${\sf Querying}$

Holger Pirk

Some housekeeping first

We've release the first coursework today

- it is due in exactly two weeks (November 3)
- the focus is on ORM, SQL and storage

You will need to use a VM on the cloudstack cluster

- Create it using the web interface
- Log in using ssh
- You can develop on the VM or locally
- Solutions will be marked against the VM

A brief demo

How to set up a VM

Supervised labtime

- Next Wednesday and Friday at 1pm in 202/206
- C++ is a b***h, let us help you

Haven't provided answers for Monday's calculations

- Sorry
- Will do it soon

What you know:

- How to design redundancy-free schemas
- How to get data into the DBMS using ORMs
- How data is stored internaly

Today you will learn how to get it out

- Relational Algebra
- Basic SQL
 - Enough to get you through the coursework
 - We will not talk about joins

The Running Example

| Сι | istor | ner | | |
|----|-------|--------|------------------|--|
| | ١D | Name | ShippingAddress | |
| | 1 | Holger | 180 Queens Gate | |
| | 2 | Sam | 32 Vassar Street | |
| | 3 | Peter | 180 Queens Gate | |

| Order | | |
|-------------------------|---------------------------|---|
| <u>D</u> 1 2 3 | Customer D 1 2 3 | - |

| OrderedItem | | | |
|-------------|--------|--|--|
| Orderld | BookID | | |
| 1 | 1 | | |
| 1 | 2 | | |
| 2 | 1 | | |
| 3 | 3 | | |
| | | | |

ь.

| Bool | < | |
|------|-----------------------------|-----------------------|
| D | Title | Author |
| 1 | Database Management Systems | Ramakrishnan & Gehrke |
| 2 | A Game of Thrones | Martin |
| 3 | Distributed Systems | van Steen & Tanenbaum |

Relational Algebra

A logical representation of relational operations

- Used to define the semantics of operations
- Used for logical optimization
- Not actually all that useful for end-users

Fundamentals

Nomenclatures

- An expression in relational algebra is composed from operators
- I will often refer to an expression as a plan
- Sometimes as a logical plan
- Cardinality is the number of tuples in a set

Handy properties

- Set-based
 - Order-invariant and duplicate eliminated
- Relational algebra is closed
 - Every operator produces a relation as output
 - Every operator accepts one or two relations as input
 - This simplifies the composition of operators into expressions
 - Note, that expressions can still be invalid

Relational Operators

A minimal set

- Select
- Project
- Cross (Carthesian) Product
- Union
- Difference

Not included

Intersection

$\mathsf{Project}\ \pi$

Intuitive semantics

- Extract one or multiple attributes from a relation
- Preserve relational semantics
- Changes the schema

| Customer | |
|----------|--------------------|
| D Nam | e ShippingAddress |
| 1 Holg | er 180 Queens Gate |
| 2 Sam | 32 Vassar Street |
| 3 Peter | 180 Queens Gate |

| $\pi_{ShippingAddress}$ Customer |
|--|
| ShippingAddress 180 Queens Gate 32 Vassar Street |

Quick Quiz: What is the cardinality of the output of a projection

- It can only be determined by evaluating it
- Cardinality of the input
- I don't know

Quick Quiz: What is the cardinality of the output of a projection

- It can only be determined by evaluating it
- Cardinality of the input
- I don't know

Quick Quiz: What is the upper bound for the cardinality of the output of a project

- It can only be determined by evaluating it
- Cardinality of the input
- I don't know

Intuitive Semantics

- Produce a new relation containing input tuples that satisfy a condition
- Does not change the schema
- Changes cardinality (i.e., the number of tuples in a relation)

| OrderedItem | | | |
|-------------|--------|-----|--|
| Orderld | BookID | | |
| 1 | 1 | - 1 | |
| 1 | 2 | - 1 | |
| 2 | 1 | - 1 | |
| 3 | 3 | - 1 | |
| | | _ | |

| $\sigma_{BookID=1} OrderedItem$ | | | | |
|---------------------------------|---------|--------|---|--|
| | Orderld | BookID | | |
| | 1 | 1 | | |
| | 2 | 1 | | |
| _ | | | _ | |

Quick Quiz: What is the cardinality of the output of a selection

- It can only be determined by evaluating it
- Cardinality of the input
- I don't know

Quick Quiz: What is the cardinality of the output of a selection

- It can only be determined by evaluating it
- Cardinality of the input
- I don't know

Quick Quiz: What is the upper bound for the cardinality of the output of a select

- It can only be determined by evaluating it
- Cardinality of the input
- I don't know

