=_1 e_2 \Downarrow \Phi}{(- \mid \phi, \eta_e, \eta_c), e_1.f := e_2 \Downarrow \Phi} \\
\text{RED-FIELD-ASN-1}(f, e_2) \\
\frac{(- \mid \phi, \eta_e, \eta_c), e_2 \Downarrow (M \mid \phi', r^\#) \quad (M \mid \phi', M(\eta_c), M(l), r^\#), .f :=_2 \cdot \Downarrow \Phi}{(- \mid \phi, \eta_c, (\eta_e, l)), .f :=_1 e_2 \Downarrow \Phi} \\
\text{RED-FIELD-ASN-2}(f) \\
\frac{}{(- \mid l \mapsto \{f : u^\#\}, \eta_c, l, (\eta_e, v^\#)), .f :=_2 \cdot \Downarrow (- \mid l \mapsto \{f : v^\#\}, \eta_e, \eta_c)} \\
\text{RED-DELETE}(e, f) \\
\frac{(- \mid \phi, \eta_e, \eta_c), e \Downarrow (M \mid \phi', r^\#) \quad (M \mid \phi', M(\eta_c), r^\#), delete_1.f \Downarrow \Phi}{(- \mid \phi, \eta_e, \eta_c), delete e.f \Downarrow \Phi} \\
\text{RED-DELETE-1}(f) \\
\frac{}{(- \mid l \mapsto \{f : u^\#\}, \eta_c, (l, \eta_e)), delete_1.f \Downarrow (- \mid l \mapsto \{f : \square\}, \eta_c, \eta_e)}
\end{array}$$

2.3 Aborting Rules

$$\begin{array}{ccc}
\text{RED-ERROR-EXPR}(e) & & \text{RED-ERROR-STAT}(s) \\
\frac{}{\Phi, e \Downarrow \Phi} \quad \text{abort } \Phi \wedge \neg \text{intercept}_e \Phi & & \frac{}{\Phi, s \Downarrow \Phi} \quad \text{abort } \Phi \\
& & \\
\frac{x^\# = C[err^\#]}{\text{abort } (M \mid \phi, x^\#)} & & \frac{}{\text{intercept}_{\text{@}_3} (M \mid \phi, ret(\eta_e, v^\#))}
\end{array}$$