Cross (Carthesian) Product imes



| OrderedItem | | | |
|-------------|--------|--|--|
| Orderld | BookID | | |
| 1 | 1 | | |
| 1 | 2 | | |
| 2 | 1 | | |
| 3 | 3 | | |
| | | | |

| $Order \times Ord$ | eredi | ltem |
|--------------------|-------|------|
|--------------------|-------|------|

| ID | Customer D | Orderld | BookID |
|----|--------------|---------|--------|
| 1 | 1 | 1 | 1 |
| 1 | 1 | 1 | 2 |
| 1 | 1 | 2 | 1 |
| 1 | 1 | 3 | 3 |
| 2 | 2 | 1 | 1 |
| 2 | 2 | 1 | 2 |
| 2 | 2 | 2 | 1 |
| 2 | 2 | 3 | 3 |
| 3 | 3 | 1 | 1 |
| 3 | 3 | 1 | 2 |
| 3 | 3 | 2 | 1 |
| 3 | 3 | 3 | 3 |

Cross (Carthesian) Product imes

Intuitive Semantics

- Takes two inputs
- Produce a new relation by combining every tuple from the left with every tuple from the right
- Changes the schema

Cross (Carthesian) Product imes

Intuitive Semantics

- Takes two inputs
- Produce a new relation by combining every tuple from the left with every tuple from the right
- Changes the schema

Quick Quiz: What is the cardinality of the output of a cross product

- It can only be determined by evaluating it
- Cardinality of left input plus cardinality of right input
- Cardinality of left input times cardinality of right input
- I don't know

Rules

- Since relational algebra is closed, operators can be combined as long as their signature is respected
 - Cross products take two inputs
 - Selections and Projections take one

Example

 $\pi_{BookID}(\sigma_{ID==OrderID}(Order \times OrderedItem))$

Worksheet

| Customer | | | | | |
|----------|-----|--------|------------------|--|--|
| | I D | Name | ShippingAddress | | |
| | 1 | Holger | 180 Queens Gate | | |
| | 2 | Sam | 32 Vassar Street | | |
| | 3 | Peter | 180 Queens Gate | | |
| | | | | | |

| Order | |
|--------------|------------|
| <u>ID</u> Cu | 1 stomer D |
| 1 | 1 |
| 2 | 2 |

| OrderedItem | Book | |
|--|---|--|
| OrderId BookID 1 1 1 2 2 1 3 3 | IDTitle1Database Management Systems2A Game of Thrones3Distributed Systems | Author Ramakrishnan & Gehrke Martin van Steen & Tanenbaum |

The task

• From the operators we have discussed, compose a query that finds the titles of all books I (holger) have ordered.

Worksheet

| Cu | istor | ner | |
|----|-------|--------|------------------|
| | ID | Name | ShippingAddress |
| | 1 | Holger | 180 Queens Gate |
| | 2 | Sam | 32 Vassar Street |
| | 3 | Peter | 180 Queens Gate |
| | | | |

| Order | | |
|--------------------------|---------------------------|--|
| <u>ID</u> 1 2 3 | Customer D 1 2 3 | |

| OrderedItem | Book | |
|--|---|--|
| OrderId BookID 1 1 1 2 2 1 3 3 | ID Title 1 Database Management Sy 2 A Game of Thrones Systems 3 Distributed Systems | Author stems Ramakrishnan & Gehrke Martin van Steen & Tanenbaum |

The task

• From the operators we have discussed, compose a query that finds the titles of all books I (holger) have ordered.

Congratulations

• You have performed a join

$\mathsf{Union}\ \cup$

Intuitive Semantics

- Produce a new relation from two relations containing any tuple that is present in one of the inputs
- Does not change the schema
 - Requires union compatibility
- Changes cardinality (i.e., the number of tuples in a relation)

| ⁷ Name=Holger | Customer | $\sigma_{Name=Sam}Cust$ | omer |
|--------------------------|---|-------------------------|-----------------------------------|
| D Name 1 Holger | ShippingAddress 180 Queens Gate | D Name S 2 Sam 3 | hippingAddress 2 Vassar Street |
| | $\sigma_{Name=Sam}$ Custome $\sigma_{Name=Holger}$ Custom | r U ner | |
| | D Name SI | ipping Address | |

32 Vassar Street

2

Sam

Quick Quiz: What is the cardinality of the output of a union

- It can only be determined by evaluating it
- Cardinality of left input plus cardinality of right input
- Cardinality of left input times cardinality of right input
- I don't know

Quick Quiz: What is the cardinality of the output of a union

- It can only be determined by evaluating it
- Cardinality of left input plus cardinality of right input
- Cardinality of left input times cardinality of right input
- I don't know

Quick Quiz: What is the upper bound for the cardinality of a union

- It can only be determined by evaluating it
- Cardinality of left input plus cardinality of right input
- Cardinality of left input times cardinality of right input
- I don't know

Difference -

Intuitive Semantics

- Produce a new relation from two relations containing any tuple that is present in the first but not the second input
- Does not change the schema
 - Requires union compatibility
- Changes cardinality (i.e., the number of tuples in a relation)

| Сι | istoi | mer | |
|----|-------|--------|------------------|
| | ID | Name | ShippingAddress |
| | 1 | Holger | 180 Queens Gate |
| | 2 | Sam | 32 Vassar Street |
| | 3 | Peter | 180 Queens Gate |
| | | | |

| $\sigma_{Name=Sam}$ Customer | | | |
|------------------------------|------|------------------|--|
| ١D | Name | ShippingAddress | |
| 2 | Sam | 32 Vassar Street | |

| $Customer - \sigma_{Name=Sam}Customer$ | | | |
|--|---|--------|-----------------|
| | D | Name | ShippingAddress |
| | 1 | Holger | 180 Queens Gate |
| | 3 | Peter | 180 Queens Gate |

The Structured Query Language

Purpose

- Not a faithful implementation of relational algebra
- Meant to be useful
- Unfortunately, the S is not for "standard"

A language people like to hate

• An anecdote about SQL

Our example query

| Cu | istor | ner | |
|----|-------|--------|------------------|
| | ID | Name | ShippingAddress |
| | 1 | Holger | 180 Queens Gate |
| | 2 | Sam | 32 Vassar Street |
| | 3 | Peter | 180 Queens Gate |
| | | | |

| Order | |
|-----------|------------|
| <u>ID</u> | CustomerID |
| 1 | 1 |
| 2 | 2 |
| 3 | 3 |

| OrderedItem | Book | |
|--|--|--|
| OrderId BookID 1 1 1 2 2 1 3 3 | ID Title 1 Database Management Systems 2 A Game of Thrones 3 Distributed Systems | Author Ramakrishnan & Gehrke Martin van Steen & Tanenbaum |

A mildly complex query

```
select book.title as t from customer, book, order,
orderedItem
where customer.name = "holger"
and book.id = orderedItem.BookID
and order.id = orderedItem.orderID
and customer.id = order.customerId
Holger Pirk Querying
```

Mapping select book.title as t from customer, book, order, orderedItem where customer.name = "holger" and book.id = orderedItem.BookID and order.id = orderedItem.orderID and customer.id = order.customerId

Mapping

select book.title as t from customer, book, order, orderedItem where customer.name = "holger" and book.id = orderedItem.BookID and order.id = orderedItem.orderID and customer.id = order.customerId

```
Mapping
select book.title as t from customer, book, order,
orderedItem
where customer.name = "holger"
and book.id = orderedItem.BookID
and order.id = orderedItem.orderID
and customer.id = order.customerId
```

Set Operations in SQL

Union in SQL

select 1 union select 2 union select 3

Difference in SQL

(select 1 union select 2 union select 3) except select 2

Intersect in SQL

```
(select 1 union select 2 union select 3)
intersect (select 1 union select 2)
```

Bag vs. Set semantics in SQL

Attention

- Tables in SQL implement bag semantics
- select implements multisets/bag

Customer

| D | Name | ShippingAddress |
|---|--------|------------------|
| 1 | Holger | 180 Queens Gate |
| 2 | Sam | 32 Vassar Street |
| 3 | Peter | 180 Queens Gate |
| | | |



| Other operators have other defaults | | | | |
|-------------------------------------|---------|---------------------|--|--|
| Operator | Default | Switching behaviour | | |
| Select | Bag | select disctinct | | |
| Union | Set | union all | | |
| Except | Set | except all | | |
| Intersect | Set | intersect all | | |

Tables can be renamed while querying

select ShippingAddress from Customer as c

| Customer | | | | |
|----------|---|--------|------------------|--|
| | D | Name | ShippingAddress | |
| | 1 | Holger | 180 Queens Gate | |
| | 2 | Sam | 32 Vassar Street | |
| | 3 | Peter | 180 Queens Gate | |

This is particularly important when crossing with itself

```
select c1.Name, c2.Name
from Customer as c1, Customer as c2
where c1.ShippingAddress = c2.ShippingAddress
and c1.Name <> c2.Name
```

| Cu | istor | ner | |
|----|-------|--------|------------------|
| | ID | Name | ShippingAddress |
| | 1 | Holger | 180 Queens Gate |
| | 2 Sam | | 32 Vassar Street |
| | 3 | Peter | 180 Queens Gate |
| | | | |

Quiz: What does a null-value mean

- Unknown We don't know the value of this attribute, may not even have one
- Missing The value exists but has not been entered into the system
- Uncertain The value exists but I have conflicting information about what it is
 - Invalid The value that was entered in the system does not conform to the rules

Not an option This column is here for technical reasons but is not to be set

Null-Semantics in SQL

Definition fuzzy

- SQL uses a three valued logic to process WHERE predicate
- Truth values are TRUE, FALSE, and UNKNOWN
- SQL standard vague, but handling of NULL is roughtly equivalent to "Unknown"

| Truth values of SQL Formulae | | | | |
|------------------------------|---|--|--|--|
| Formula | Result | | | |
| x =null | UNKNOWN | | | |
| null=null | UNKNOWN | | | |
| x IS NULL | TRUE if x has a null value, FALSE otherwise | | | |
| x IS NOT NULL | TRUE if x does not have a null value, FALSE otherwise | | | |

Expression on null values produce null values

select c from (select null as c) as d where (1/c) is null

Aggregate functions - the five ansi SQL ones

• Sum, Count, Min, Max, Avg



Min, Max, Avg, Sum

- Null values do not contribute
- They are simply ignored

Count is weird

count(attribute) null values are ignored

count(*) null values are included in count

| Cι | istor | ner | | |
|----|-------|--------|------------------|--|
| | ID | Name | ShippingAddress | |
| | 1 | Holger | 180 Queens Gate | |
| | 2 Sam | | 32 Vassar Street | |
| | 3 | Peter | 180 Queens Gate | |
| | | | | |

select count(*) as c
from Customer group by
ShippingAddress

2

select max(len(Name)) as
length from Customer group by
ShippingAddress

| length | |
|--------|--|
| 6 | |
| 3 | |

Rule

• Null values form a group of their own

| Customer | | | | |
|----------|-----------------------|--|---|--|
| | D 1 2 3 3 | Name Holger Sam Peter Daniel | ShippingAddress 180 Queens Gate 32 Vassar Street 180 Queens Gate NULL | |

```
select max(len(Name))
as l from Customer
group by
ShippingAddress
```

| Question | | | | |
|-----------------------------|---------------------------------------|------------------------------|----------|--|
| What is | the output of the fo | llowing query | | |
| select coun union all s | t(*) from (selec select null as c) | t null as c as t group by | с | |
| Option 1 | Option 2 | Option 3 | Option 4 | |

NULL

empty

1

2

Question

• What is the output of the following query

```
select v, count(c) from (
select 1 as v, cast(null as int) as c union all
select 2 as v, cast(null as int) as c) as t group by v
```

| Option 1 | Option 2 | Option 3 | Ontion 4 |
|------------|------------------|---|----------|
| 1 0 2 0 | 1 NULL 2 NULL | $\begin{array}{ccc}1&1\\2&1\end{array}$ | Empty |

| Сι | istor | ner | | |
|----|-------|--------|------------------|--|
| | ID | Name | ShippingAddress | |
| | 1 | Holger | 180 Queens Gate | |
| | 2 Sam | | 32 Vassar Street | |
| | 3 | Peter | 180 Queens Gate | |
| | | | | |

select count(*) as c from Customer group by ShippingAddress having max(len(Name)) < 5 $\frac{c}{1}$

| Cι | istor | ner | |
|-------------------|--------|--------|-------------------------------------|
| | D Name | | ShippingAddress |
| 1 Holger 2 Sam | | Holger | 180 Queens Gate 32 Vassar Street |
| | | Sam | |
| | 3 | Peter | 180 Queens Gate |
| | | | |

| se | lect | ; * | from Customer |
|----|------|-------|-------------------|
| or | der | by | len(Name) |
| | ID | Name | ShippingAddress |
| | 2 | Sam | 32 Vassar Street |
| | 3 | Peter | 180 Queens Gate |
| | 1 | Holge | r 180 Queens Gate |

| ١D | Name | Shipping Address |
|----|--------|------------------|
| 1 | Holger | 180 Queens Gate |
| 2 | Sam | 32 Vassar Street |
| 3 | Peter | 180 Queens Gate |

Remember bag semantics!

• top n without order by is not properly defined

| Operators | |
|--------------|--|
| • left join | |
| • right join | |
| • in | |
| • some | |
| • all | |
| • exists | |

The End